

# 15-317 Lecture 12: Theorem proving

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- Proof search (and related problems)
- Reduced sequent calculus
- Contraction-free sequent calculus

# Proof Search

- Given  $A$ , find a proof of  $\Rightarrow A$
  - Or justify that  $\Rightarrow A$  is not provable
  - Second idea: Given  $\Gamma, A$ ,  
prove  $\Gamma \Rightarrow A$  or justify  $\Gamma \not\Rightarrow A$
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Similar problems:

- Given a proof, find prop./sequent
- Given a proof and a sequent: Do they match?

Proof	Sequent
Find	Given
Given	Given
Given	Find

Proof search / Program synthesis

Proof checking / Type checking

Type inference

# Reduced Sequent Calculus

- Ideally we never apply the same rule to the same prop

⋮

$$\frac{\Gamma, A, A, A, A \wedge B \Rightarrow C}{\Gamma, A \wedge B, A, A \Rightarrow C} \wedge L_1$$
$$\frac{\Gamma, A \wedge B, A, A \Rightarrow C}{\Gamma, A \wedge B, A \Rightarrow C} \wedge L_1$$
$$\frac{\Gamma, A \wedge B, A \Rightarrow C}{\Gamma, A \wedge B \Rightarrow C} \wedge L_1$$

$$\frac{\Gamma \rightarrow A \quad \Gamma \rightarrow B}{\Gamma \rightarrow A \wedge B} \wedge R$$

$$\frac{\Gamma, A, B \rightarrow C}{\Gamma, A \wedge B \rightarrow C} \wedge L$$

Ensure  $\Gamma \Rightarrow A$  iff  $\Gamma \rightarrow A$

Soundness of  $\rightarrow$  w.r.t.  $\Rightarrow$

$$\Gamma, A, B \Rightarrow C$$

$$\Gamma, \overline{A \wedge B}, \overline{A}, \overline{B} \Rightarrow C \quad \text{- weakening} \Rightarrow$$

$$\Gamma, A \wedge B, A \Rightarrow C$$

$$\Gamma, A \wedge B \Rightarrow C$$

Completeness of  $\rightarrow$  w.r.t.  $\Rightarrow$

$$\text{If } \mathcal{D} = \frac{\Gamma, A \wedge B, A \Rightarrow C}{\Gamma, A \wedge B \Rightarrow C} \wedge I,$$

$$\overline{\Gamma}, \overline{A}, \overline{B} \rightarrow A \quad \text{id}_A$$

$$\overline{\Gamma}, \overline{A \wedge B} \rightarrow A \quad \wedge I$$

$$\overline{\Gamma}, \overline{A \wedge B} \rightarrow \overline{C}$$

i.h. ( $\mathcal{D}_1$ )

$$\Gamma, A \wedge B, A \rightarrow C$$

-cut<sub>A</sub>

T/F

\_\_\_\_\_ FL  
 $\Gamma, F \rightarrow C$

\_\_\_\_\_ TR  
 $\Gamma \rightarrow T$

$\Gamma \rightarrow C$   
\_\_\_\_\_ TL  
 $\Gamma, T \rightarrow C$

# Disjunction

$$\frac{\Gamma \rightarrow A}{\Gamma \rightarrow A \vee B} \vee R_1 \quad \frac{\Gamma \rightarrow B}{\Gamma \rightarrow A \vee B} \vee R_2$$

$$\frac{\Gamma, A \rightarrow C \quad \Gamma, B \rightarrow C}{\Gamma, A \vee B \rightarrow C} \vee L$$

$$\frac{\Gamma, A \rightarrow C \quad \Gamma, B \rightarrow C}{\Gamma, A \vee B \rightarrow C} \vee L$$

$$\mathcal{D} = \frac{\frac{\Gamma, A \vee B, A \Rightarrow C}{\Gamma, A \vee B \Rightarrow C} \mathcal{D}_1 \quad \Gamma, A \vee B, B \Rightarrow C}{\Gamma, A \vee B \Rightarrow C} \vee L$$

$$\frac{\frac{\Gamma, A \rightarrow A}{\Gamma, A \rightarrow A \vee B} \vee R, \text{ i.h. } (\mathcal{D}_1) \quad \Gamma, A \vee B, A \rightarrow C}{\Gamma, A \rightarrow C} \text{cut}_A \quad \Gamma, B \rightarrow C}{\Gamma, A \vee B \rightarrow C} \vee L$$



# Implication

$$\frac{\Gamma, A \rightarrow B}{\Gamma \rightarrow A \supset B} \supset R$$

$$\frac{\Gamma, A \supset B \rightarrow A \quad \Gamma, B \rightarrow C}{\Gamma, A \supset B \rightarrow C} \supset L$$

# Properties of $\rightarrow$

- Weakening: If  $\Gamma \rightarrow C$ , then  $\Gamma, A \rightarrow C$

- Contraction: If  $\Gamma, A, A \rightarrow C$  then  $\Gamma, A \rightarrow C$

- Identity:  $\Gamma, A \rightarrow A$

- Cut: If  $\Gamma \rightarrow A$  and  $\Gamma, A \rightarrow C$ , then  $\Gamma \rightarrow C$

- Soundness: If  $\Gamma \rightarrow A$  then  $\Gamma \Rightarrow A$

- Completeness: If  $\Gamma \Rightarrow A$  then  $\Gamma \rightarrow A$

$$\frac{\Gamma, A \supset B \rightarrow A \quad \Gamma, B \rightarrow C}{\Gamma, A \supset B \rightarrow C} \supset I$$

Consider cases of what  $A$  can be

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$$A = T$$

$$\frac{[\Gamma, T \supset B \rightarrow T]^{TR} \quad \Gamma, B \rightarrow C}{\Gamma, T \supset B \rightarrow C} T \supset I$$

$$\frac{\Gamma, B \rightarrow C}{\Gamma, T \supset B \rightarrow C} T \supset I$$

$$A = F$$

$$\frac{\Gamma, F \supset B \rightarrow F \quad \Gamma, B \rightarrow C}{\Gamma, F \supset B \rightarrow C} \text{FOL}$$

$$\frac{\Gamma, B \rightarrow C}{\Gamma, F \supset B \rightarrow C} \text{FOL}$$

$$\Gamma, F \supset B \rightarrow F \rightarrow \Gamma, T \rightarrow F$$

$$\frac{\Gamma \rightarrow C}{\Gamma, F \supset B \rightarrow C} \text{FOL}$$

          
T

$$A = A_1 \vee A_2$$

$$\frac{\Gamma, (A_1 \vee A_2) \supset B \rightarrow A_1 \vee A_2 \quad \Gamma, B \rightarrow C}{\Gamma, (A_1 \vee A_2) \supset B \rightarrow C} \vee \supset L$$

$$\Gamma, (A_1 \vee A_2) \supset B \rightarrow C$$



$$(A_1 \vee A_2) \supset B \equiv (A_1 \supset B) \wedge (A_2 \supset B)$$



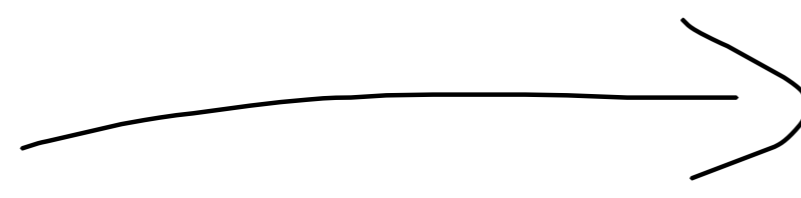
$$\frac{\Gamma, A_1 \supset B, A_2 \supset B \rightarrow A_1 \vee A_2 \quad \Gamma, B \rightarrow C}{\Gamma, (A_1 \vee A_2) \supset B \rightarrow C} \vee \supset L$$

$$\Gamma, (A_1 \vee A_2) \supset B \rightarrow C$$

$\frac{\Gamma, A_1 \supset B, A_2 \supset B \rightarrow C}{\Gamma, (A_1 \vee A_2) \supset B \rightarrow C} \vee \supset L$
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$$A = P$$

$$\frac{\Gamma, P \supset B \rightarrow P \quad \Gamma, B \rightarrow C}{\Gamma, P \supset B \rightarrow C}$$



$$\frac{P \in \Gamma \quad \Gamma, B \rightarrow C}{\Gamma, P \supset B \rightarrow C} P \supset L$$

$$A = A_1 \wedge A_2$$

$$\frac{\Gamma, A_1 \supset (A_2 \supset B) \rightarrow C}{\Gamma, (A_1 \wedge A_2) \supset B \rightarrow C}$$

- Justified by letting  
 $\wedge$  have higher weight than  $\supset$   
in the induction.

$$A = A_1 \supset A_2$$

$$A_1 \wedge ((A_1 \supset A_2) \supset B)$$

III

$$A_1 \wedge (A_2 \supset B)$$

$$\Gamma, A_2 \supset B \rightarrow A_1 \supset A_2$$

[

$$\Gamma, A_2 \supset B, A_1 \rightarrow A_2$$

$$\Gamma, B \rightarrow C$$

]

$$\Gamma, (A_1 \supset A_2) \supset B \rightarrow C$$

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$$\frac{\Gamma, (A_1 \supset A_2) \supset B, A_1 \rightarrow A_2}{\Gamma, (A_1 \supset A_2) \supset B \rightarrow A_1 \supset A_2} \supset R$$

$$\Gamma, (A_1 \supset A_2) \supset B \rightarrow A_1 \supset A_2$$

$$\Gamma, B \rightarrow C$$

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$$\Gamma, (A_1 \supset A_2) \supset B \rightarrow C$$



# Induction notes

- Multiset ordering:

Given a set  $A$  with an order  $\leq$ .

Suppose  $M_1, M_2$  are multisets of elements of  $A$ .

We say  $M_1 \leq M_2$  if

-  $M_1 \subseteq M_2$

-  $M_2 = M_2', a$  and  $M_1 = M_1' \cup \{b_1, \dots, b_n\}$ ,

where each  $b_i < a$  and  $M_1' \leq M_2'$