

The Case for Explainability of Real-Time Systems and their Analyses

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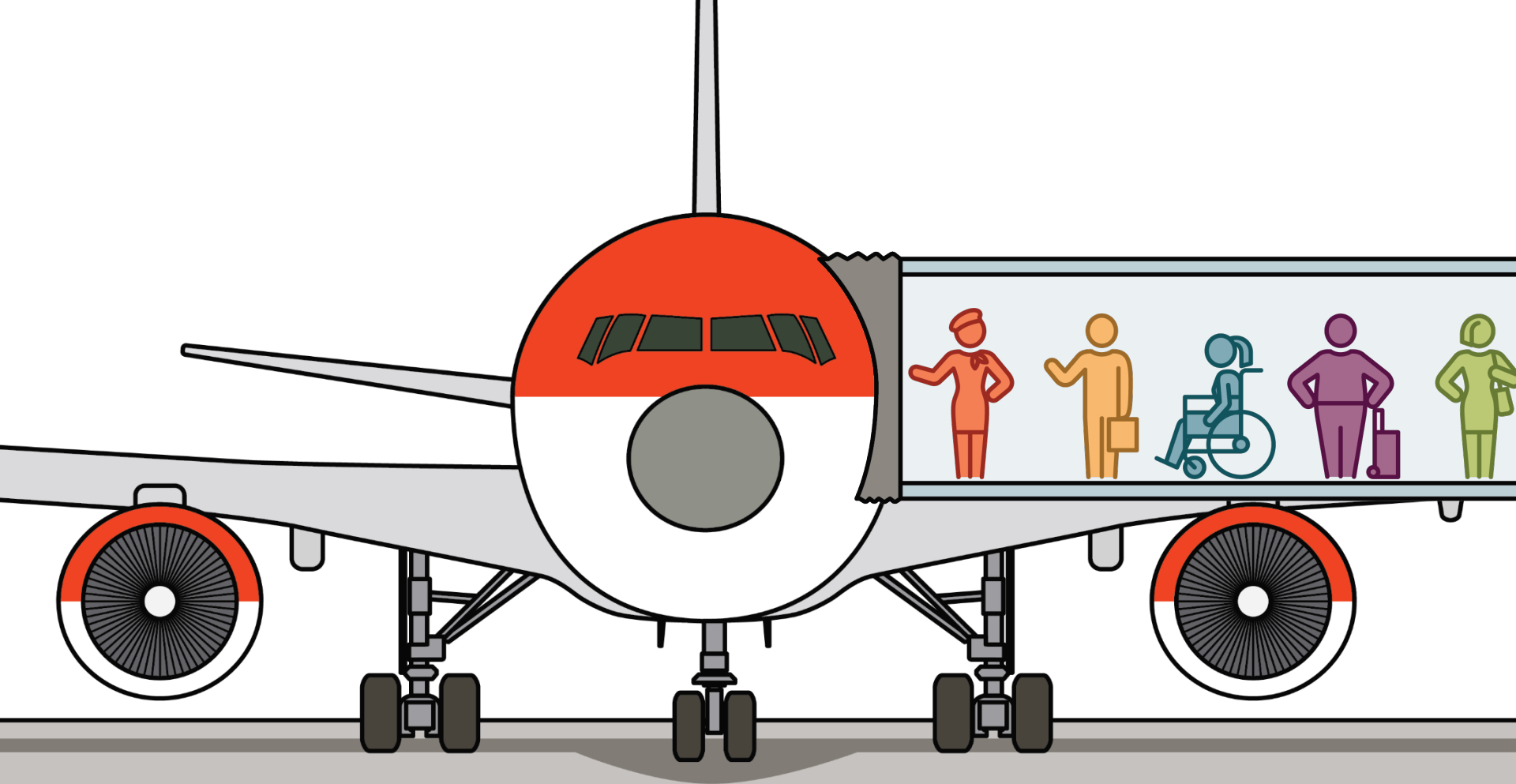
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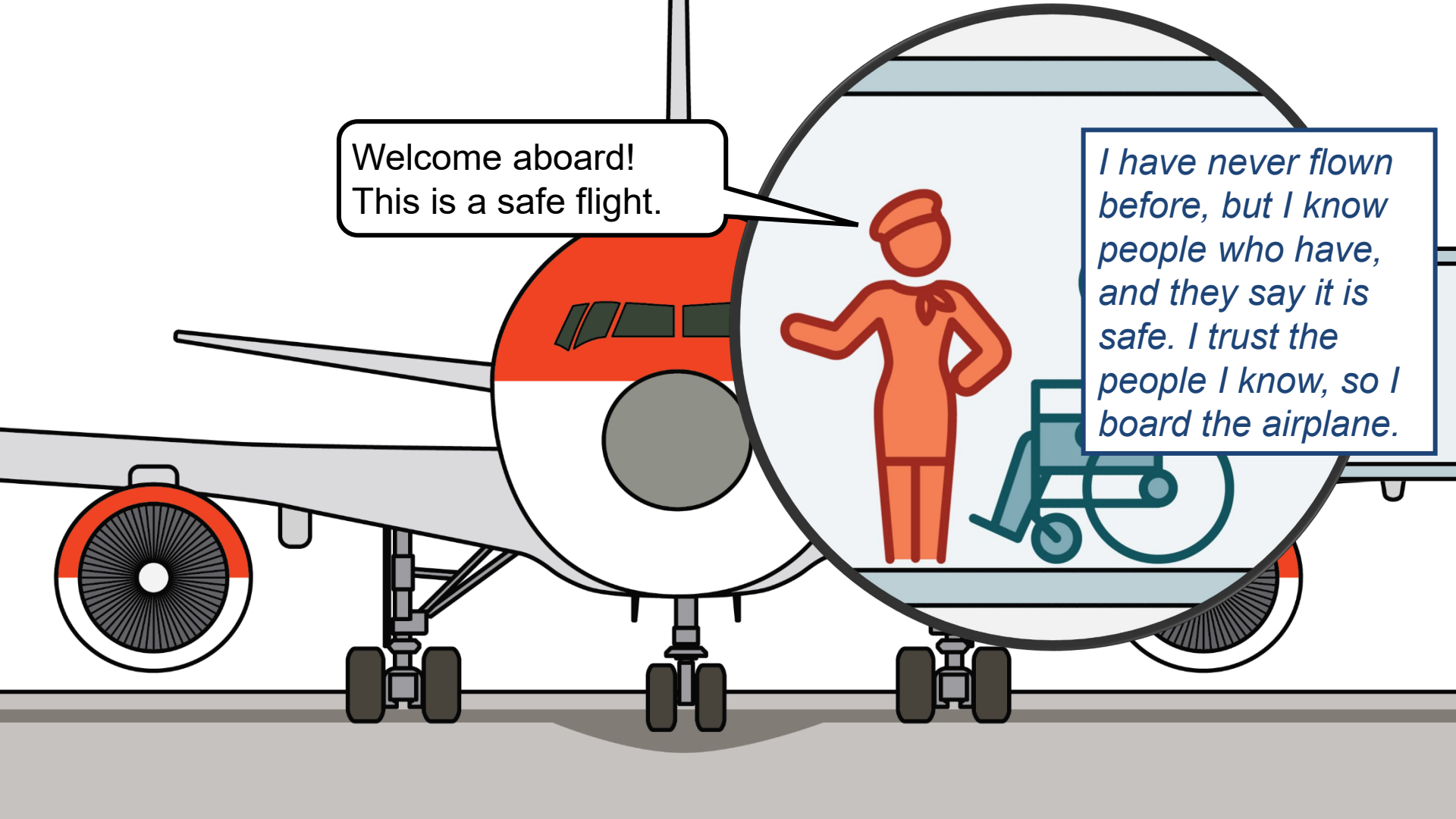
Motivation





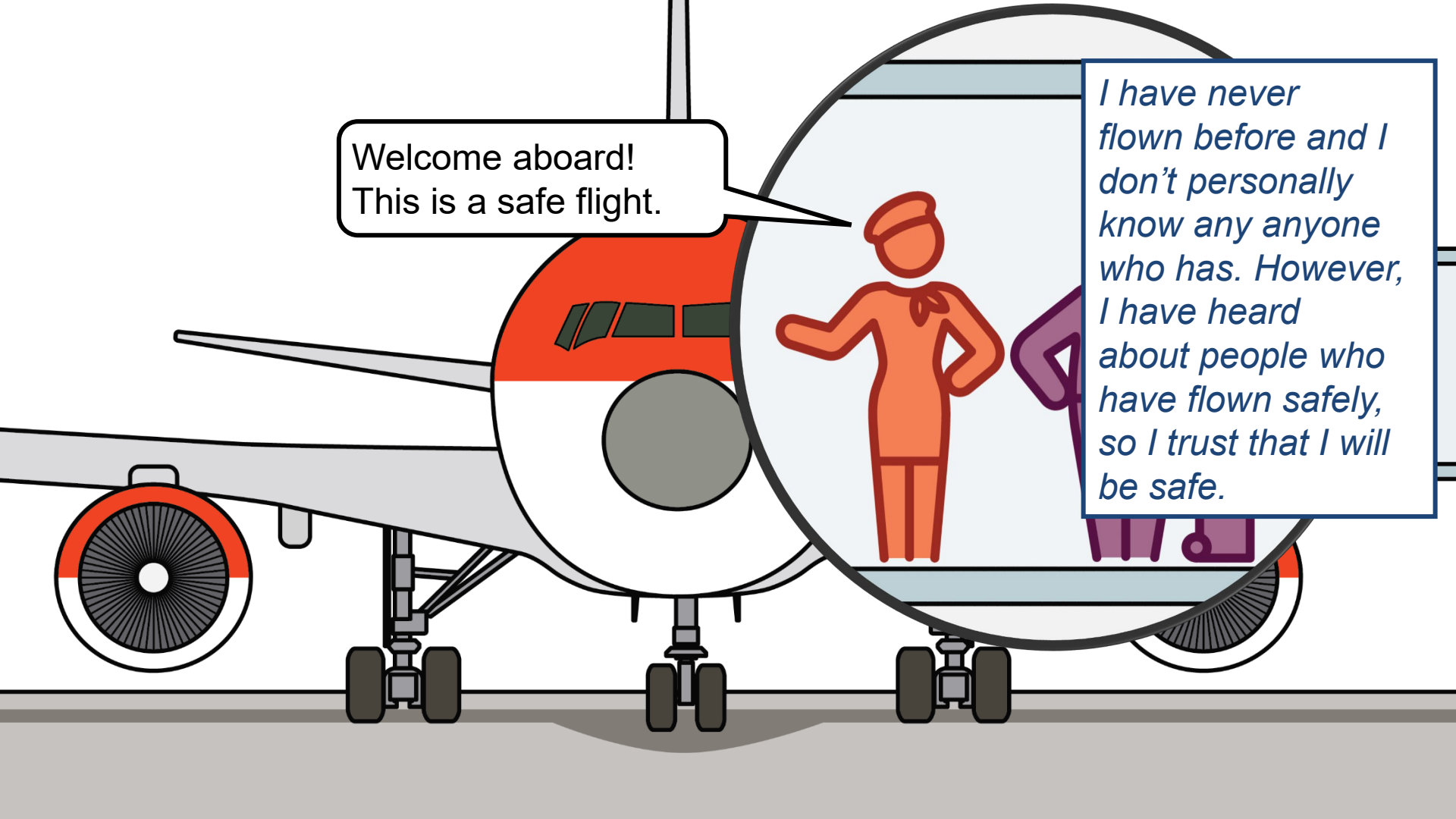
Welcome aboard!
This is a safe flight.

*I have flown many
times, so I trust
that this flight will
be safe. I board
the airplane.*



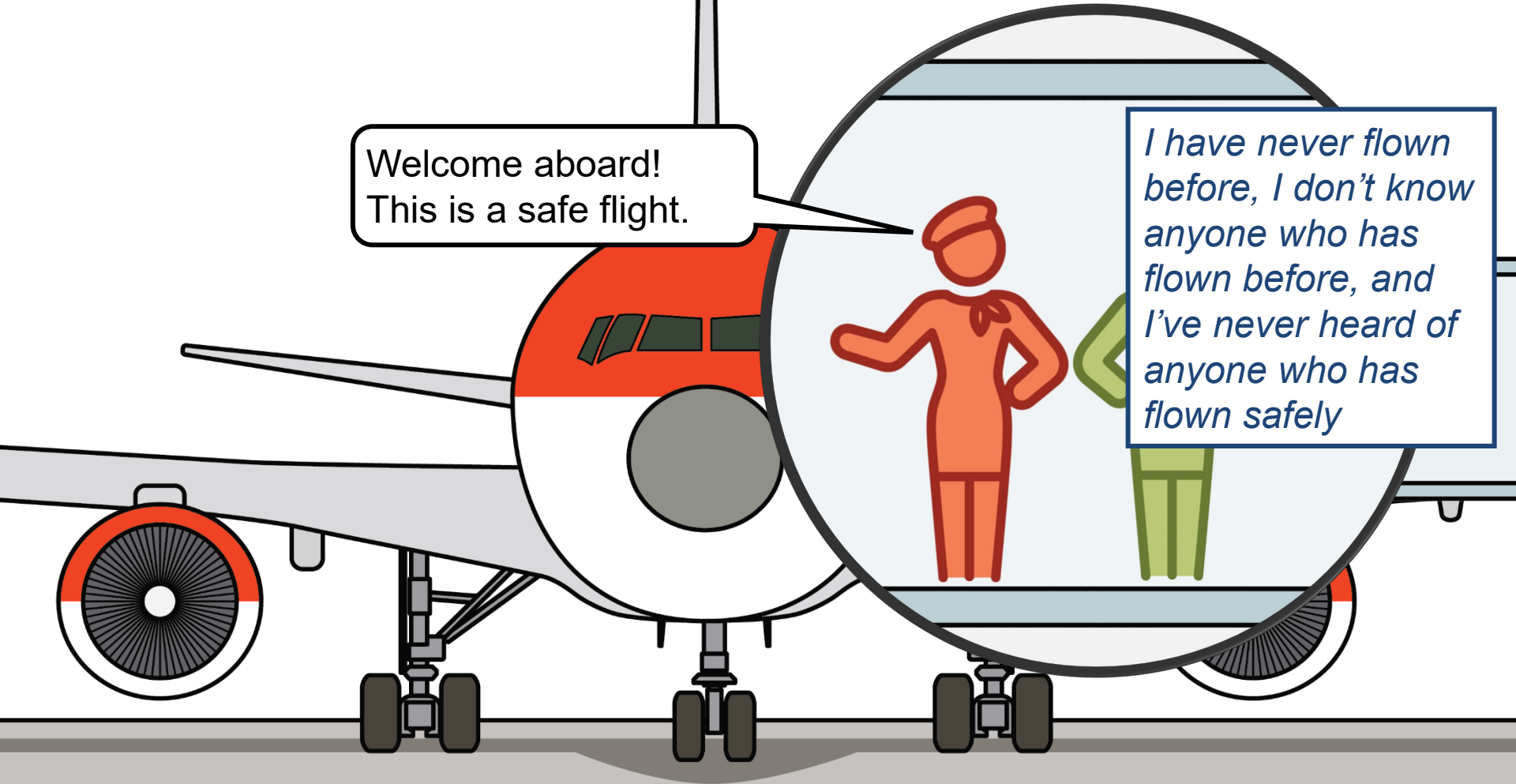
Welcome aboard!
This is a safe flight.

*I have never flown
before, but I know
people who have,
and they say it is
safe. I trust the
people I know, so I
board the airplane.*



Welcome aboard!
This is a safe flight.

*I have never
flown before and I
don't personally
know any anyone
who has. However,
I have heard
about people who
have flown safely,
so I trust that I will
be safe.*

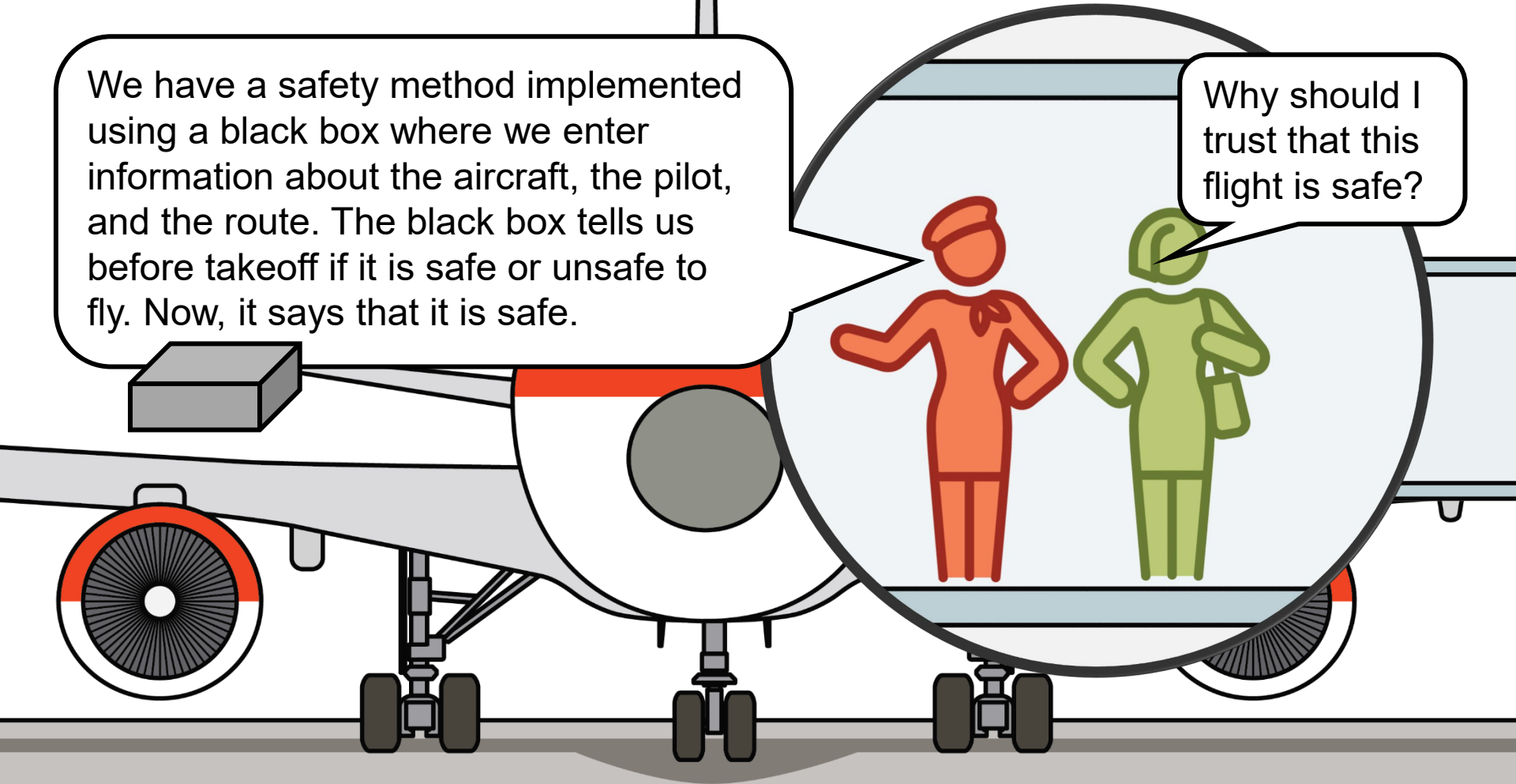


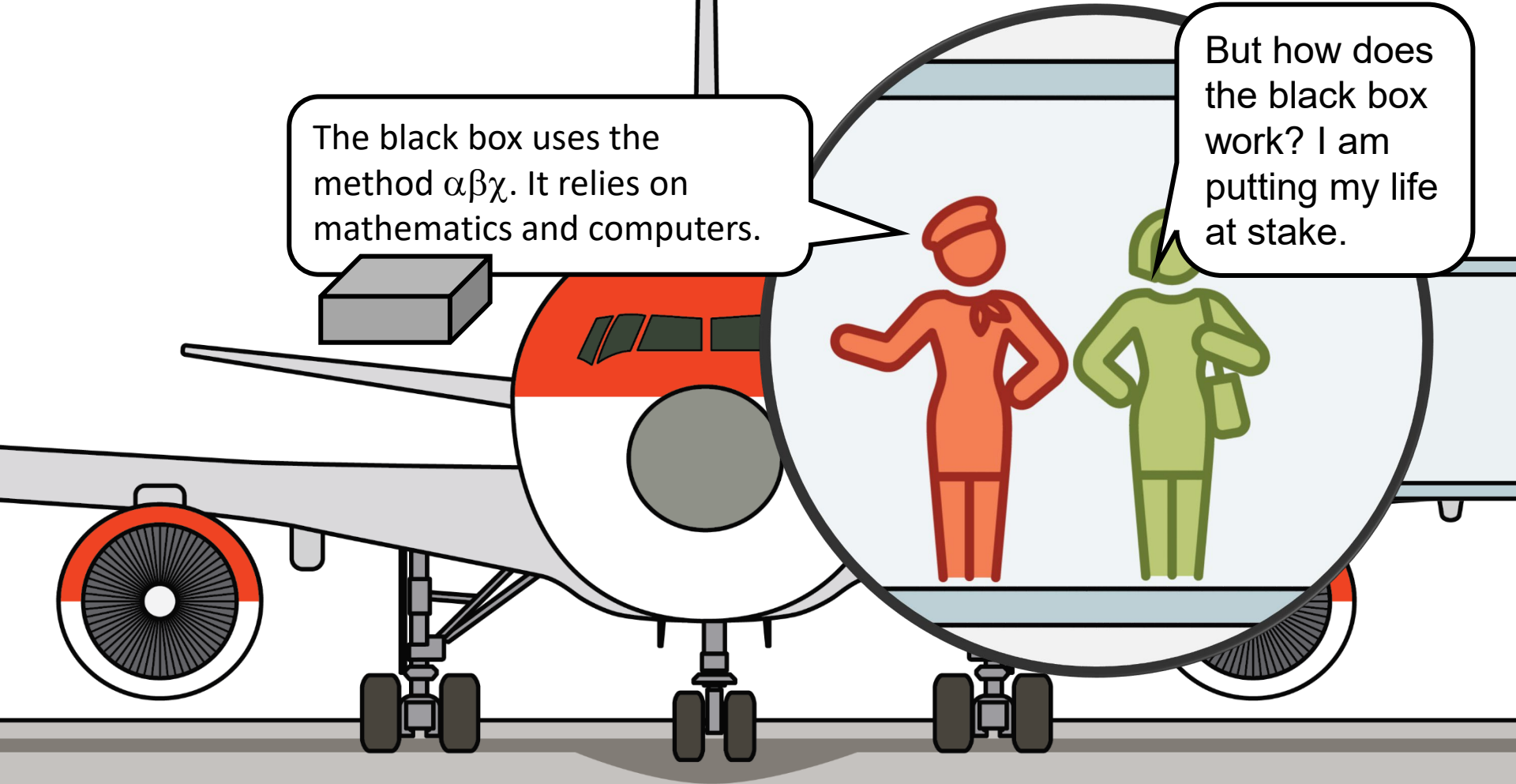
Welcome aboard!
This is a safe flight.

*I have never flown
before, I don't know
anyone who has
flown before, and
I've never heard of
anyone who has
flown safely*

We have a safety method implemented using a black box where we enter information about the aircraft, the pilot, and the route. The black box tells us before takeoff if it is safe or unsafe to fly. Now, it says that it is safe.

Why should I trust that this flight is safe?

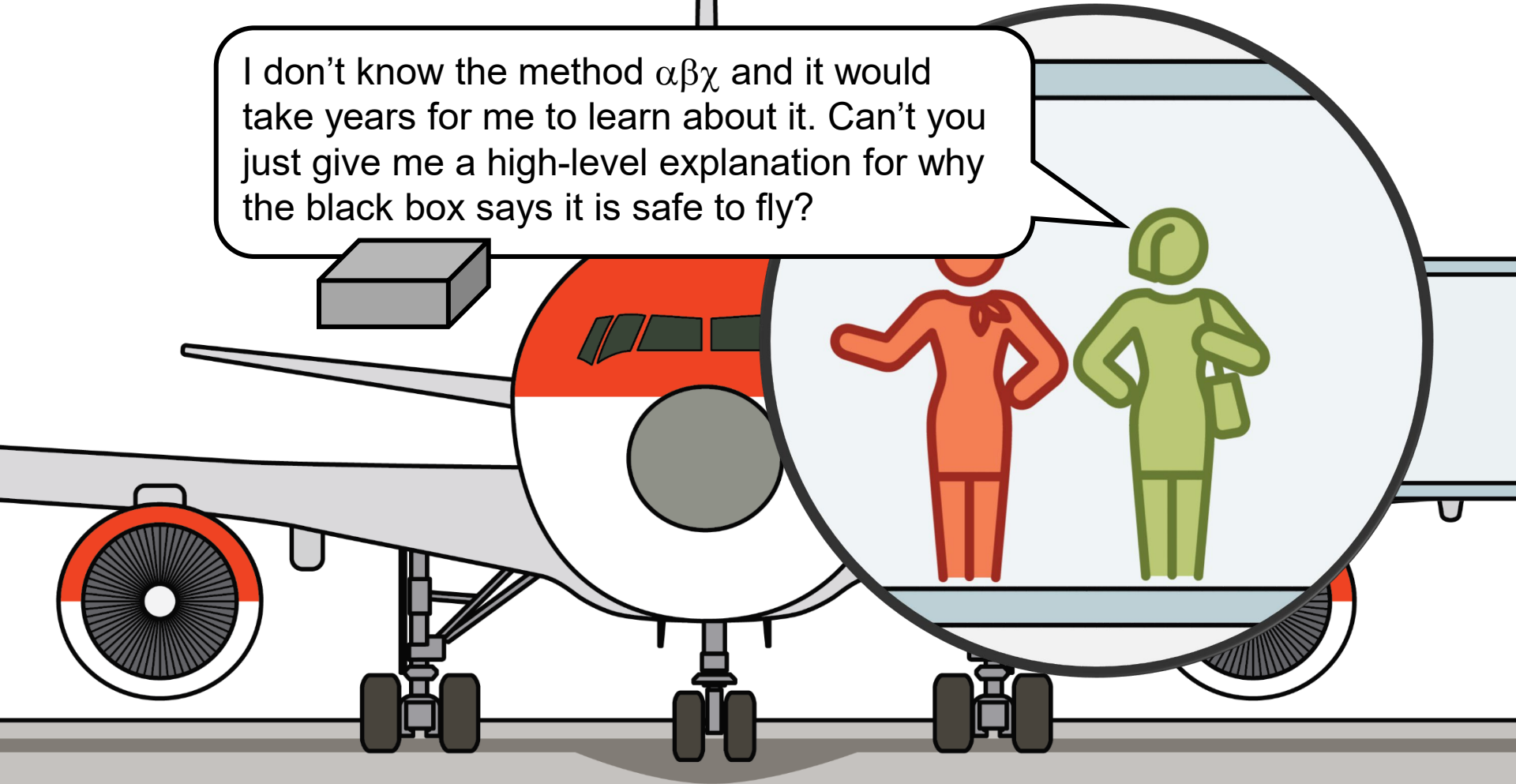




The black box uses the method $\alpha\beta\chi$. It relies on mathematics and computers.

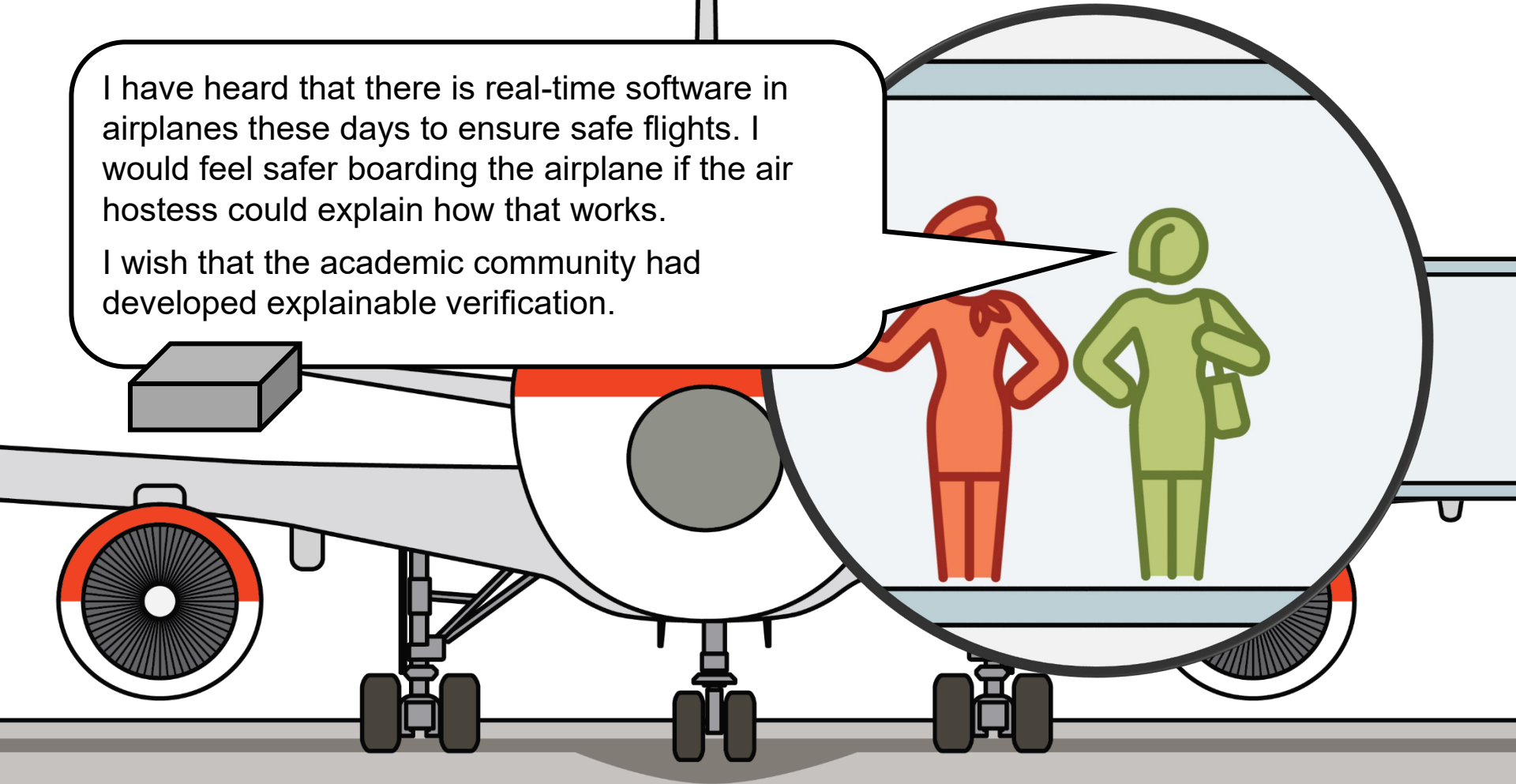
But how does the black box work? I am putting my life at stake.

I don't know the method $\alpha\beta\chi$ and it would take years for me to learn about it. Can't you just give me a high-level explanation for why the black box says it is safe to fly?

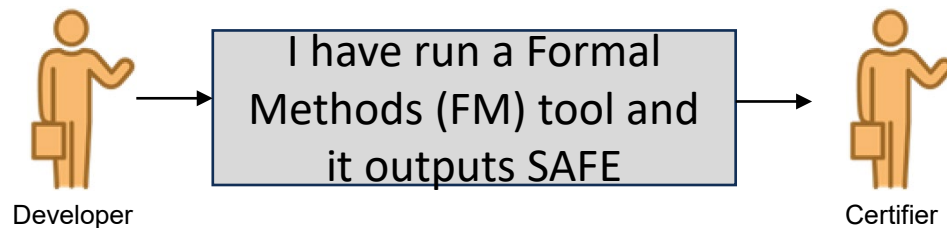


I have heard that there is real-time software in airplanes these days to ensure safe flights. I would feel safer boarding the airplane if the air hostess could explain how that works.

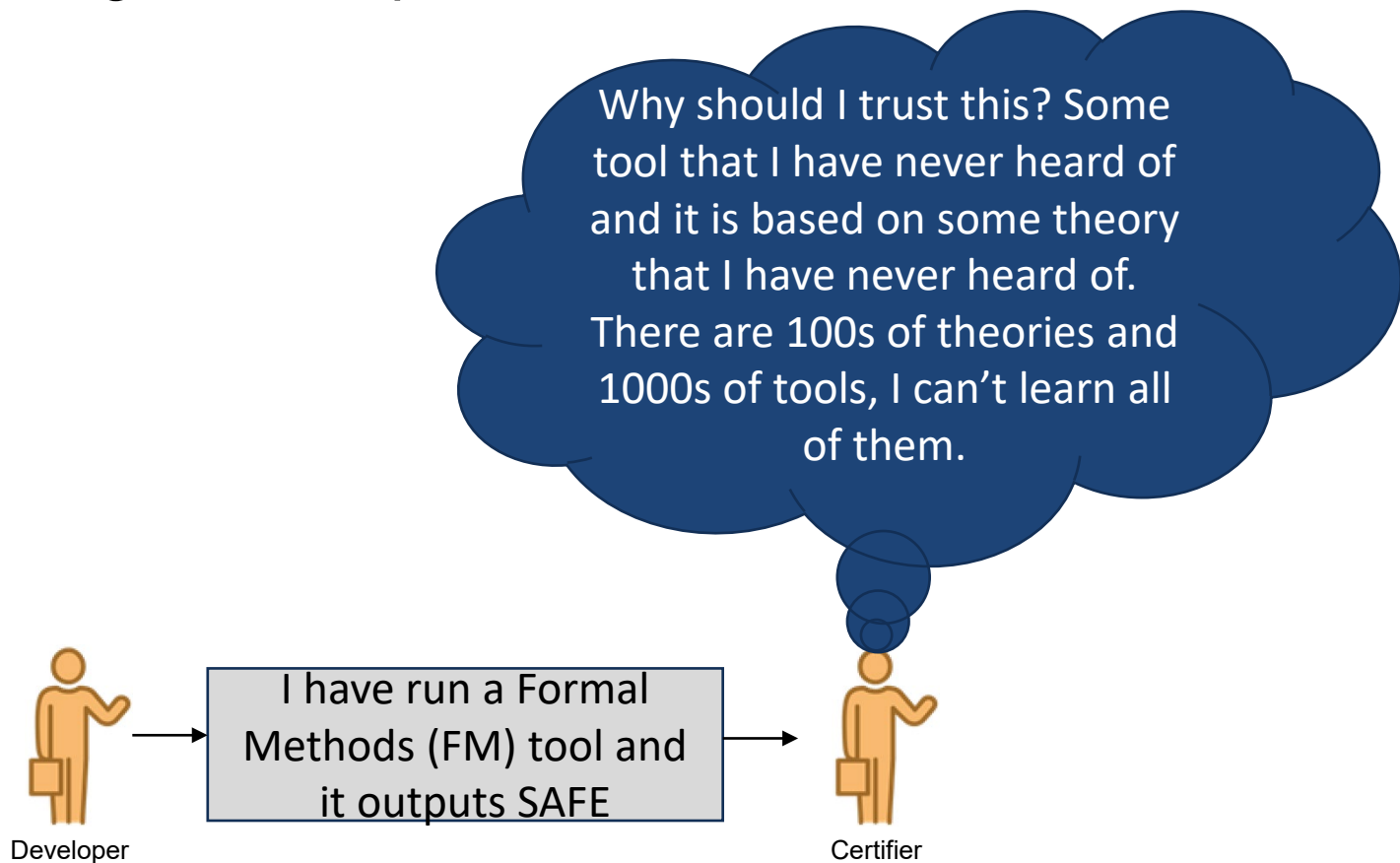
I wish that the academic community had developed explainable verification.



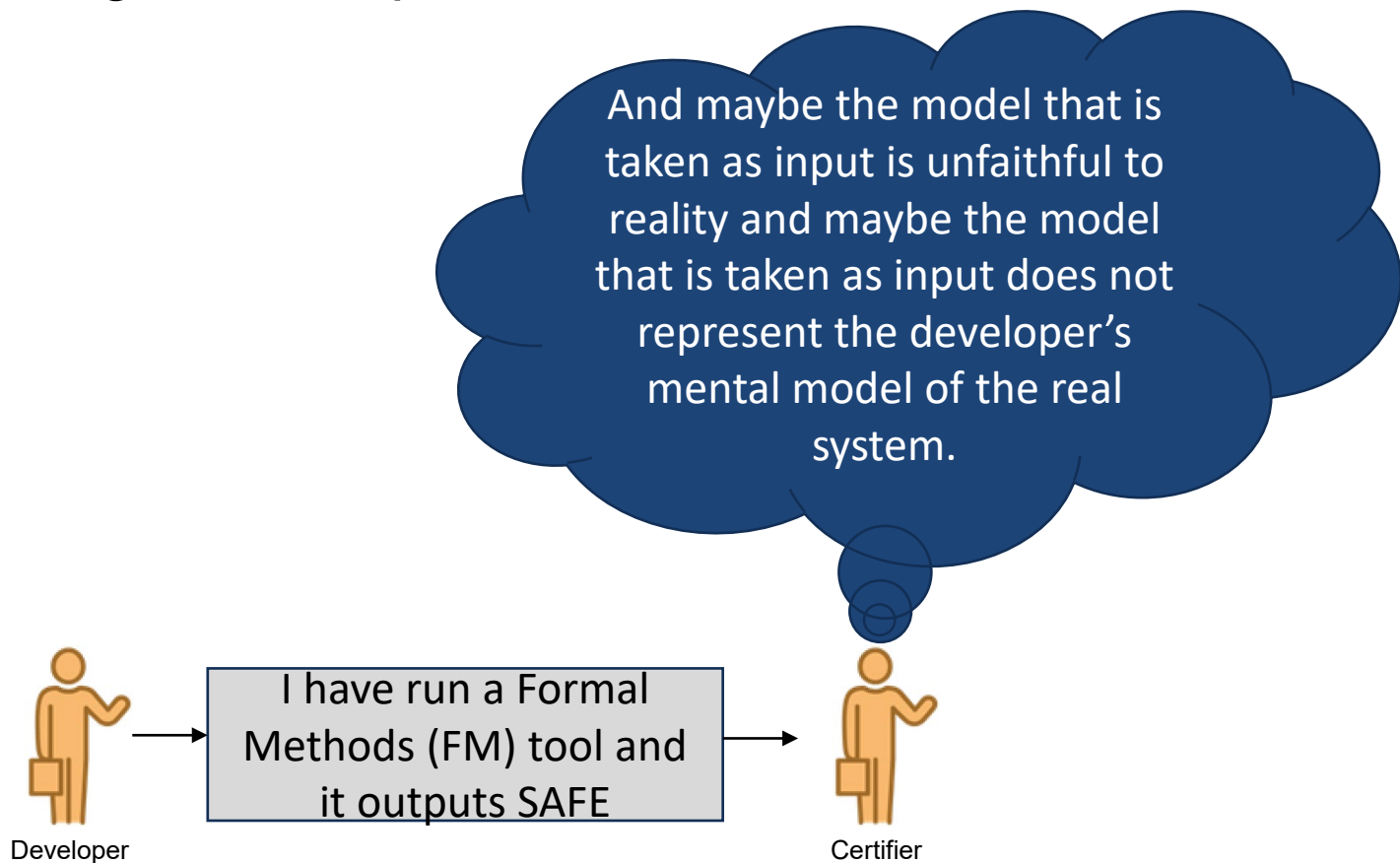
Problem: How to get developers and certifiers to trust FMs



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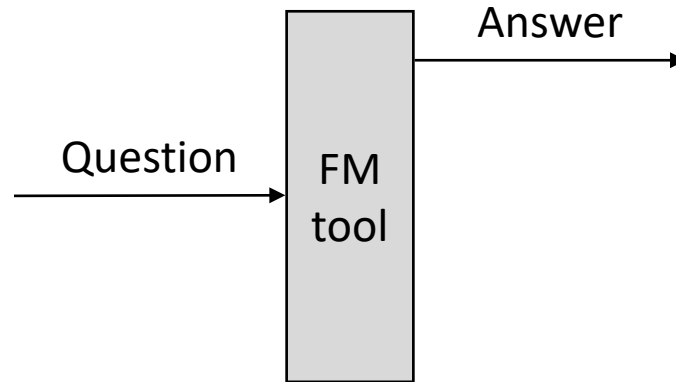




Explanation of Analysis

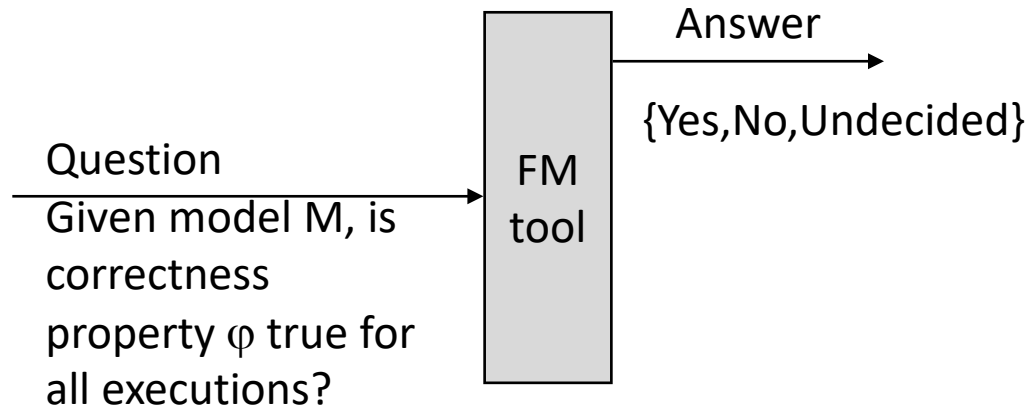


The Present: A World Without Explainability





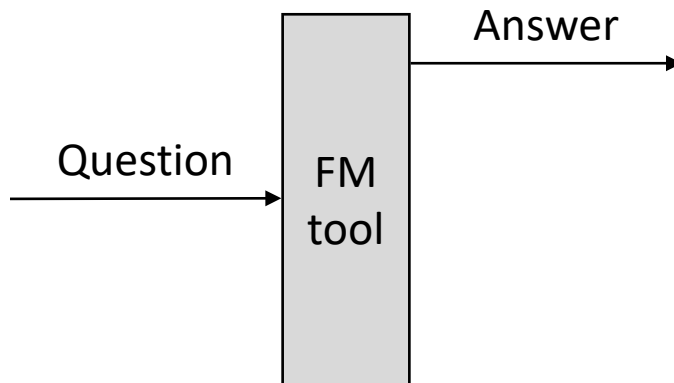
The Present: A World Without Explainability





The Present: A World Without Explainability

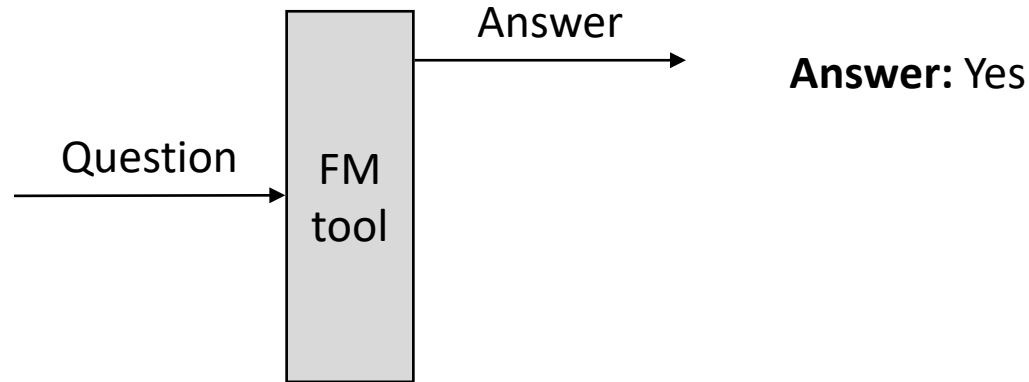
Question: Given taskset τ , the following run-time scheduler S , and assumptions A , is it schedulable?





The Present: A World Without Explainability

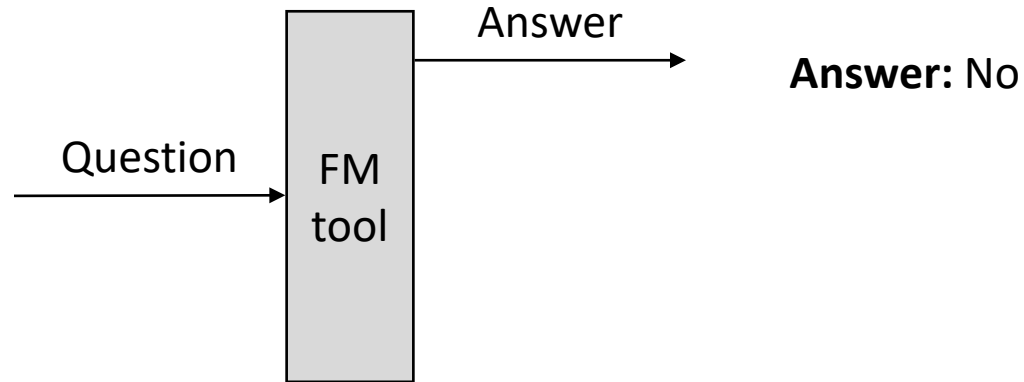
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The Present: A World Without Explainability

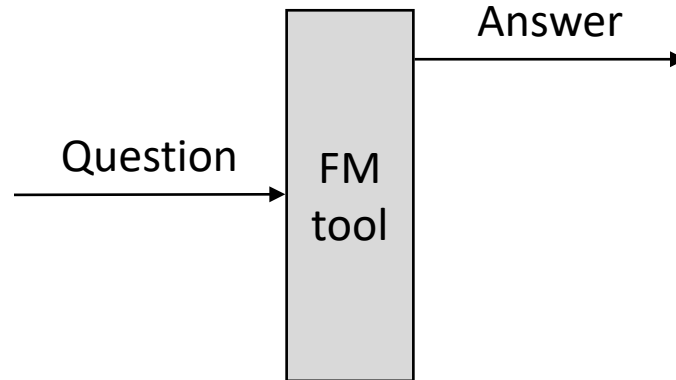
Question: Given taskset τ , the following run-time scheduler S , and assumptions A , is it unschedulable?





The Present: A World Without Explainability

Question: Given program P and restriction R on inputs, what is its worst-case execution time?

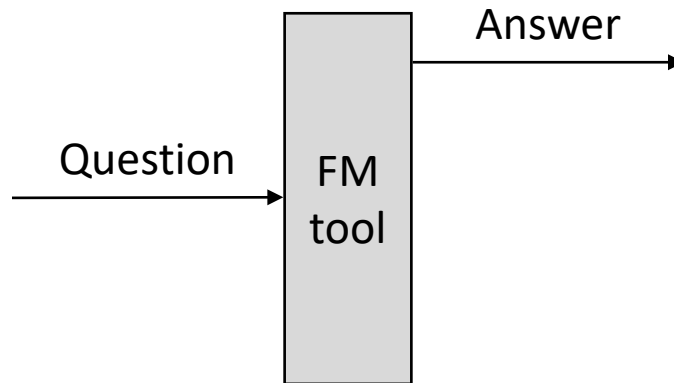


Answer: 42 milliseconds



The Present: A World Without Explainability

Question: Given program P and restriction R on inputs, does it hold for all executions that the execution time is less than or equal to TARGET_WCET_BOUND?

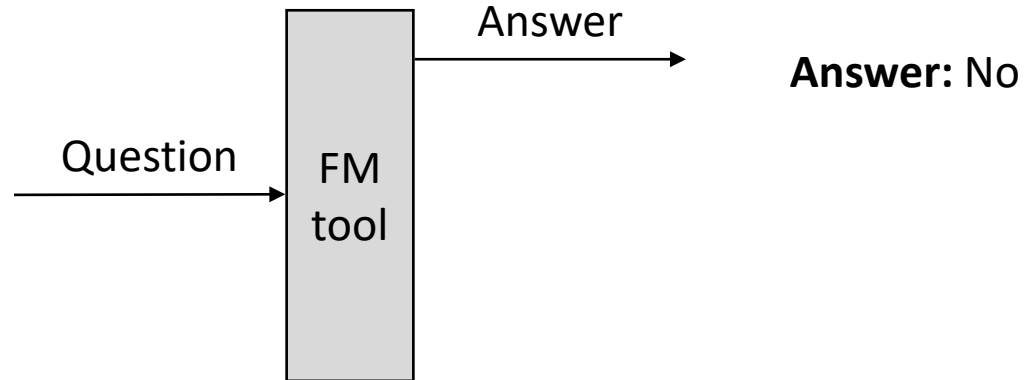


Answer: Yes



The Present: A World Without Explainability

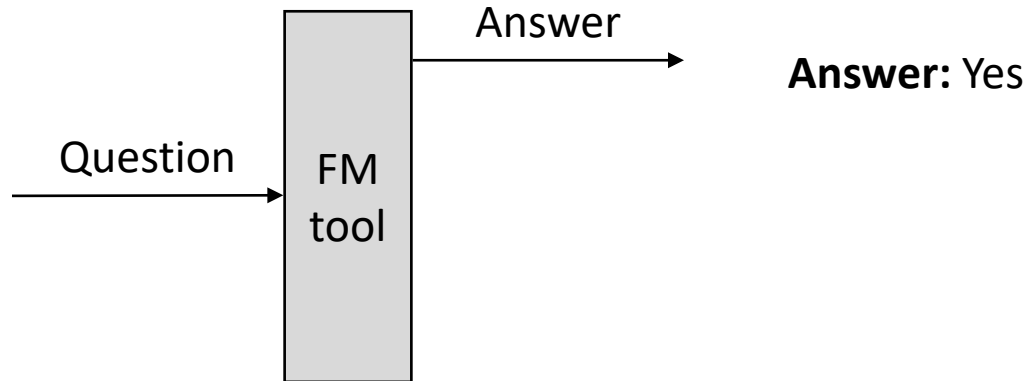
Question: Given program P ,
is assertion φ true for all
executions?





The Present: A World Without Explainability

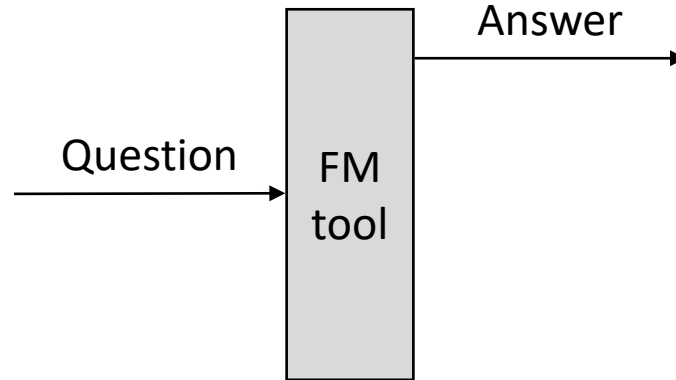
Question: Given program P ,
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The Present: A World Without Explainability

Question: Given the following executable code, is the assumption “no-task-suspension-can-occur-at-runtime” true?

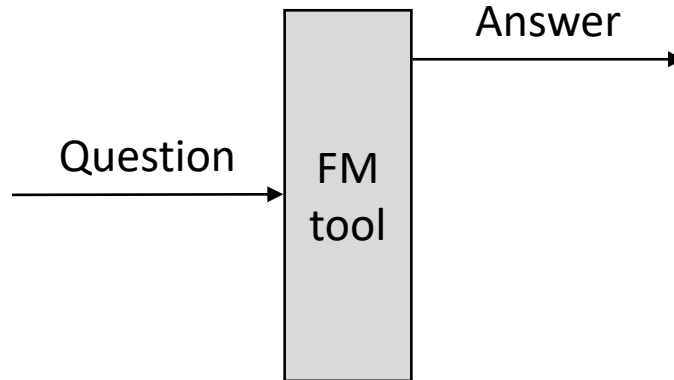


Answer: Yes



The Present: A World Without Explainability

Question: Given the following plant and controller " $x^{\text{new}} = A*x+B*u$, $u=K*x$ " does the schedule S of the task that executes the controller yield control performance J?

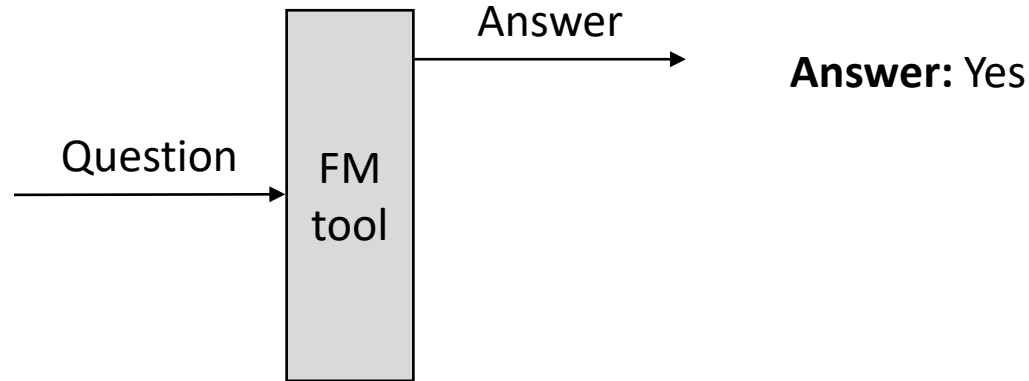


Answer: Yes



The Present: A World Without Explainability

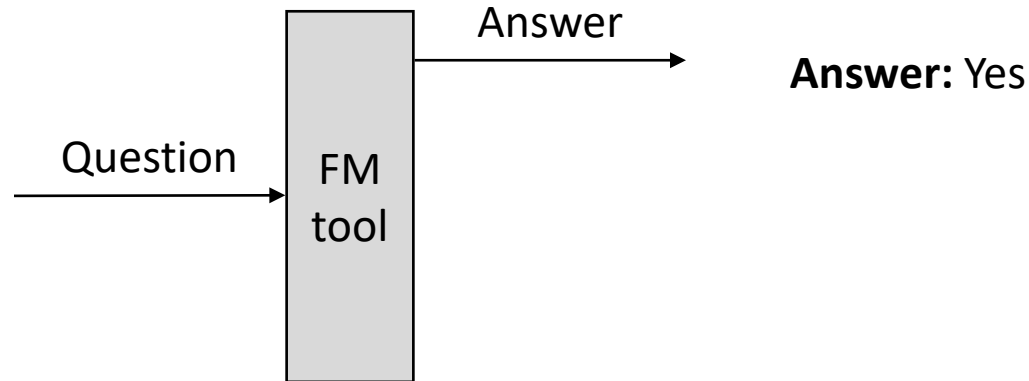
Question: Given the following plant and controller " $x^{\text{new}} = A*x+B*u$, $u=K*x$ " does the set of schedules S of the task that executes the controller yield control performance J or better?





The Present: A World Without Explainability

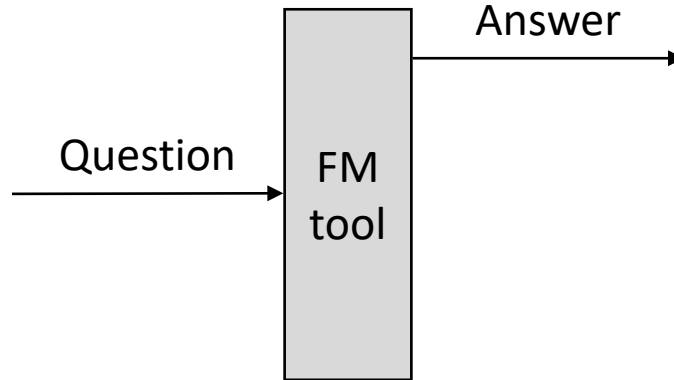
Question: Given the following human operator H, plant P, and control software C, does it hold that the mode perceived by H, P, and C differ by at most J?





The Present: A World Without Explainability

Question: Given the following distributed computer system D with message flow models S, and buffers B, does it hold that buffer overrun never occurs?

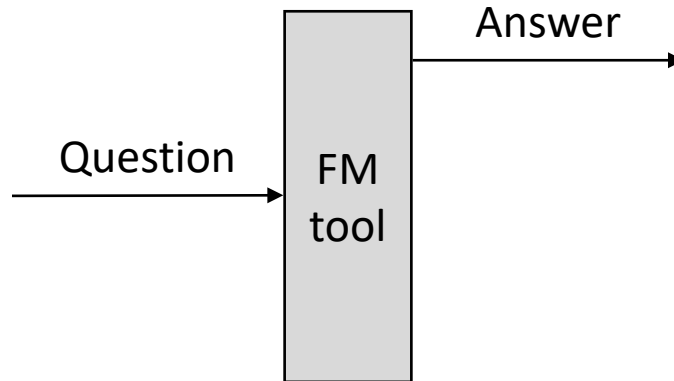


Answer: Yes



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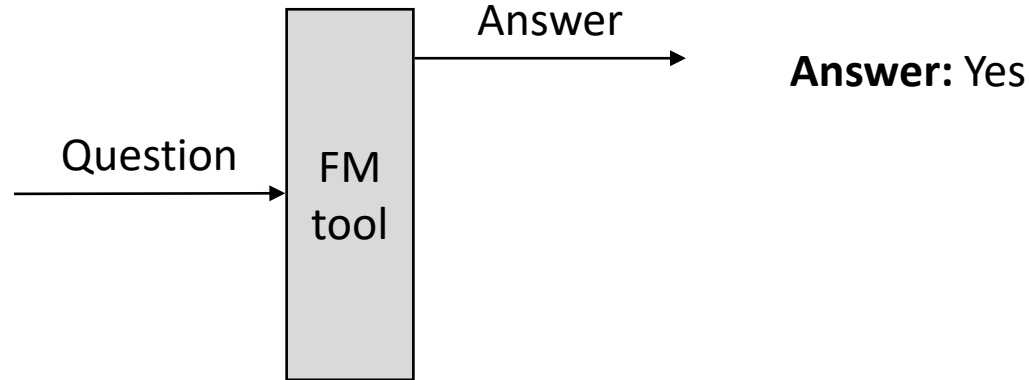
Answer: Yes

How much should one trust the FM tool that outputs answers?



The Present: A World Without Explainability

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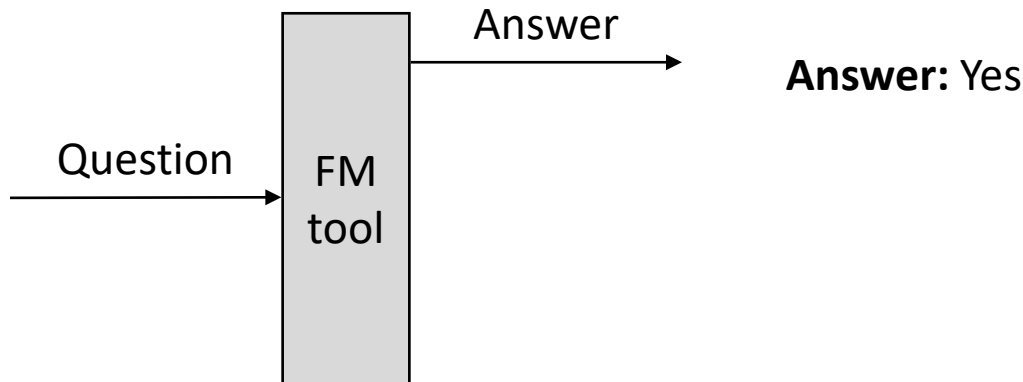


How much should one trust the FM tool that outputs answers?
If one does not know how the FM tool works, why should one trust it?



The Present: A World Without Explainability

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How much should one trust the FM tool that outputs answers?

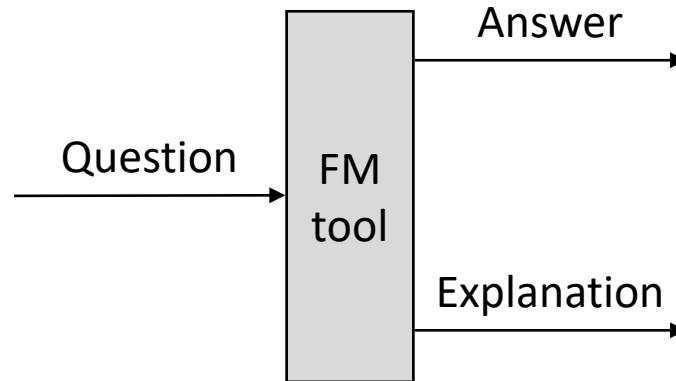
If one does not know how the FM tool works, why should one trust it?

If practitioners do not trust the FM tools that researchers produce, how much value is there in the research?



A Potential Future: A World With Explainability

Question: Given taskset τ , the following run-time scheduler S , and assumptions A , is it schedulable?

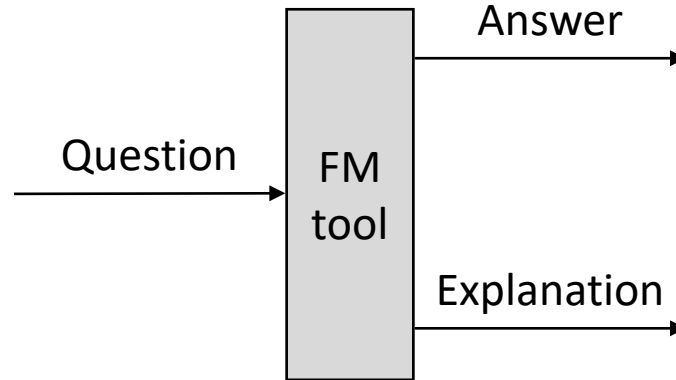


Answer: Yes



A Potential Future: A World With Explainability

Question: Given program P and restriction R on inputs, what is its worst-case execution time?

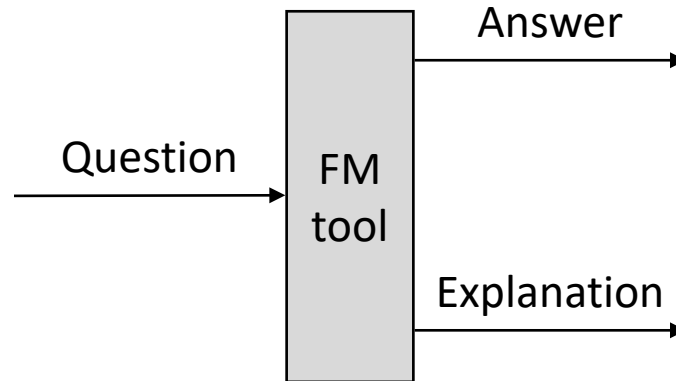


Answer: 42 milliseconds



A Potential Future: A World With Explainability

Question: Given program P and restriction R on inputs, does it hold for all executions that the execution time is less than or equal to TARGET_WCET_BOUND?

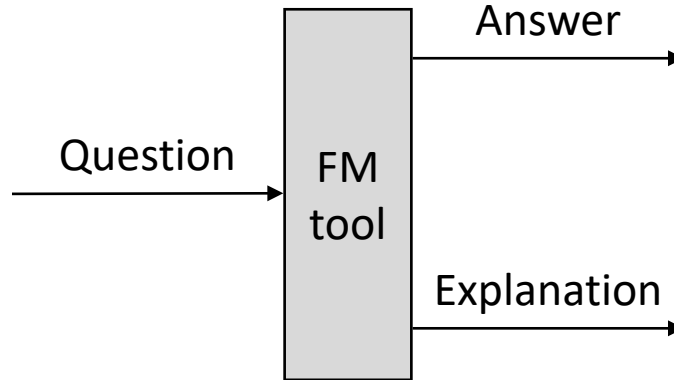


Answer: Yes



A Potential Future: A World With Explainability

Question: Given program P ,
is assertion φ true for all
executions?

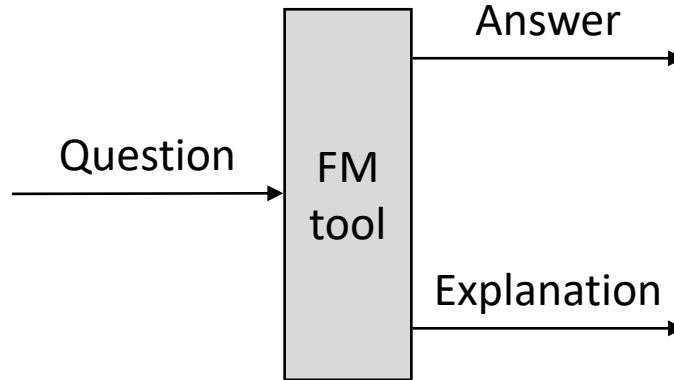


Answer: Yes



A Potential Future: A World With Explainability

Question: Given program P ,
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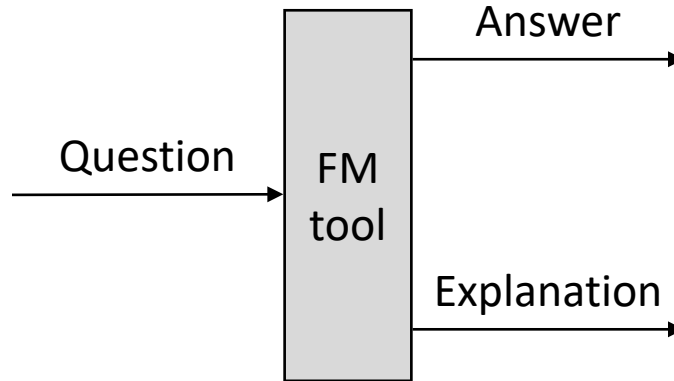


Answer: No



A Potential Future: A World With Explainability

Question: Given the following taskset, is the assumption “no-task-suspension-can-occur-at-runtime” true?

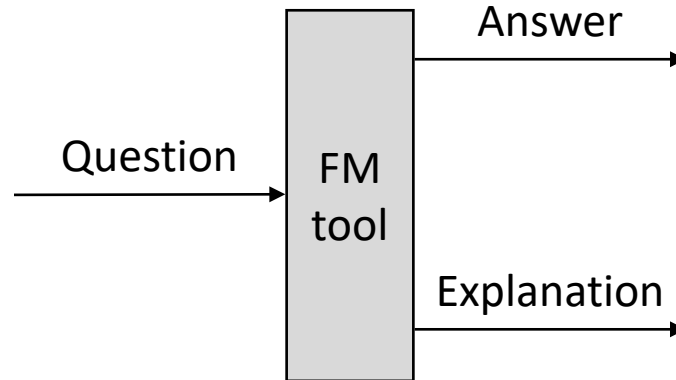


Answer: Yes



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Question: Given the following plant and controller " $x^{\text{new}} = A*x+B*u$, $u=K*x$ " does the schedule S of the task that executes the controller yield control performance J?

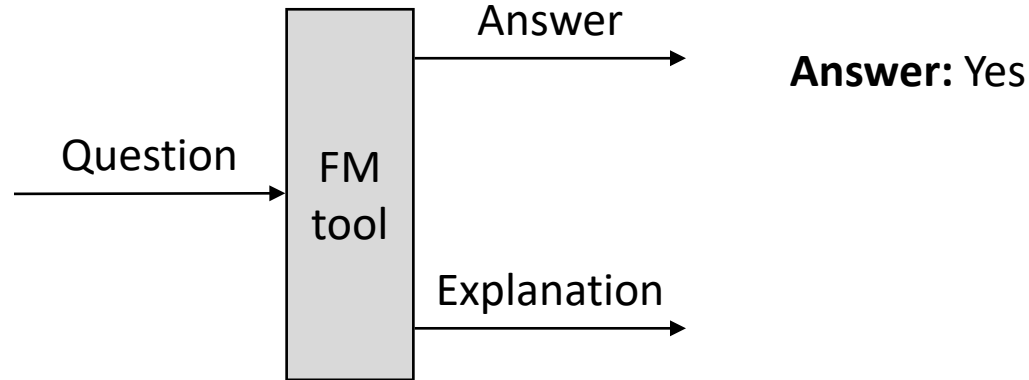


Answer: Yes



A Potential Future: A World With Explainability

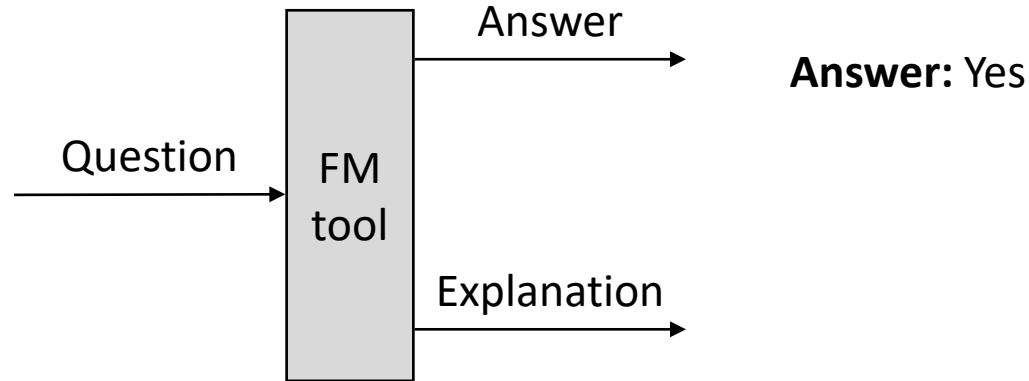
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A Potential Future: A World With Explainability

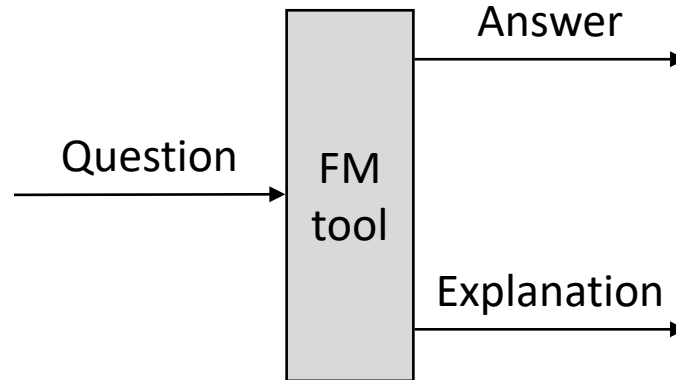
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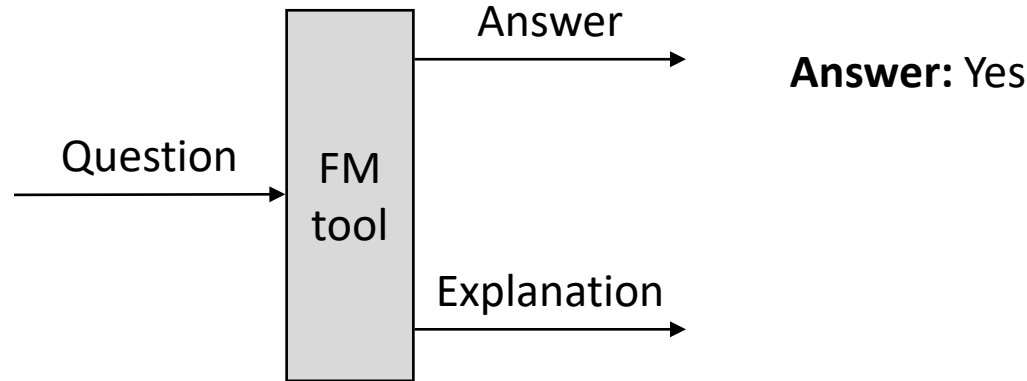


Answer: Yes



A Potential Future: A World With Explainability

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Explainability has the potential to make practitioners trust FM tools

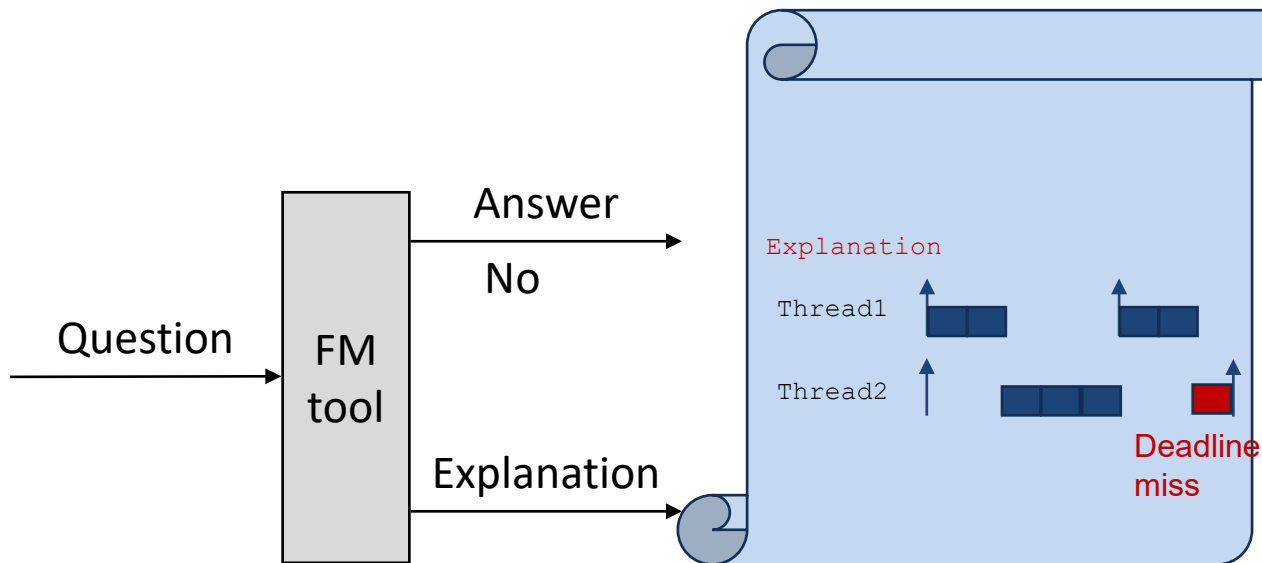
Explainability has the potential to make sure our research becomes more valuable to practitioners



What Could an Explanation Look Like?

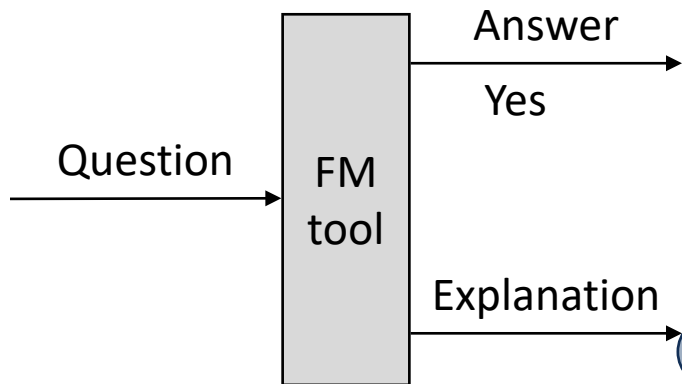


Question: Given taskset τ , the following run-time scheduler S , and assumptions A , is it schedulable?



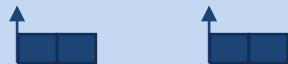


Question: Given taskset τ' , the following run-time scheduler S , and assumptions A , is it schedulable?



Explanation

Thread1



Thread2

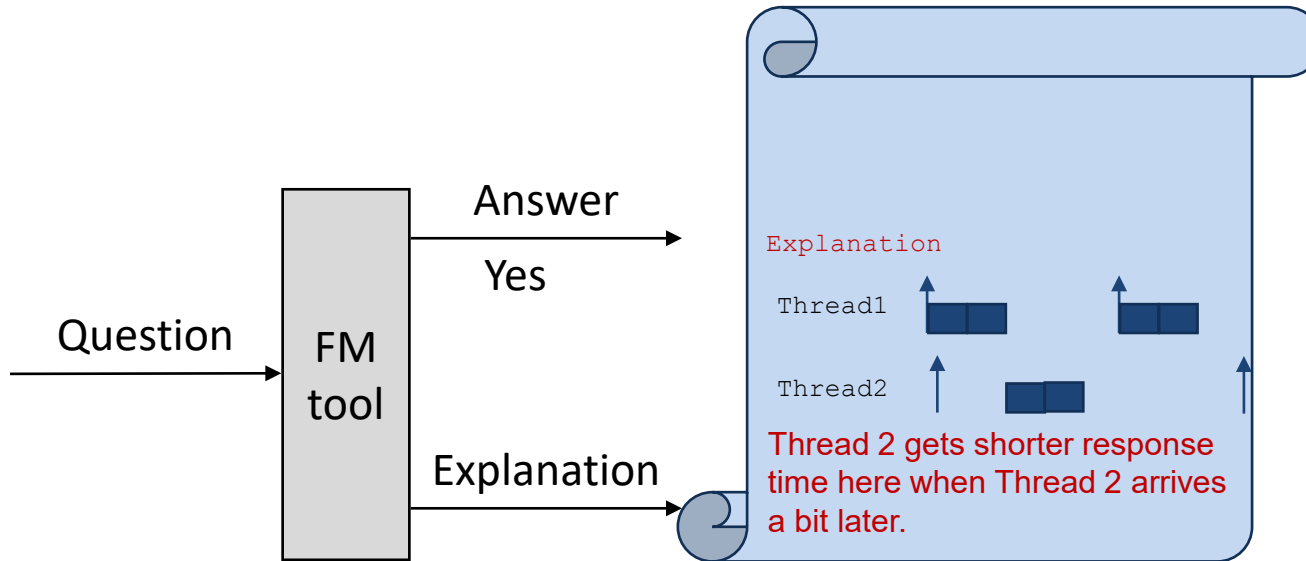


Thread 2 meets its deadline for this situation. This is the worst-case situation for Thread 2.

Click here to understand why this is the worst-case situation

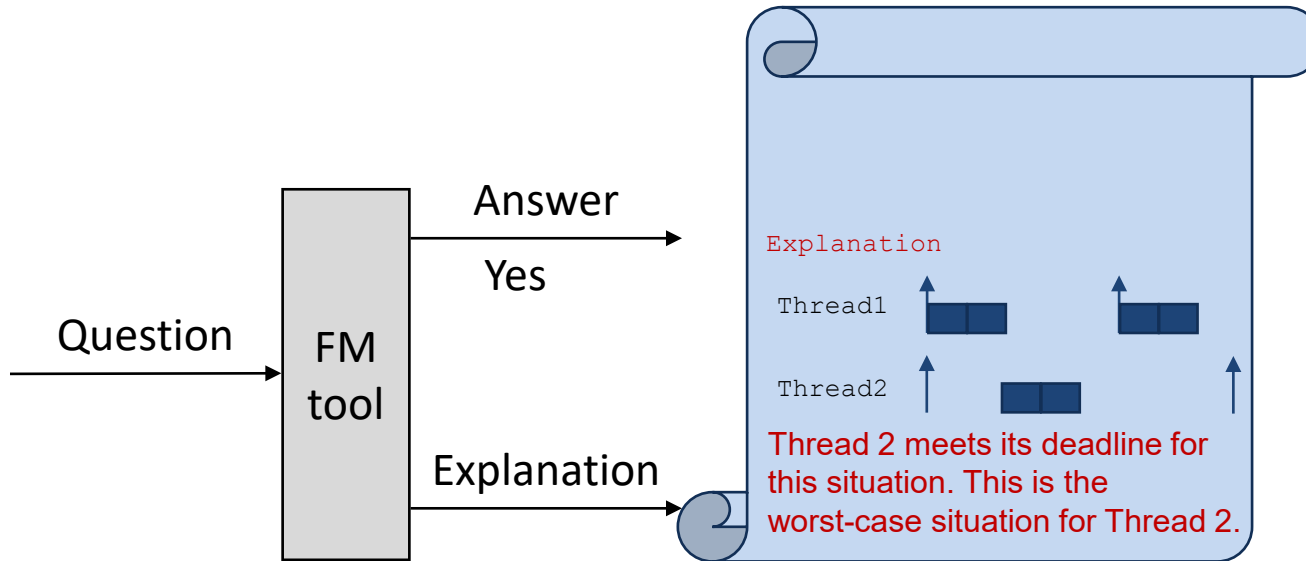


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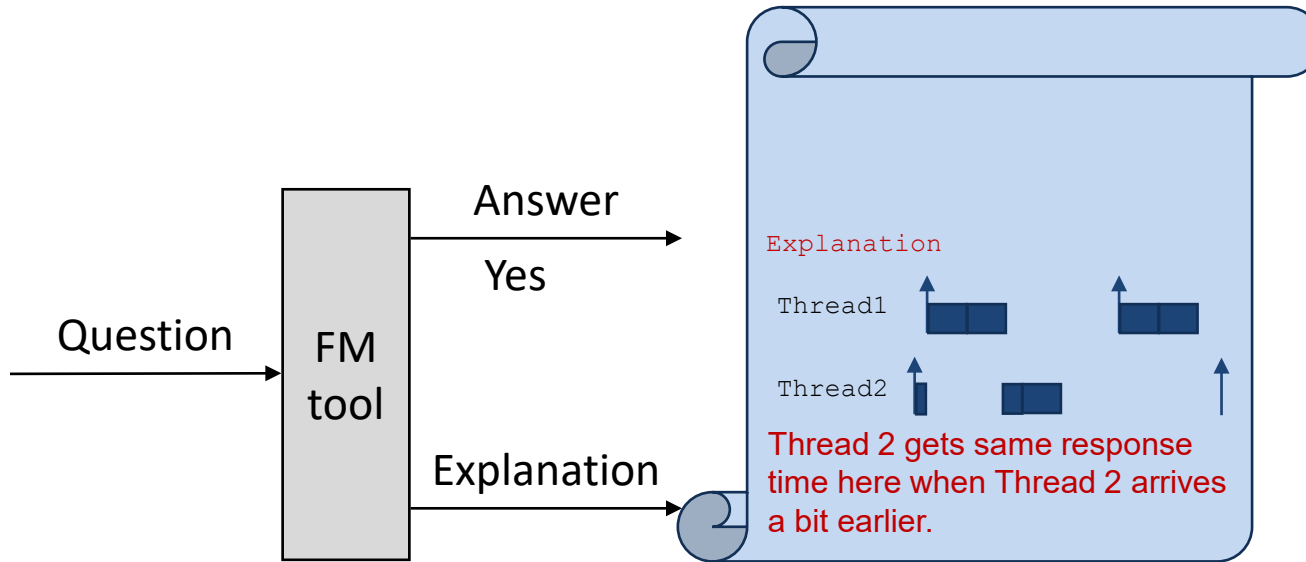


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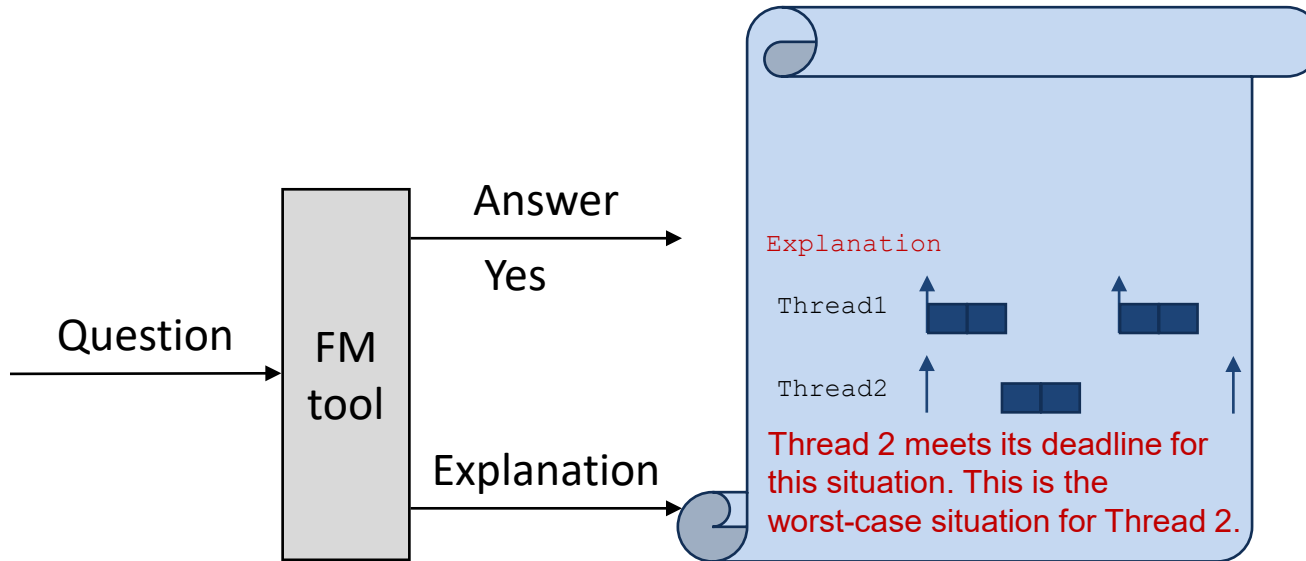


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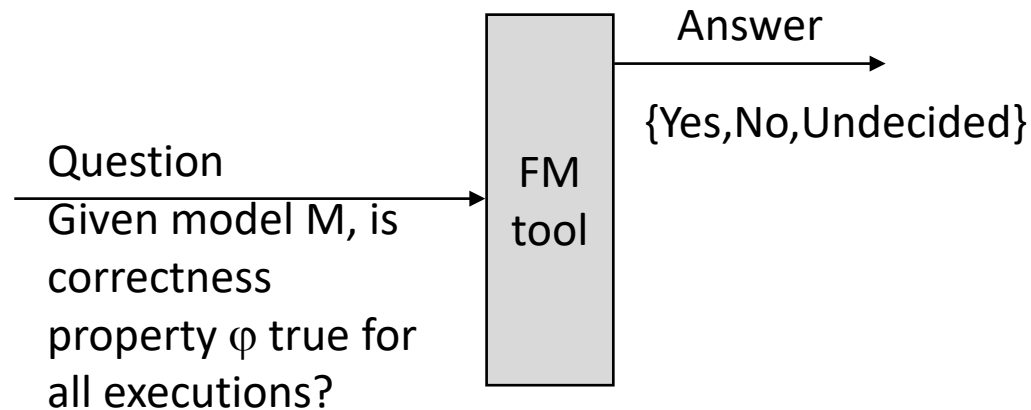
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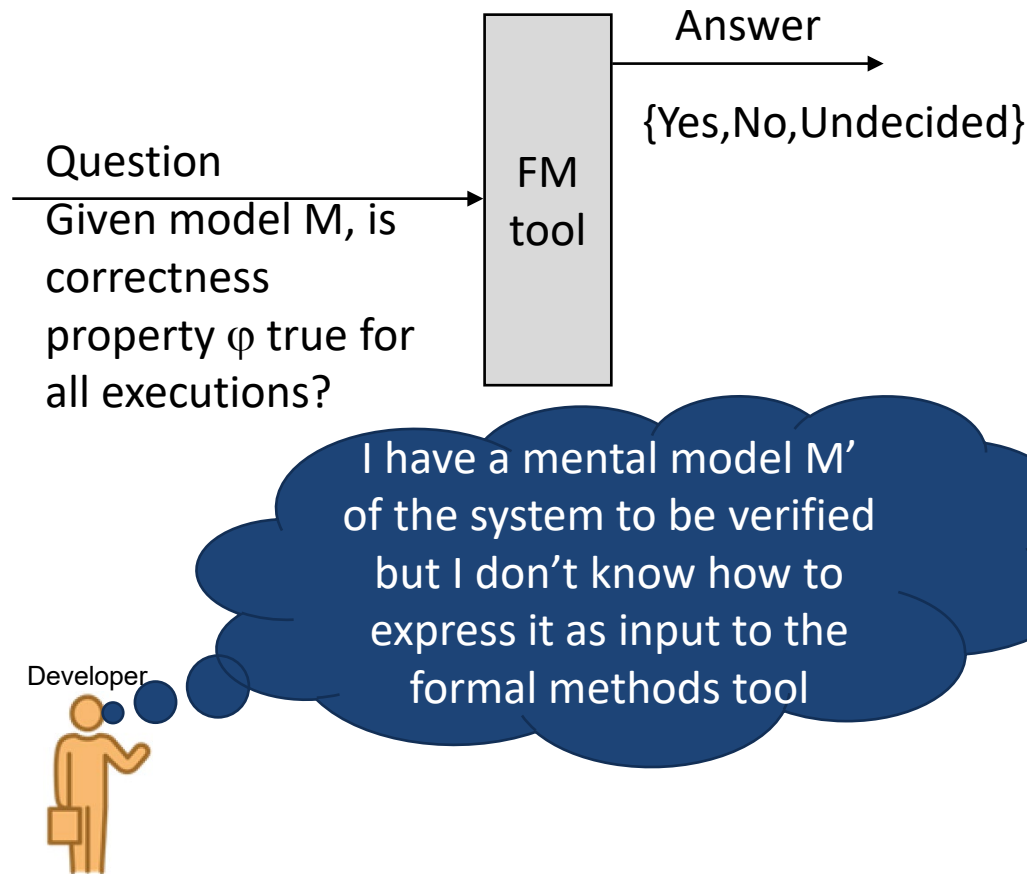


Explanation of Input

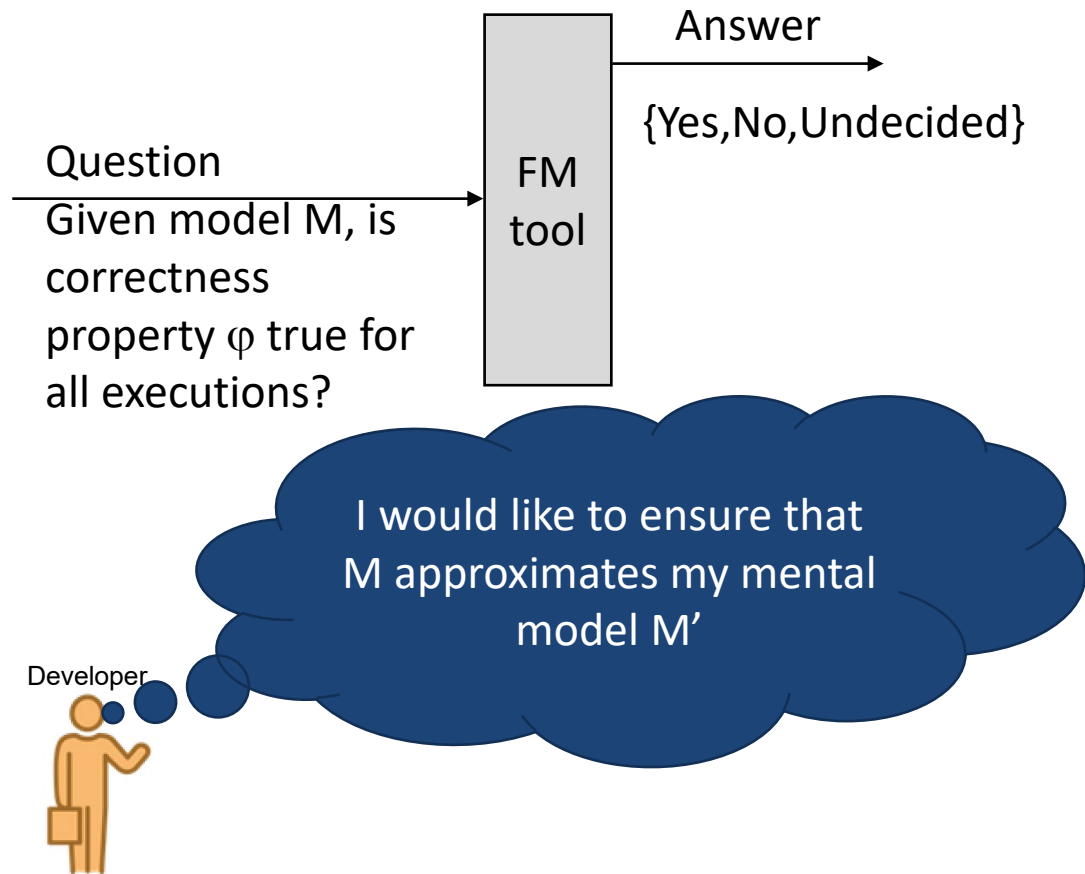
Explainability of Input



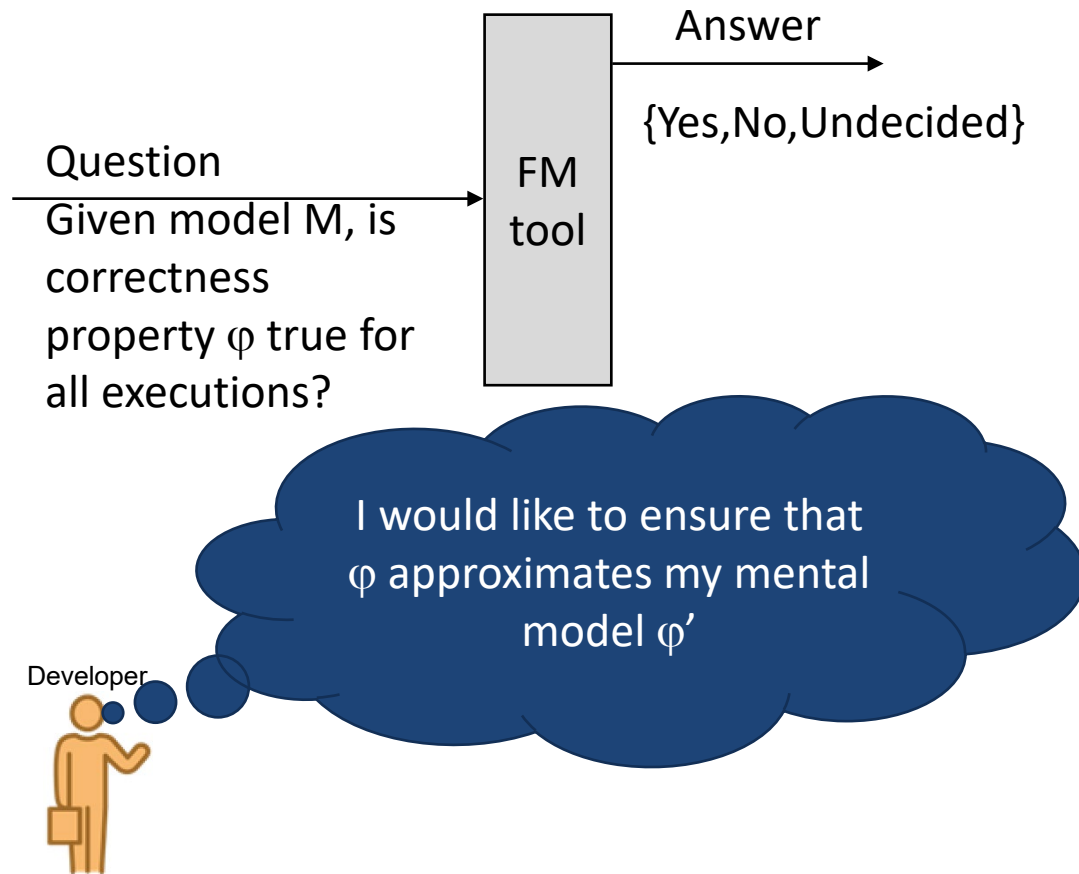
Explainability of Input



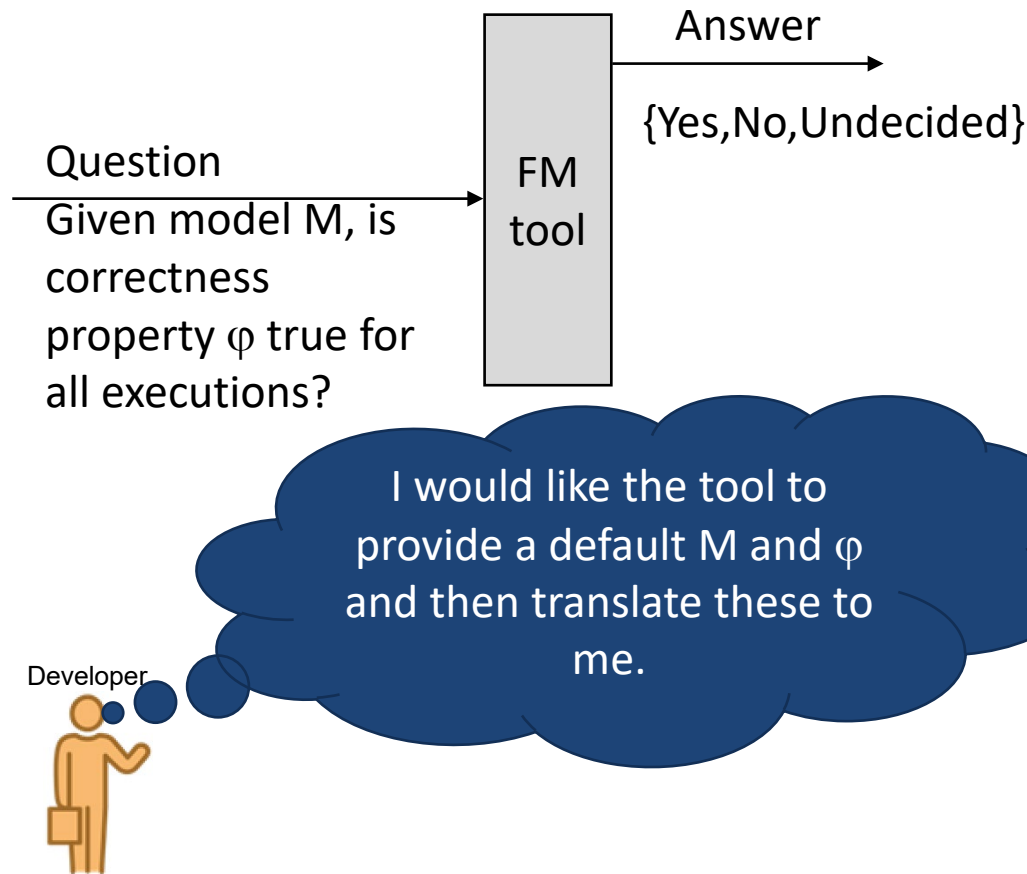
Explainability of Input



Explainability of Input

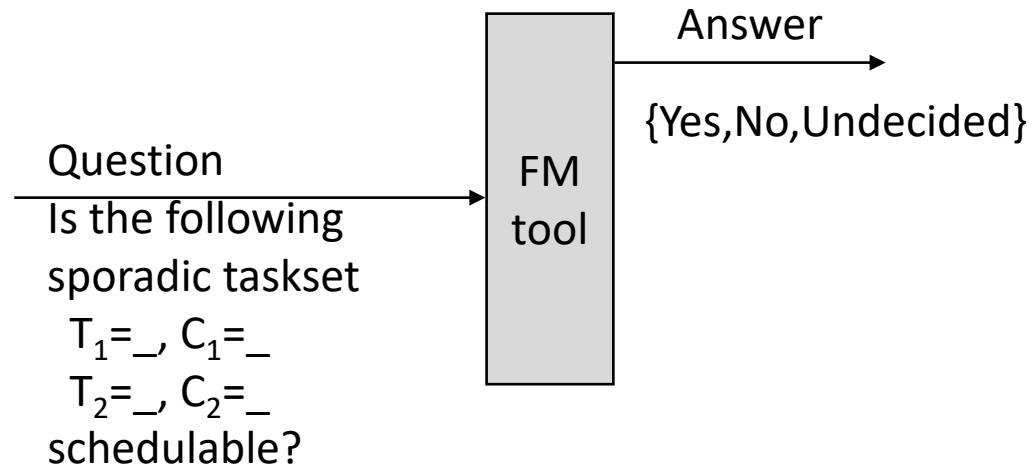


Explainability of Input





Explainability of Input: Real-Time Scheduling

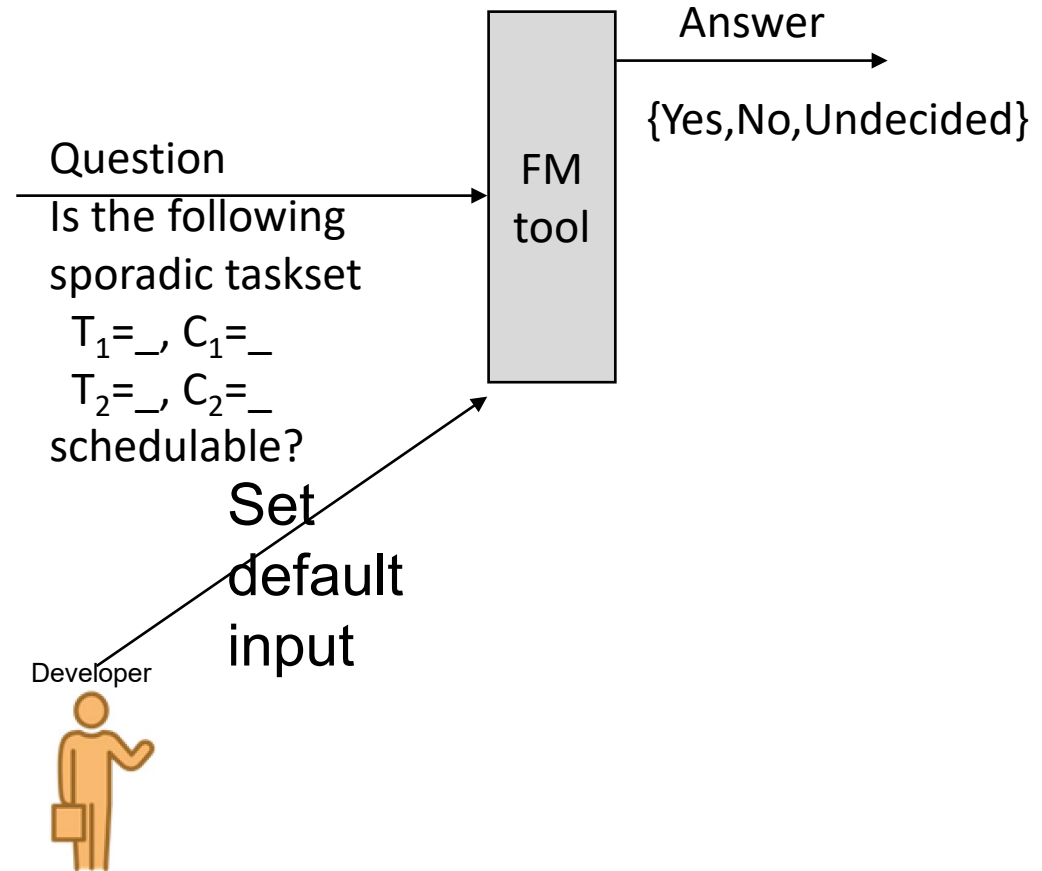


Developer



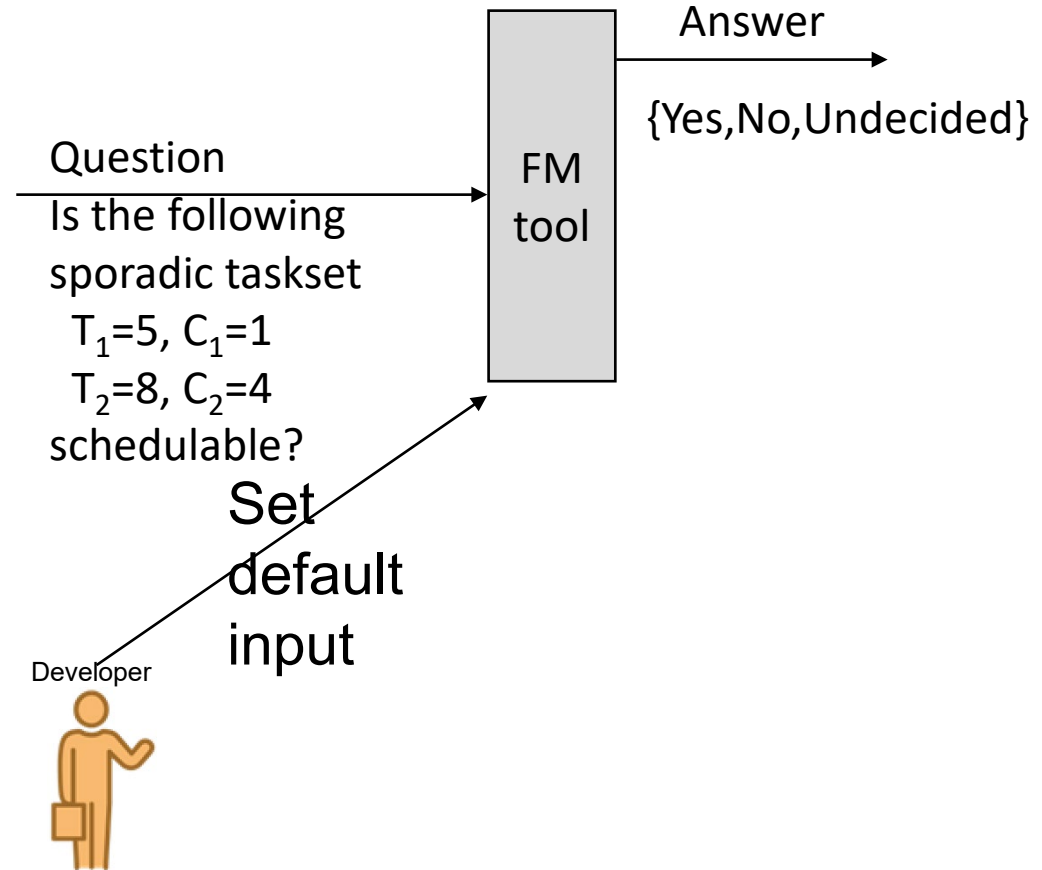


Explainability of Input: Real-Time Scheduling



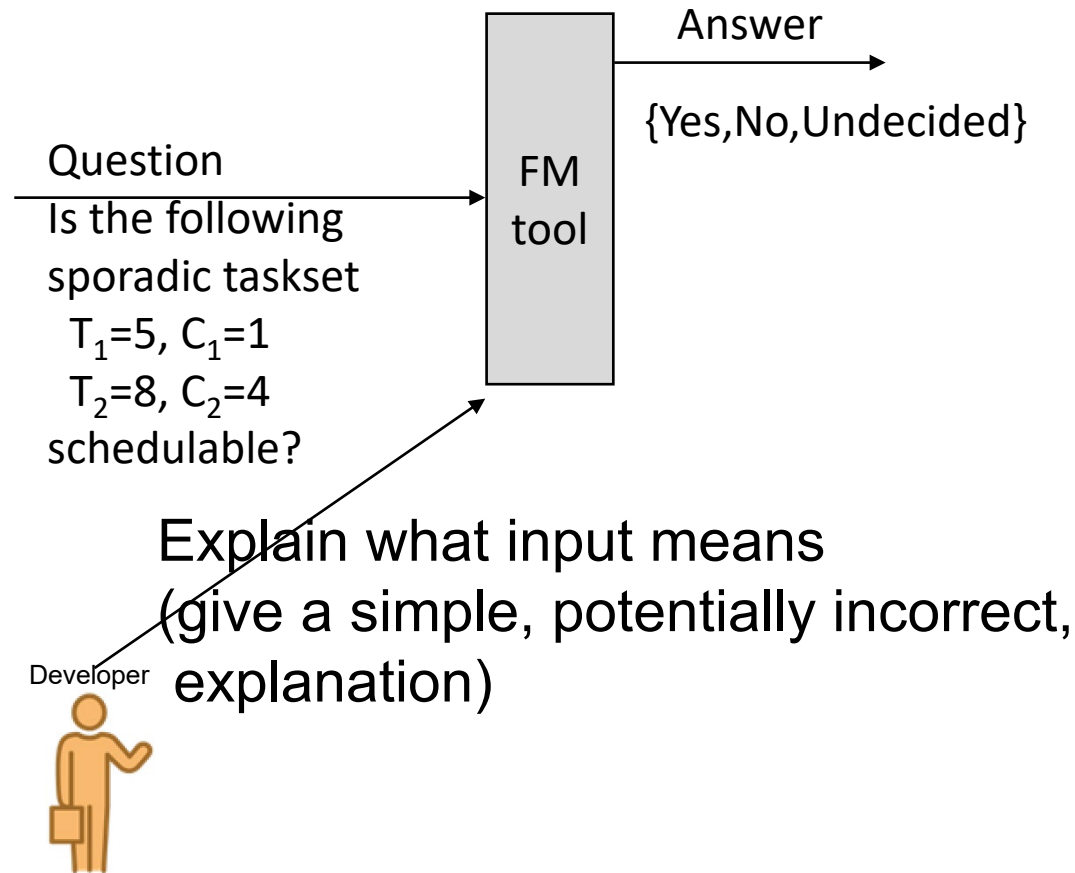


Explainability of Input: Real-Time Scheduling





Explainability of Input: Real-Time Scheduling





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5
while (1) {
    delay_until(t)
    do_work(1)
    t=t+5
}
```

```
Thread 2
t=now()+8
while (1) {
    delay_until(t)
    do_work(4)
    t=t+8
}
```

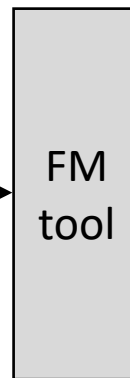
Question

Is the following
sporadic taskset

$T_1=5, C_1=1$

$T_2=8, C_2=4$

schedulable?



Answer

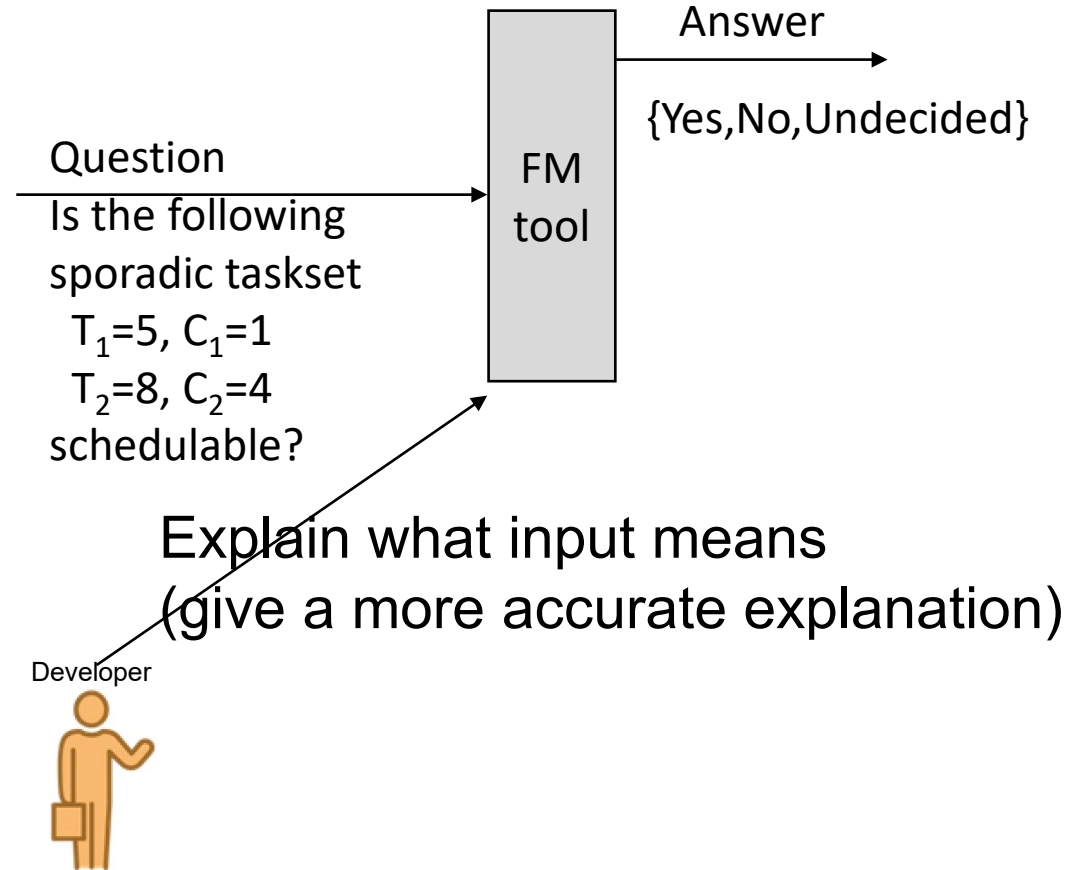
{Yes,No,Undecided}

Developer





Explainability of Input: Real-Time Scheduling





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5
}
```

```
Thread 2
t=now()+8
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8
}
```

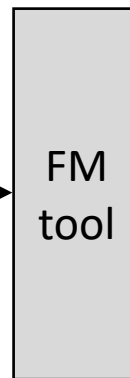
Question

Is the following
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$T_2=8, C_2=4$

schedulable?



Answer

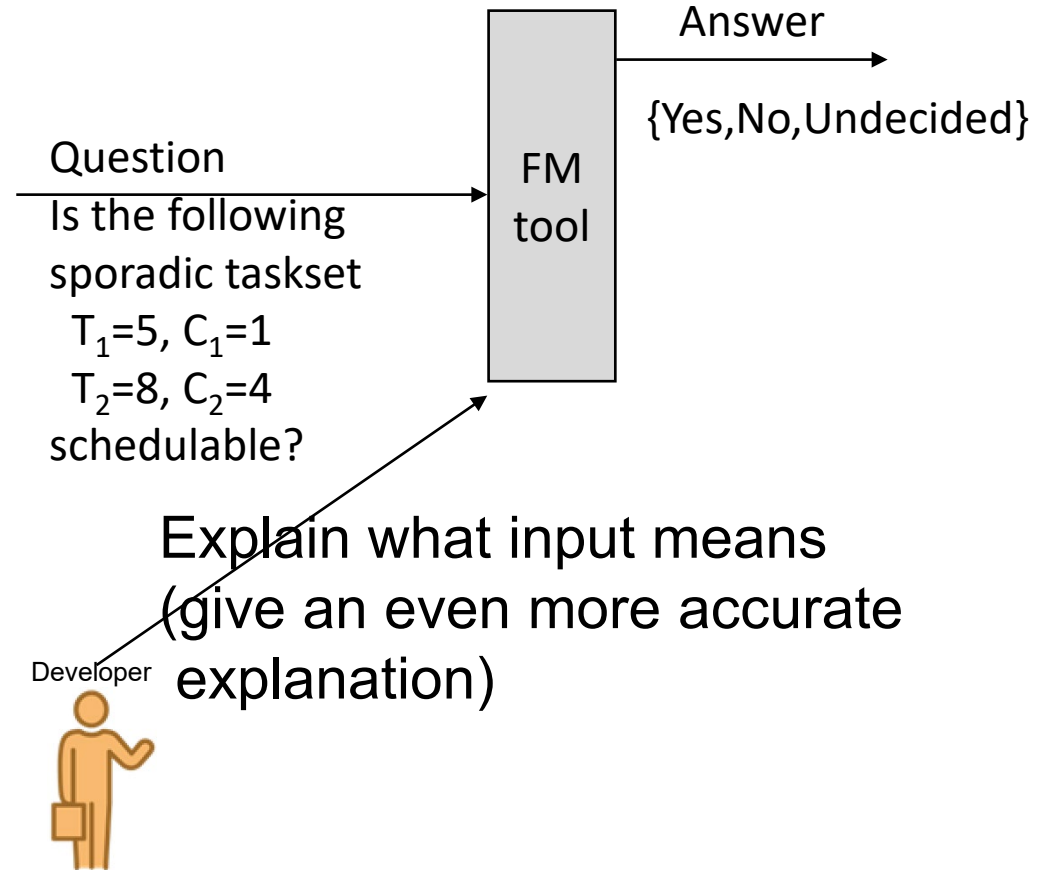
{Yes,No,Undecided}

Developer





Explainability of Input: Real-Time Scheduling





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```

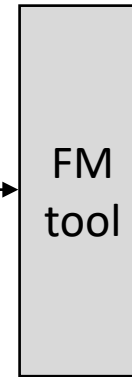
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Answer

{Yes,No,Undecided}

Developer

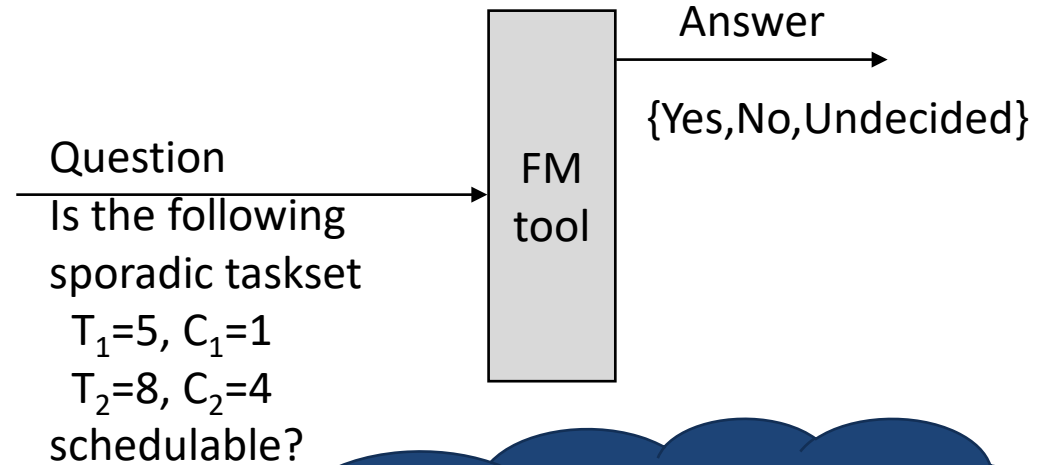




Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```



Developer



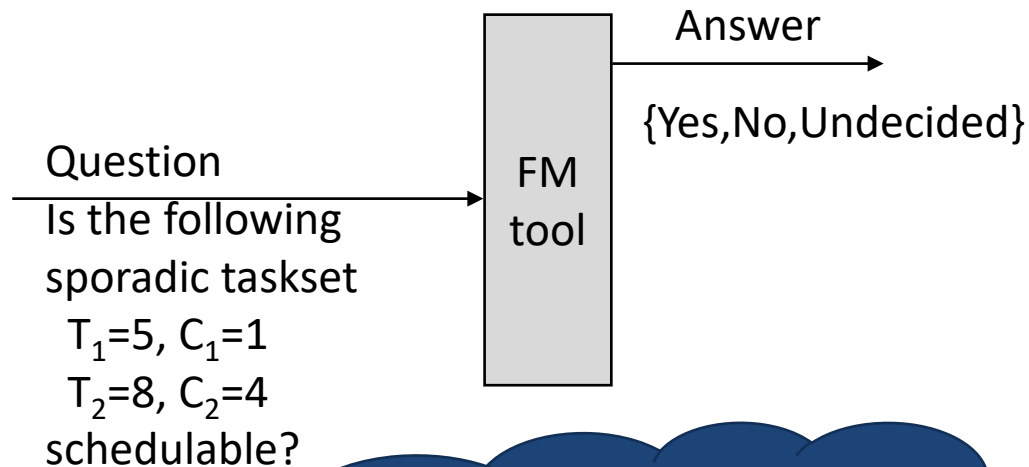
Now, I understand what these parameters mean.



Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```



Developer



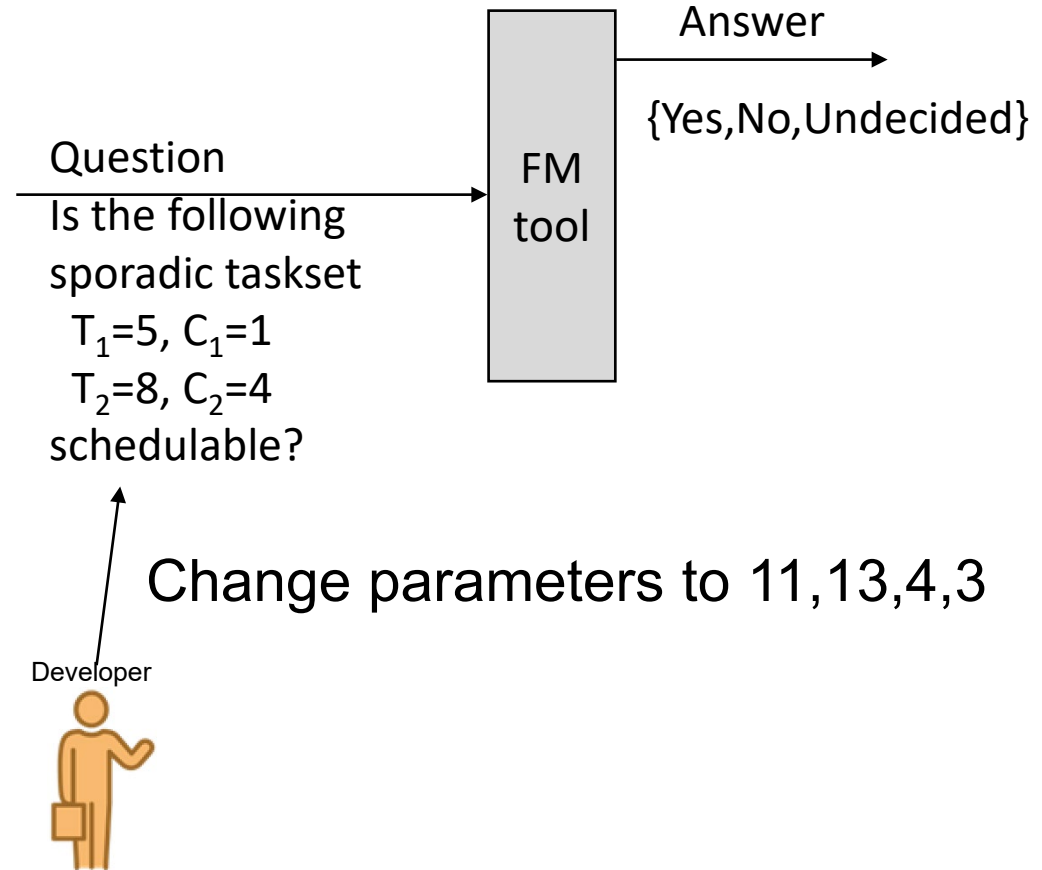
I want to change these parameters so that they reflect the system I want to analyze



Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```

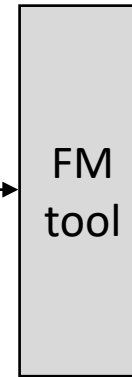
Question

Is the following
sporadic taskset

$T_1=11, C_1=4$

$T_2=13, C_2=3$

schedulable?



Answer

{Yes,No,Undecided}

Developer

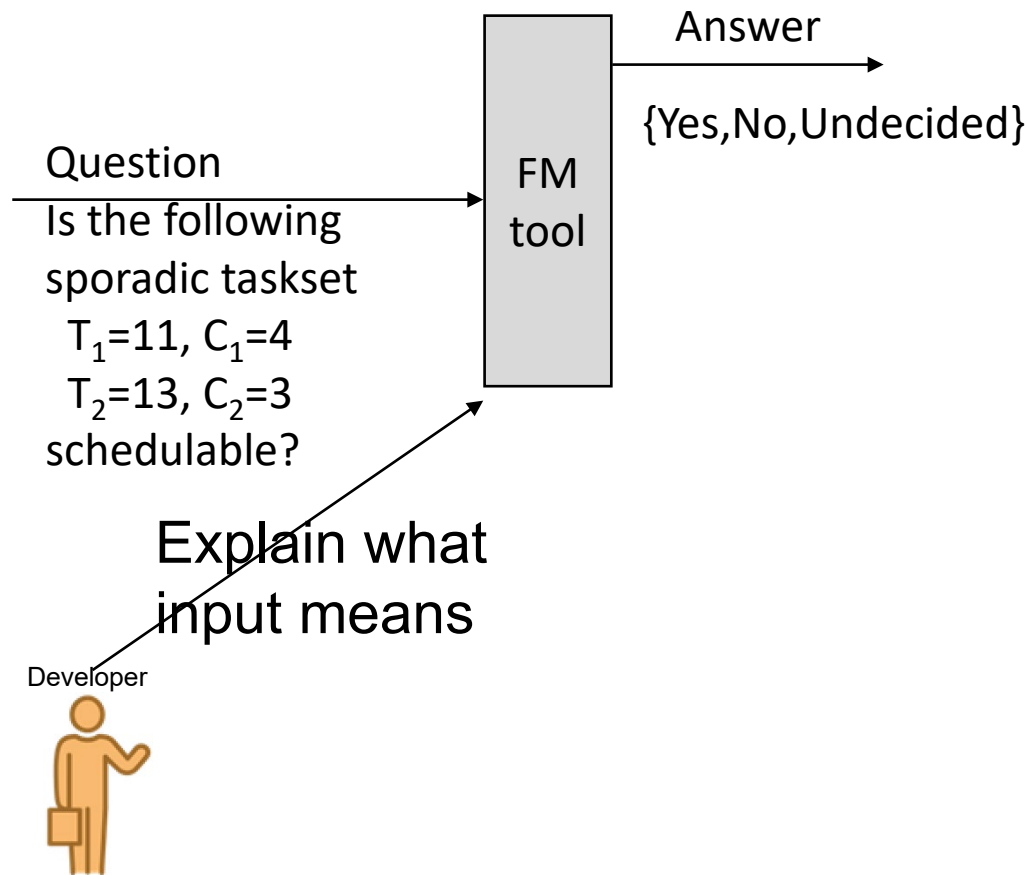




Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+5+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
}
```

```
Thread 2
t=now()+8+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
}
```





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+11+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
}
```

```
Thread 2
t=now()+13+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
}
```

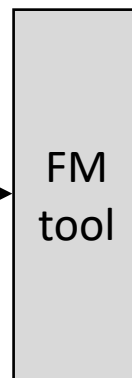
Question

Is the following
sporadic taskset

$T_1=11, C_1=4$

$T_2=13, C_2=3$

schedulable?



Answer

{Yes,No,Undecided}

Developer

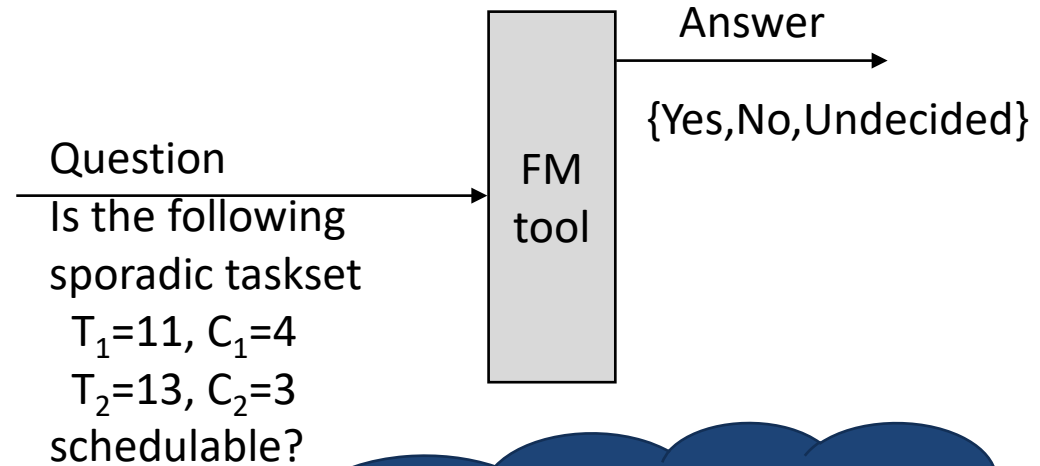




Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+11+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
}

Thread 2
t=now()+13+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
}
```



Developer



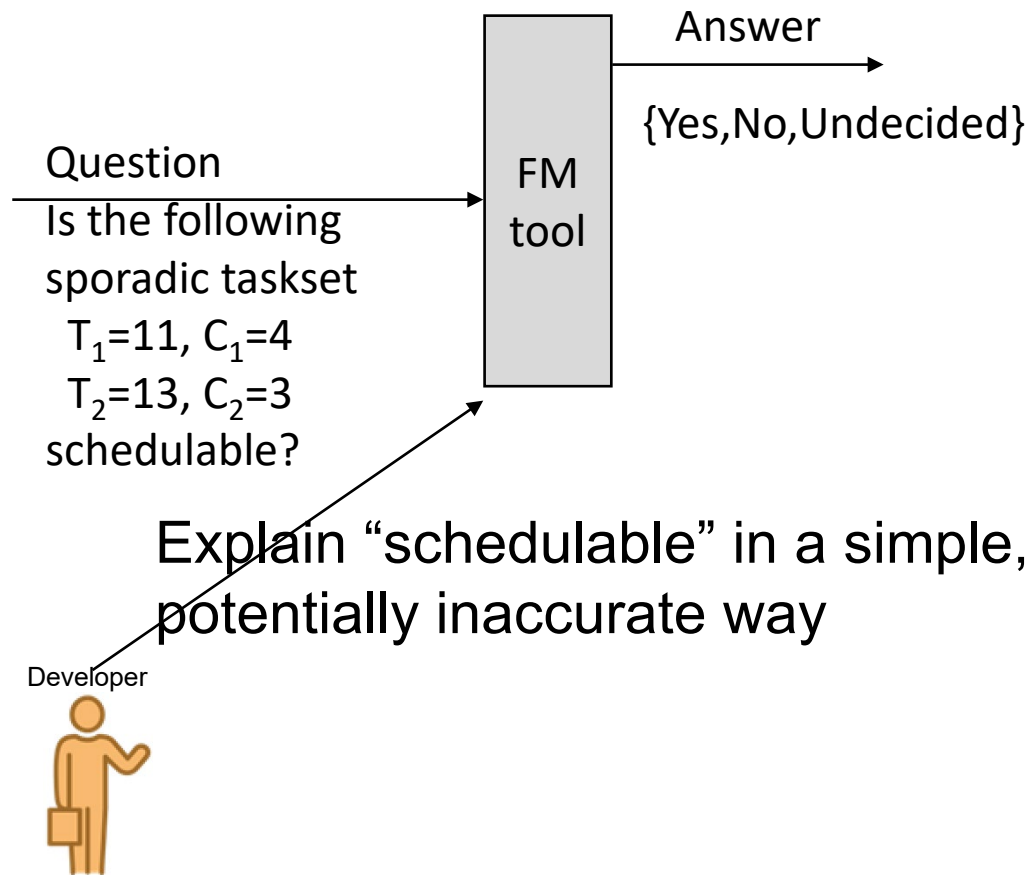
I don't know the meaning of the word "schedulable"



Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+11+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
}
```

```
Thread 2
t=now()+13+rand(0,inf)
while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
}
```





Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+11+rand(0,inf)
while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
}
```

```
Thread 2
t=now()+13+rand(0,inf)
while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
}
```

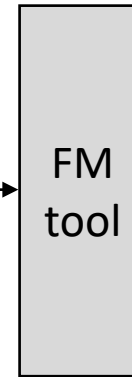
Question

Is the following
sporadic taskset

$T_1=11, C_1=4$

$T_2=13, C_2=3$

schedulable?



Answer

{Yes,No,Undecided}

Developer

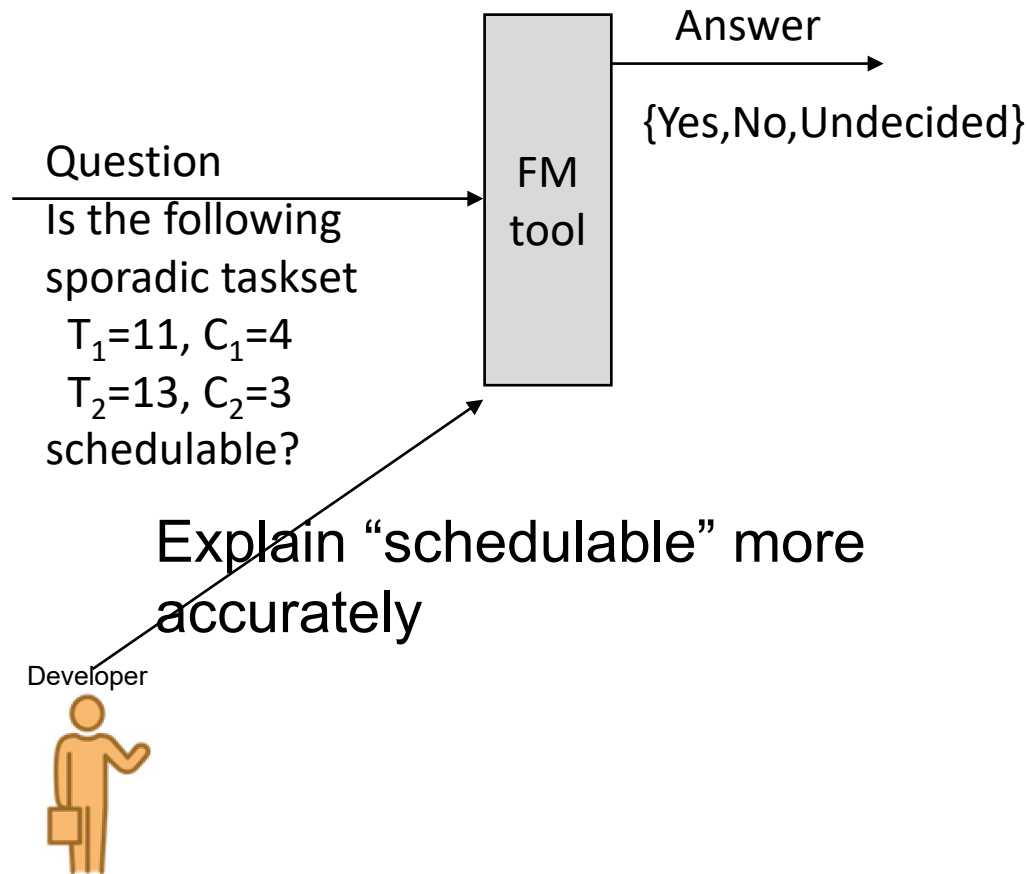




Explainability of Input: Real-Time Scheduling

```
Thread 1
t=now()+11+rand(0,inf)
while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
}
```

```
Thread 2
t=now()+13+rand(0,inf)
while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
}
```





Explainability of Input: Real-Time Scheduling

```
Thread 1
gap=rand(0,inf)
t=now()+11+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
}
```

```
Thread 2
gap=rand(0,inf)
t=now()+13+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
}
```

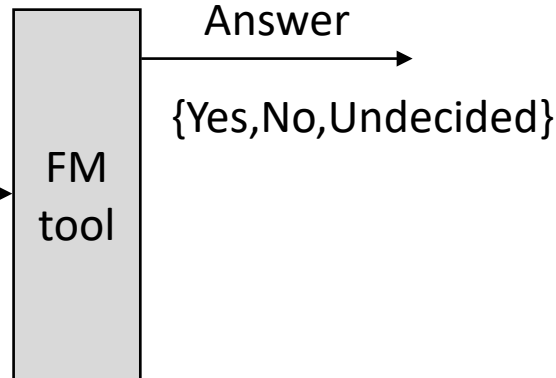
Question

Is the following
sporadic taskset

$T_1=11, C_1=4$

$T_2=13, C_2=3$

schedulable?



Developer

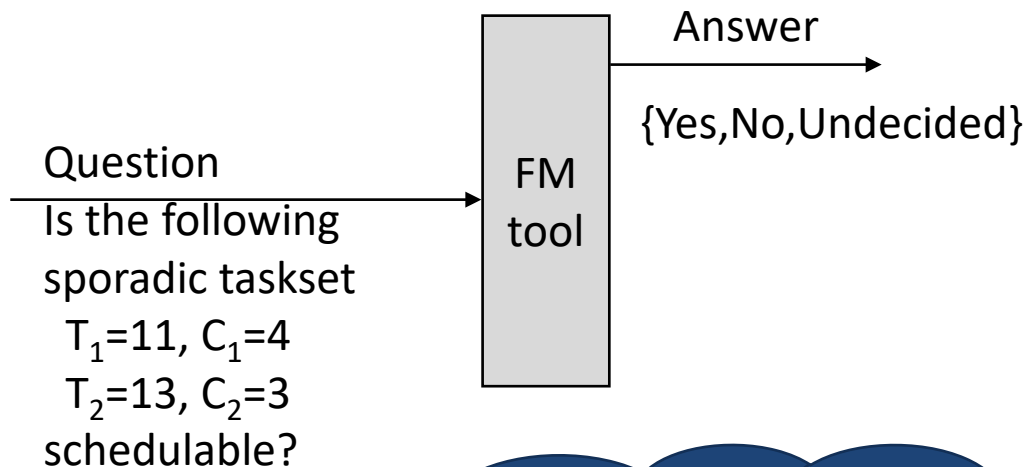




Explainability of Input: Real-Time Scheduling

```
Thread 1
gap=rand(0,inf)
t=now()+11+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
}
```

```
Thread 2
gap=rand(0,inf)
t=now()+13+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
}
```



Developer

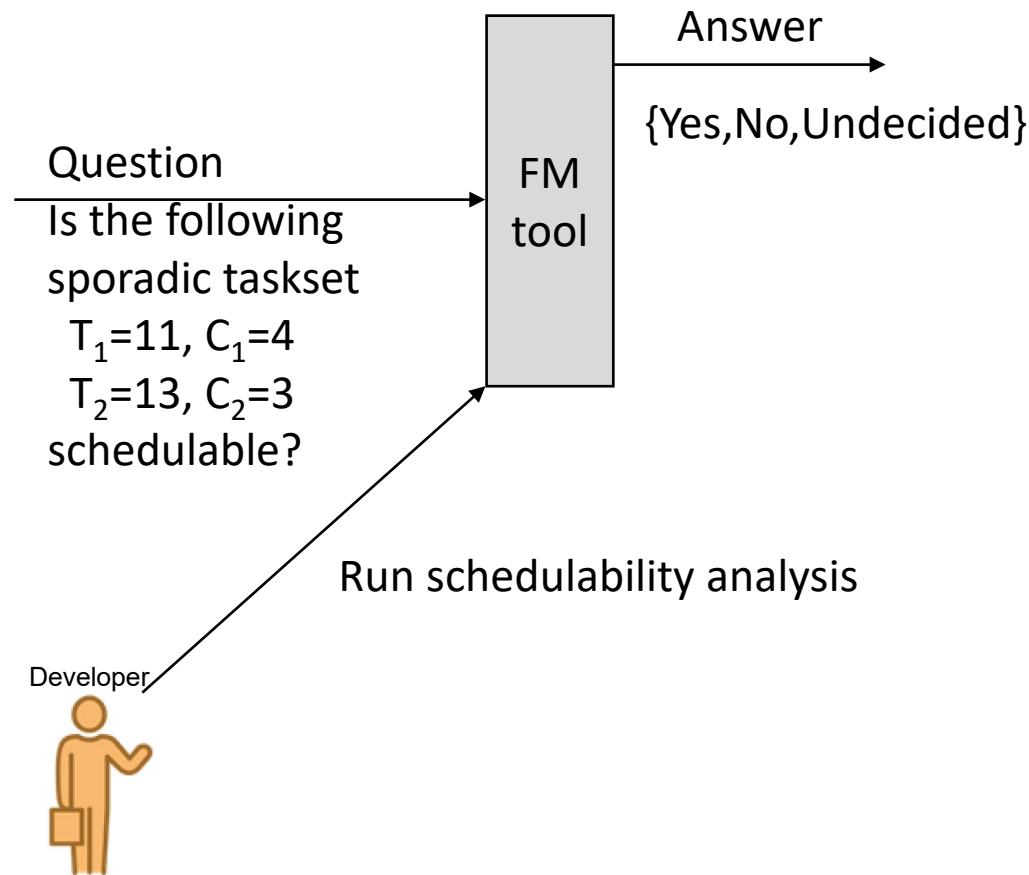
Now, I know the meaning of the taskset parameters and the word "schedulable." And the taskset parameters reflect the system I want to analyze



Explainability of Input: Real-Time Scheduling

```
Thread 1
gap=rand(0,inf)
t=now()+11+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
}
```

```
Thread 2
gap=rand(0,inf)
t=now()+13+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
}
```





Explainability of Input: Real-Time Scheduling

```
Thread 1
gap=rand(0,inf)
t=now()+11+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
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Thread 2
gap=rand(0,inf)
t=now()+13+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
}
```

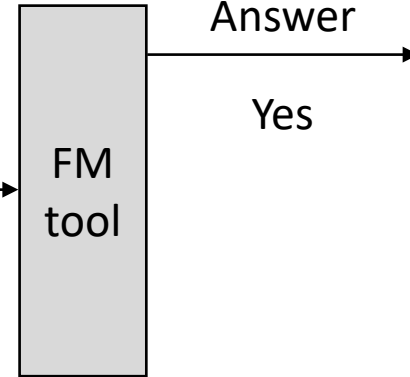
Question

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sporadic taskset

$T_1=11, C_1=4$

$T_2=13, C_2=3$

schedulable?



Developer

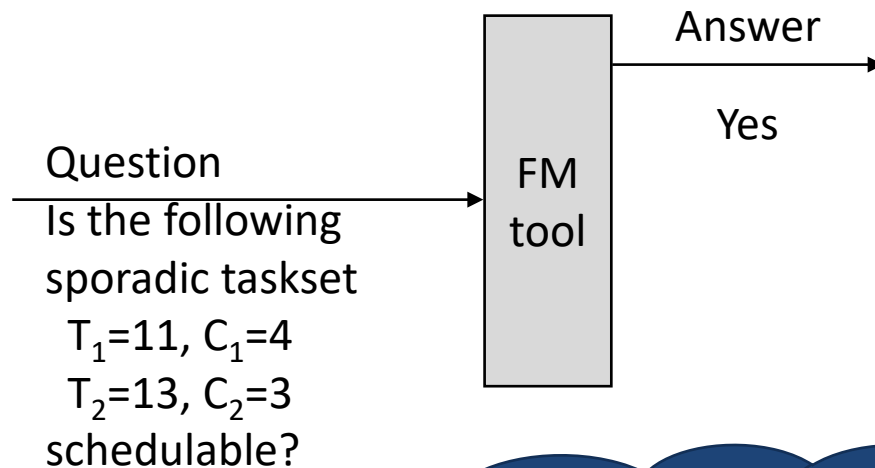




Explainability of Input: Real-Time Scheduling

```
Thread 1
gap=rand(0,inf)
t=now()+11+gap
while (1) {
    assert(now()<=t-gap)
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    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
}
```

```
Thread 2
gap=rand(0,inf)
t=now()+13+gap
while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
}
```

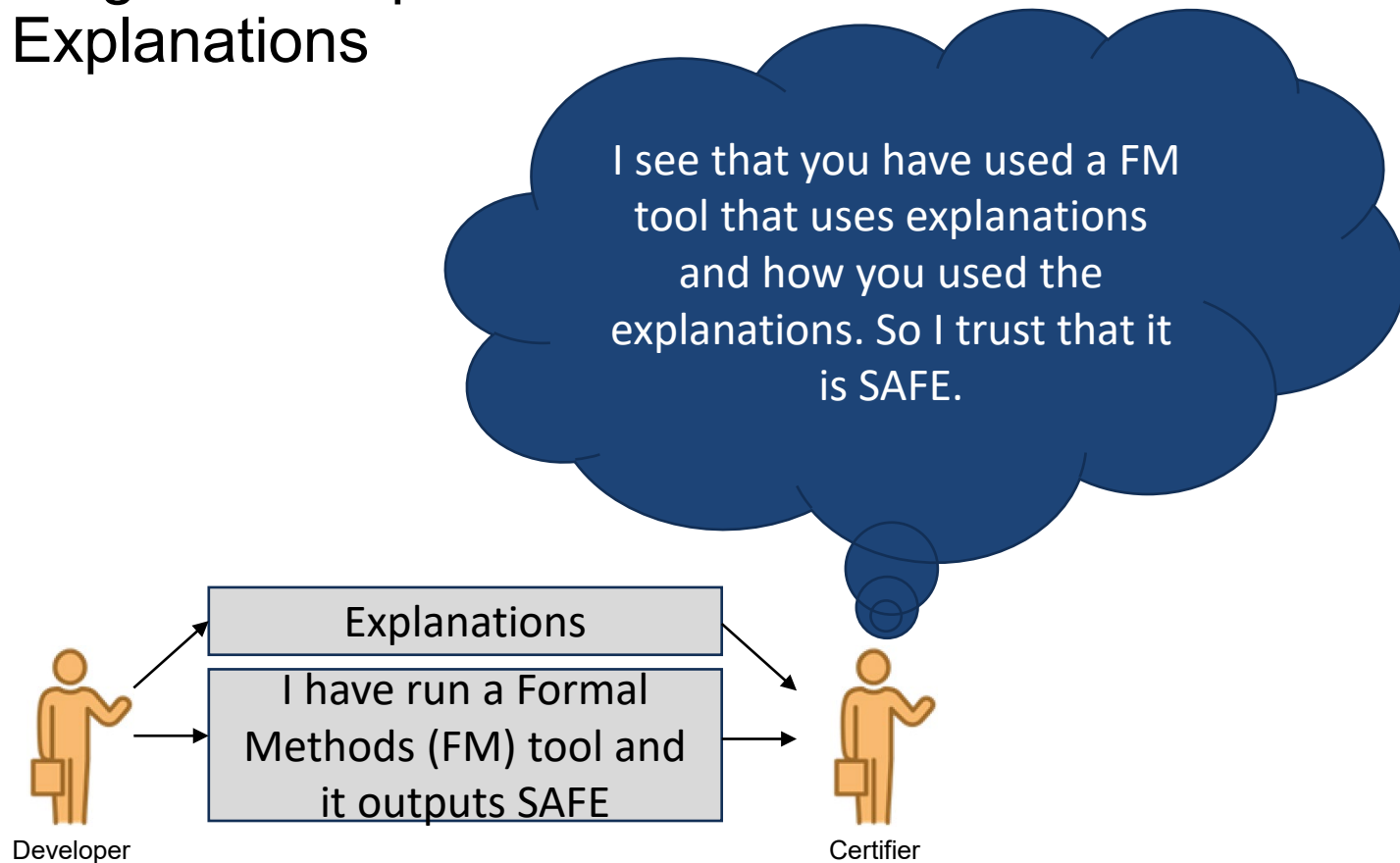


Developer

Now, I trust this result "Yes."

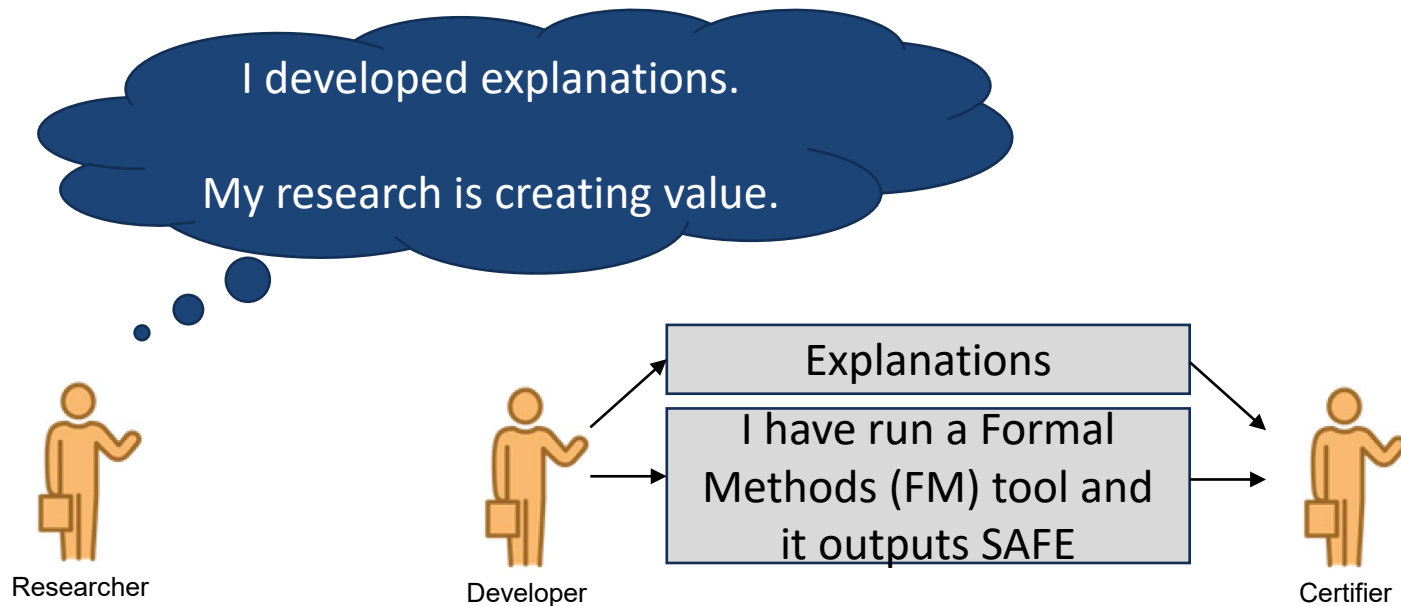
Problem: How to get developers and certifiers to trust FMs

Solution: Use Explanations



Problem: How to get developers and certifiers to trust FMs

Solution: Use Explanations





Conclusion

Explainability of real-time systems and their analysis:

has the potential to make research results valuable to practitioners

is a new, rich, area of research