

Hybrid Optimization for Time-Dependent Sequencing

Willem-Jan van Hoeve

Tepper School of Business
Carnegie Mellon University

www.andrew.cmu.edu/user/vanhoeve/



Joris Kinable

Robotics Institute & Tepper School of Business
Carnegie Mellon University



Andre A. Cire

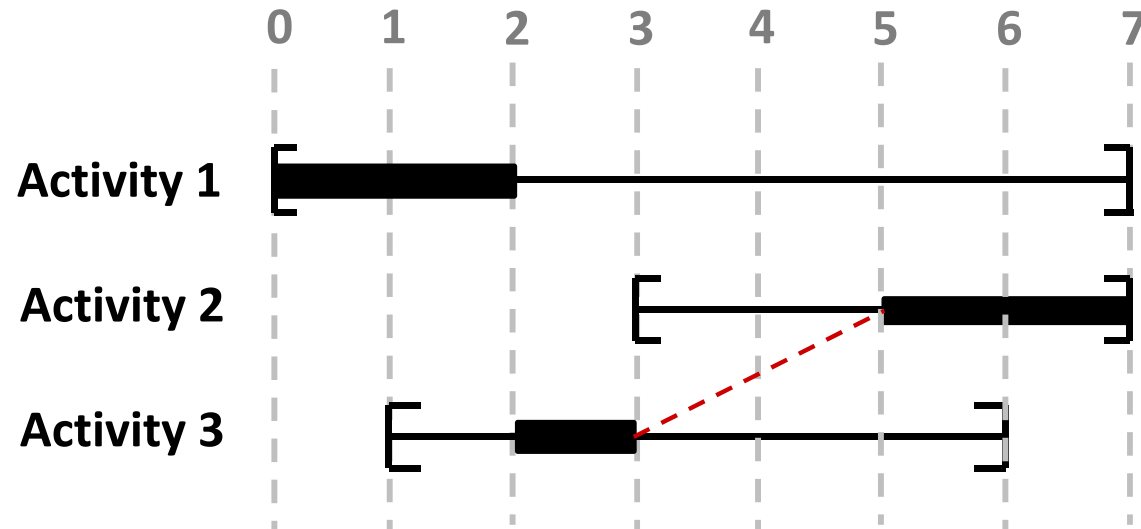
Rotman School of Management & U Toronto Scarborough
University of Toronto

- Time-dependent sequencing
 - machine scheduling, routing
- Challenging problem
 - best results with dedicated methods
 - not easy to extend with side constraints
- Utilize constraint programming framework?
 - strengthened constraint propagation with **MDDs**
 - improved bounds via **additive bounding** with LP
 - evaluate on TD-TSP and TD-SOP



- Activities

- processing time p_i
- released date r_i
- deadline d_i



- Resource

- non-preemptive
- process one activity at a time
- sequence-dependent setup times: **also depend on position!**

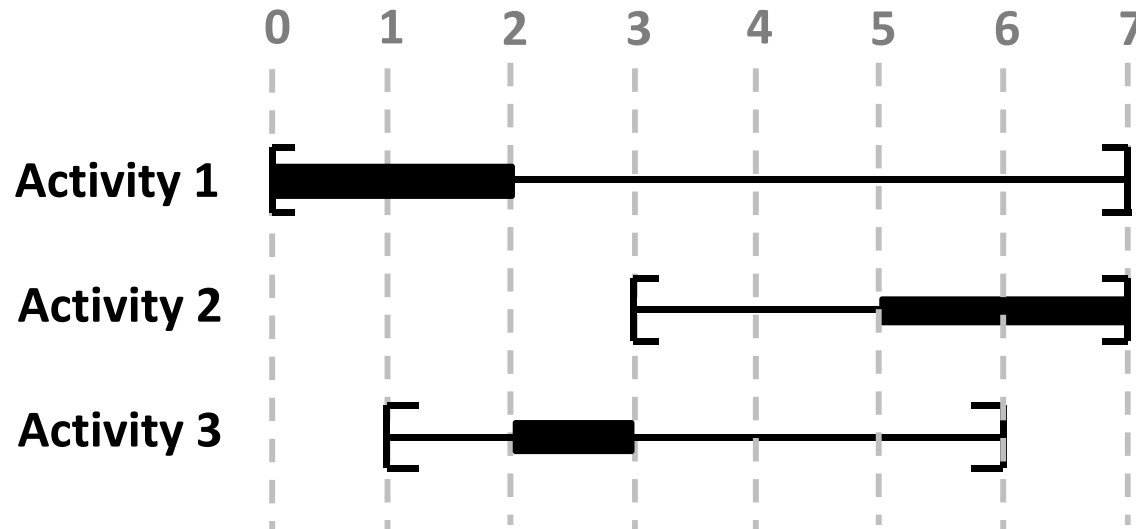
$\delta_{i,j}^t$ = setup time between i and j if i is at position t

- Precedence relations between activities
- Various objective functions
 - Sum of setup times
 - Makespan
 - (Weighted) sum of completion times
 - (Weighted) tardiness
 - number of late jobs
 - ...

- MIP-based
 - Picard and Queyranne (1978)
 - Gouveia and Voss (1995)
 - Abeledo et al. (2013)
 - ...
- Many more approaches to *time*-dependent TSP
 - Ichoua et al. (2003)
 - Cordeau et al. (2014)
 - Melgarejo et al. (2015)
 - ...

- Every solution can be written as a **permutation π**

$\pi_1, \pi_2, \dots, \pi_n$: activity sequencing



$$\pi_1 = 1, \pi_2 = 3, \pi_3 = 2$$

- Variables π_i : label of i^{th} activity in the sequence
 L_i : position of activity i in the sequence

$$\min \sum_{i=0}^n \delta_{\pi_i, \pi_{i+1}}^i$$

$$\text{s.t. AllDiff}(\pi_1, \dots, \pi_n)$$

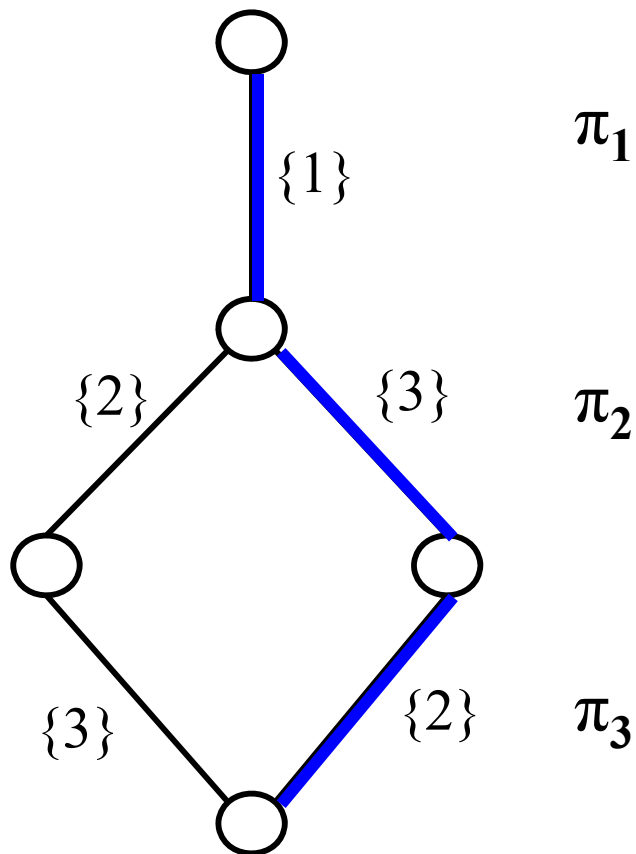
$$L_{\pi_i} = i \quad \forall i = 1, \dots, n$$

$$L_i < L_j \quad \forall (i \ll j) \in P$$

$$L_i \in \{1, \dots, n\} \quad \forall i = 1, \dots, n$$

$$\pi_i \in \{1, \dots, n\} \quad \forall i = 1, \dots, n$$

- **Weak model:** objective and AllDiff are decoupled



Act	r_i	p_i	d_i
1	0	2	3
2	4	2	9
3	3	3	8

Path {1} – {3} – {2} :

$$0 \leq \text{start}_1 \leq 1$$

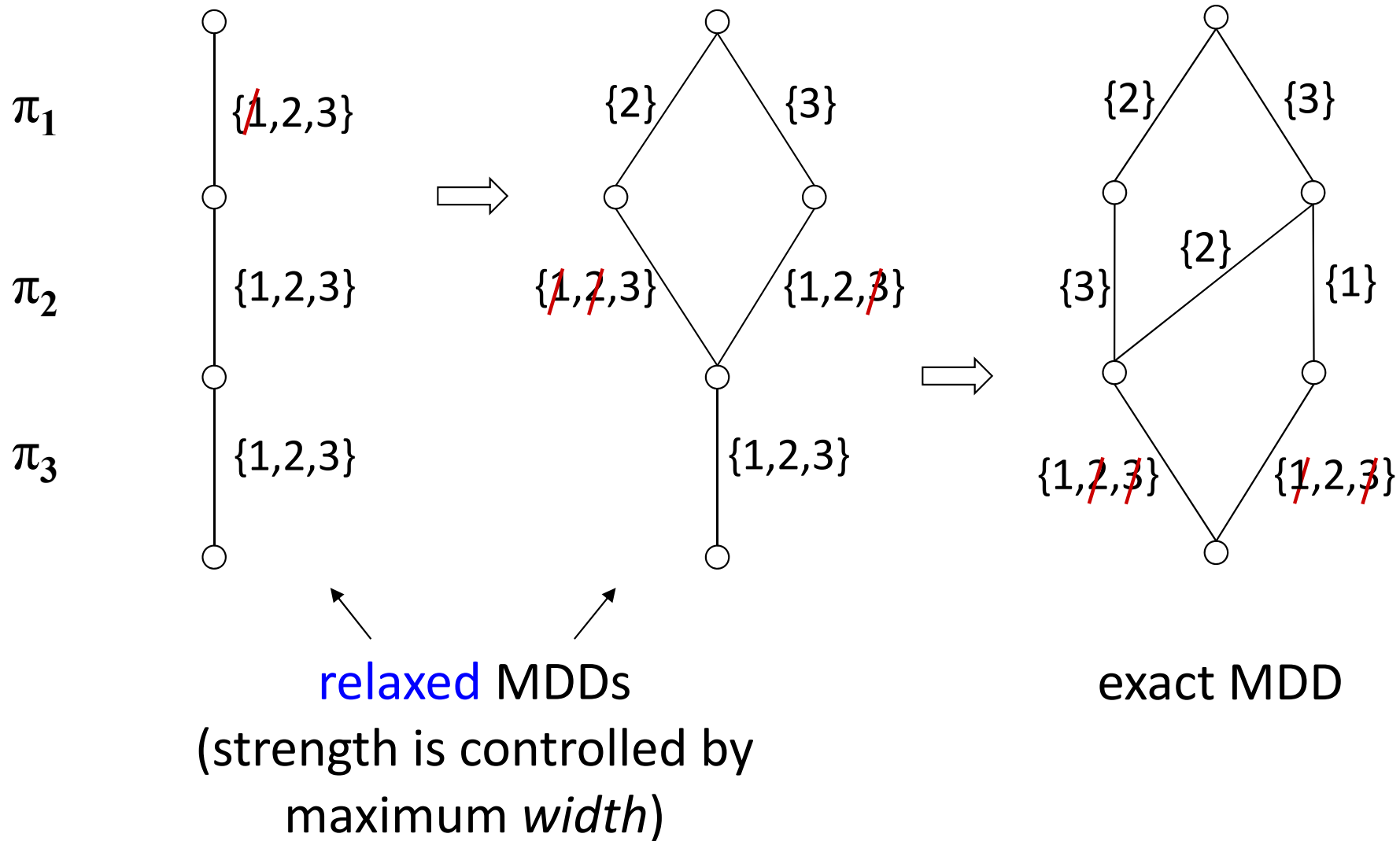
$$6 \leq \text{start}_2 \leq 7$$

$$3 \leq \text{start}_3 \leq 5$$

Cire and v.H. (2013)

Top-down MDD compilation

precedence: $3 \ll 1$



Propagation: remove infeasible arcs from the MDD

We can utilize several structures/constraints:

- *Alldifferent* for the permutation structure
- Precedence relations
- Earliest start time and latest end time

Propagating MDDs rather than variable domains can yield orders of magnitude speedup

Andersen et al. (2007), Hoda et al. (2010), Cire&v.H. (2013),
Bergman et al. (2015)

Propagation: remove infeasible arcs from the MDD

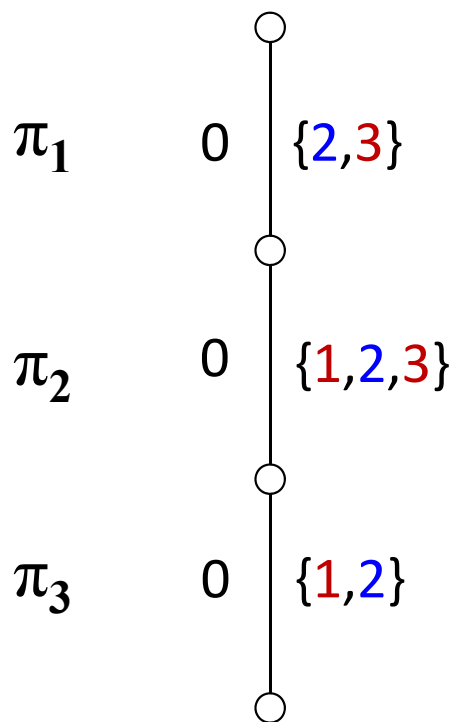
We can utilize several structures/constraints:

- *Alldifferent* for the permutation structure
- Precedence relations
- Earliest start time and latest end time
 - adapt rule: $\delta_{i,j}$ becomes $\delta_{i,j}^t$
- Also needed for **objective**
 - minimize sum of setup times

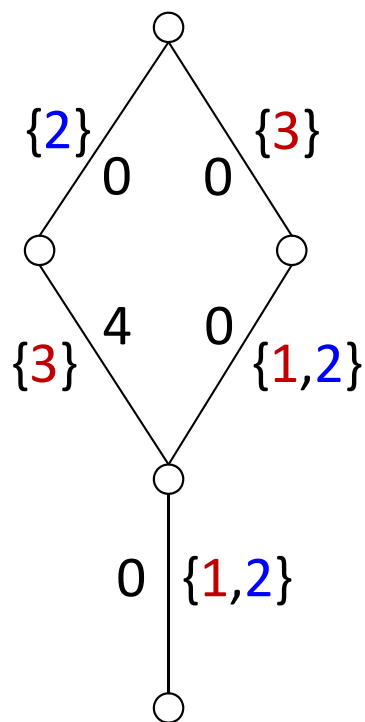
$$\begin{aligned} \min \quad & z \\ \text{s.t.} \quad & \text{AllDiff}(\pi_1, \dots, \pi_n) \\ & \text{MDDconstr}(\pi_1, \dots, \pi_n, W, z, \delta^t, P) \\ & L_{\pi_i} = i \quad \forall i = 1, \dots, n \\ & L_i < L_j \quad \forall (i \ll j) \in P \\ & L_i \in \{1, \dots, n\} \quad \forall i = 1, \dots, n \\ & \pi_i \in \{1, \dots, n\} \quad \forall i = 1, \dots, n \\ & z \in \{0, \dots, \infty\} \end{aligned}$$

Stronger model: objective handled within MDD constraint

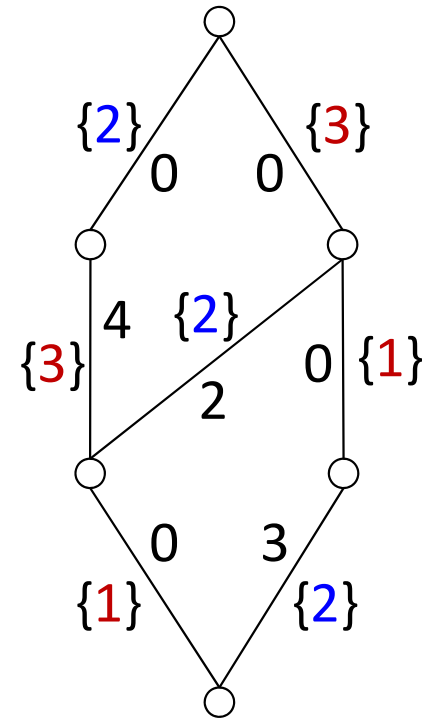
Bounds from relaxed MDDs



objective: 0



objective: 0



objective: 3

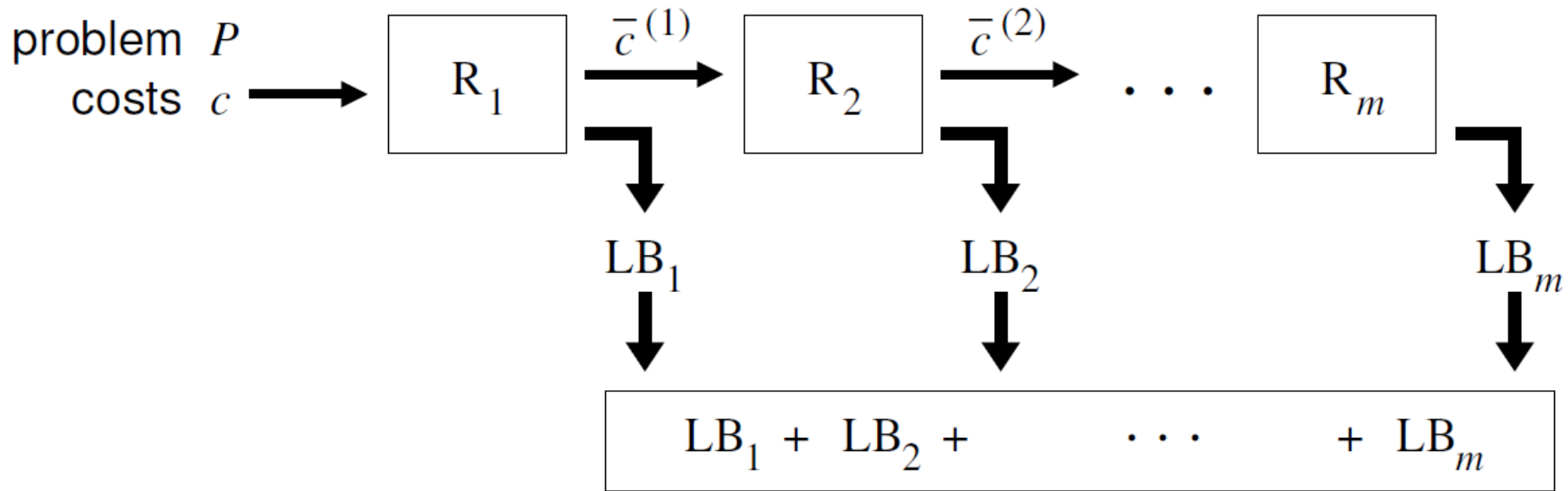
act 1: red
act 2: blue
act 3: red

$$\delta_{\text{red,blue}}^1 = 2$$

$$\delta_{\text{red,blue}}^2 = 3$$

$$\delta_{\text{blue,red}}^1 = 4$$

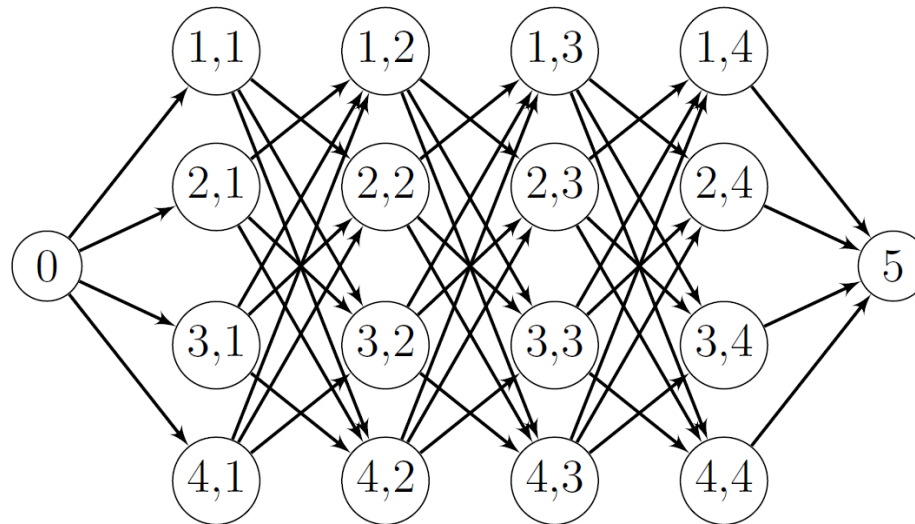
$$\delta_{\text{blue,red}}^2 = 6$$



(Fischetti & Toth, 1989)

Add LP reduced costs to MDD relaxation

- Continuous LP relaxation 'discretized' through MDD
- Stronger bounds
- Improved cost-based filtering



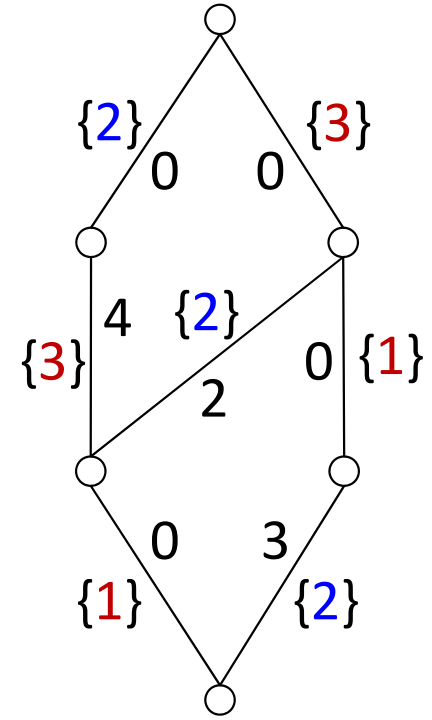
- Time-space network model (Picard & Queyranne, 1978)

- Variables

$$x_{i,j}^t = \begin{cases} 1 & \text{if } i \text{ is performed at } t \text{ and followed by } j \\ 0 & \text{otherwise} \end{cases}$$

- Constraints: flow conservation; perform each activity

- Approach
 - solve LP relaxation
 - in MDD, replace $\delta_{i,j}^t$ with $\bar{c}_{i,j}^t$
- Since MDD is relaxation, shortest path is valid bound



- Time-dependent TSP benchmark
 - 38 instances from TSPLIB (14-107 jobs)
 - $\delta_{i,j}^t = (n-t) * \delta_{i,j}$
- Time limit: 30 minutes
- IBM ILOG CPLEX and CP Optimizer 12.4
- MDD added to CP Optimizer (Cire & v.H., 2013)
 - maximum width 512

- Compare root node bound improvement

	LP	MDD	MDD+LP	percentage improvement	
				w.r.t. LP	w.r.t. MDD
berlin52	112,350.0	49,056.0	119,694.0	6.54%	143.99%
dTSP50.0	10,465.7	5,537.0	10,646.0	1.72%	92.27%
kroA100	693,870.0	223,039.0	719,992.0	3.76%	222.81%
pr76	2,496,050.0	2,116,910.0	2,679,143.0	7.34%	26.56%
Average improvement (38 instances):				5.96%	103.34%

	#Solved	Avg end gap*
(Picard & Queyrrane, 1978) (MIP)	6/38	33.37%
(Gouveia & Voss, 1995) (MIP)	6/38	31.64%
(Abeledo et al., 2013) (BPC**)	35/38	1.64%
Pure CP	0/38	45.02%
CP + MDD	7/38	7.42%
CP + MDD + Additive Bounding	12/38	6.49%

* Average end gap w.r.t. overall best bound

** Dedicated method; much longer time limit

Sequential Ordering Problem

	#Solved	Avg end gap
(Picard & Queyrrane, 1978) (MIP)	6/30	29.85%
(Gouveia & Voss, 1995) (MIP)	6/30	29.17%
Pure CP	5/30	25.68%
CP + MDD + Additive Bounding	10/30	21.22%

On average, additive MDD+LP bound improves

- LP root node bound by 51.41%
- MDD root node bound by 9.54%

- Hybrid optimization method for time-dependent sequencing
 - CP framework
 - MDD relaxation for improved propagation
 - Additive bounding with LP for stronger bounds
 - Side constraints are easily added
- Experiments
 - Competitive generic approach