The Case for Explainability of Real-Time Systems and their Analyses

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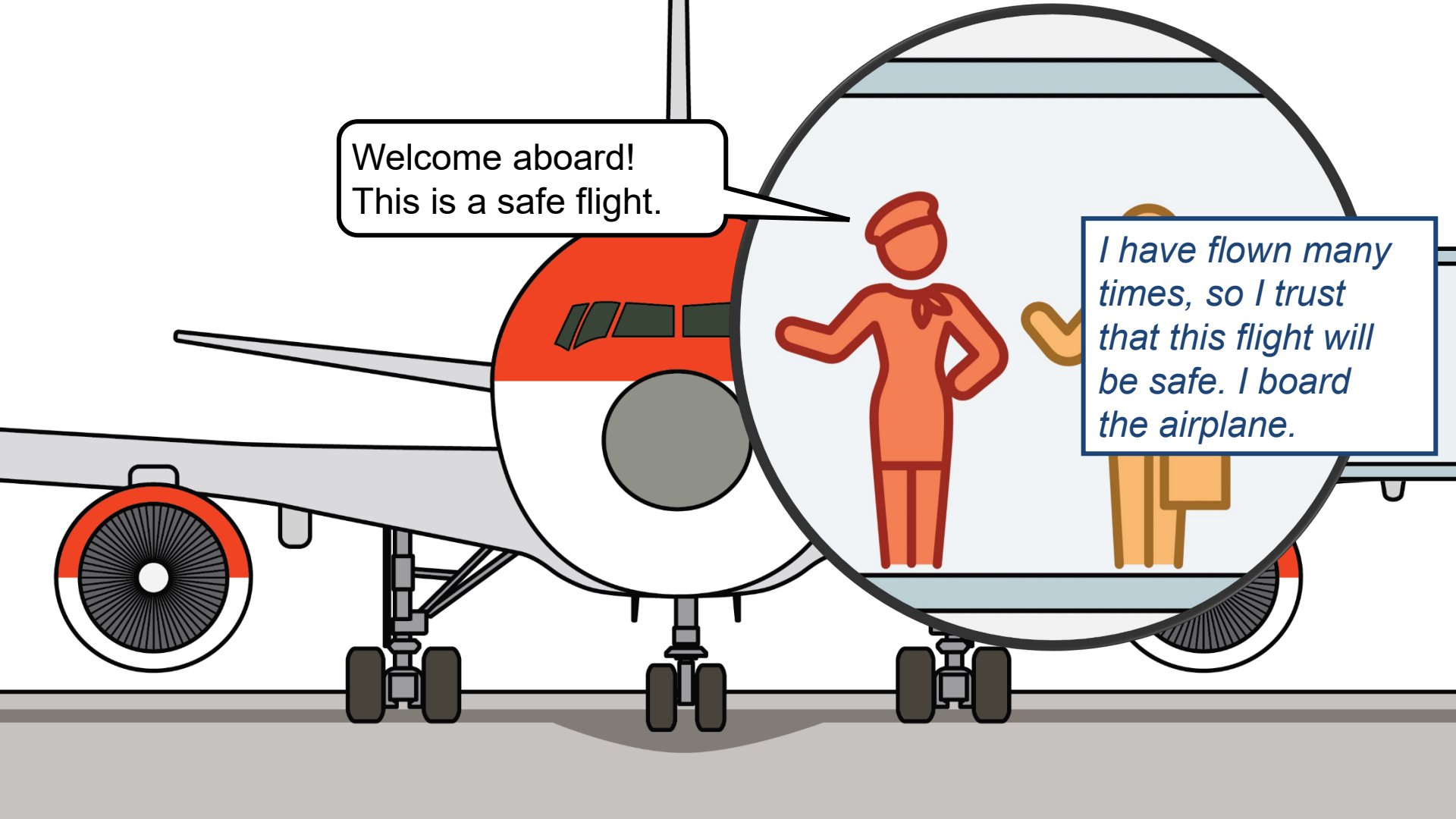
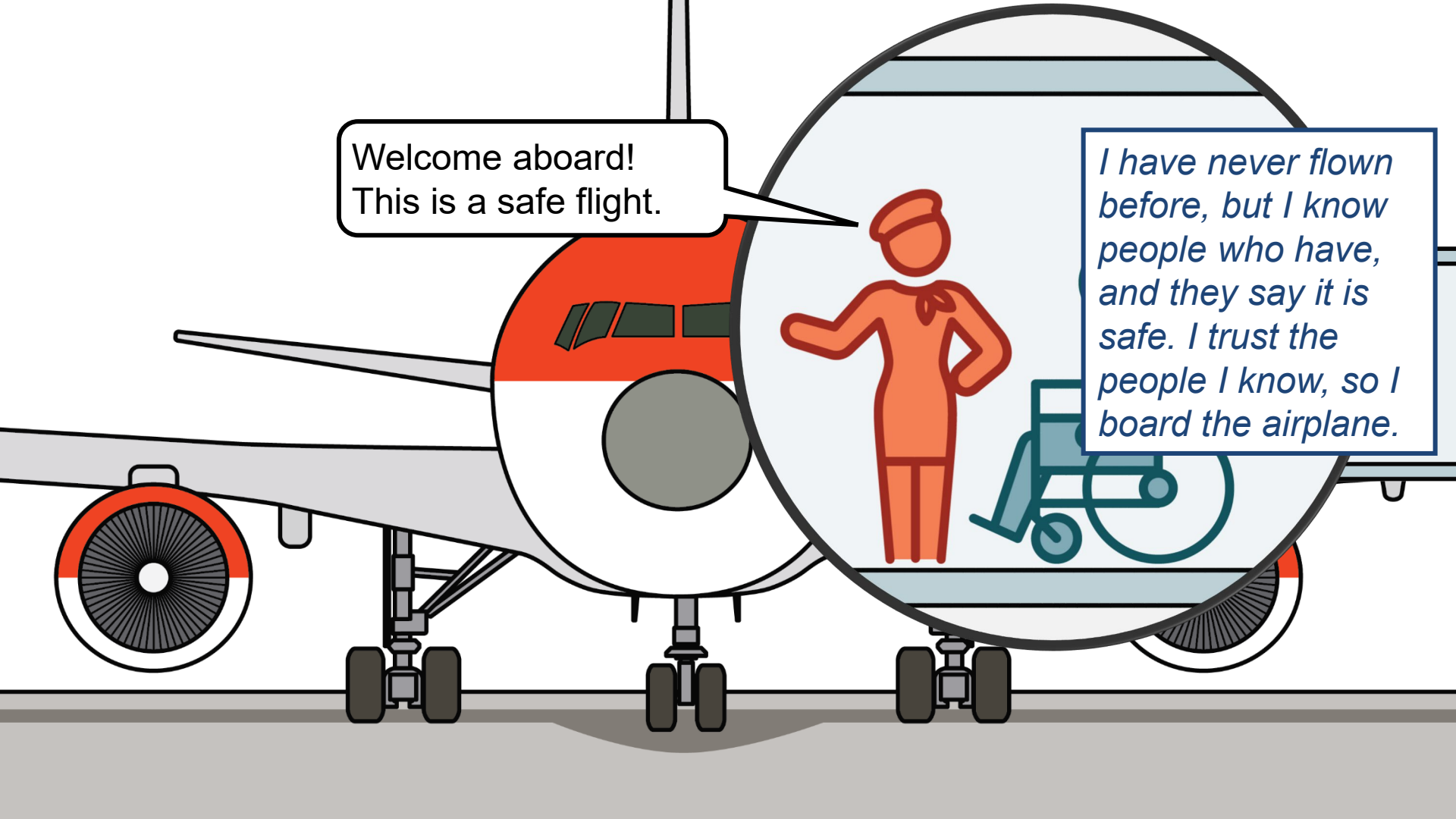
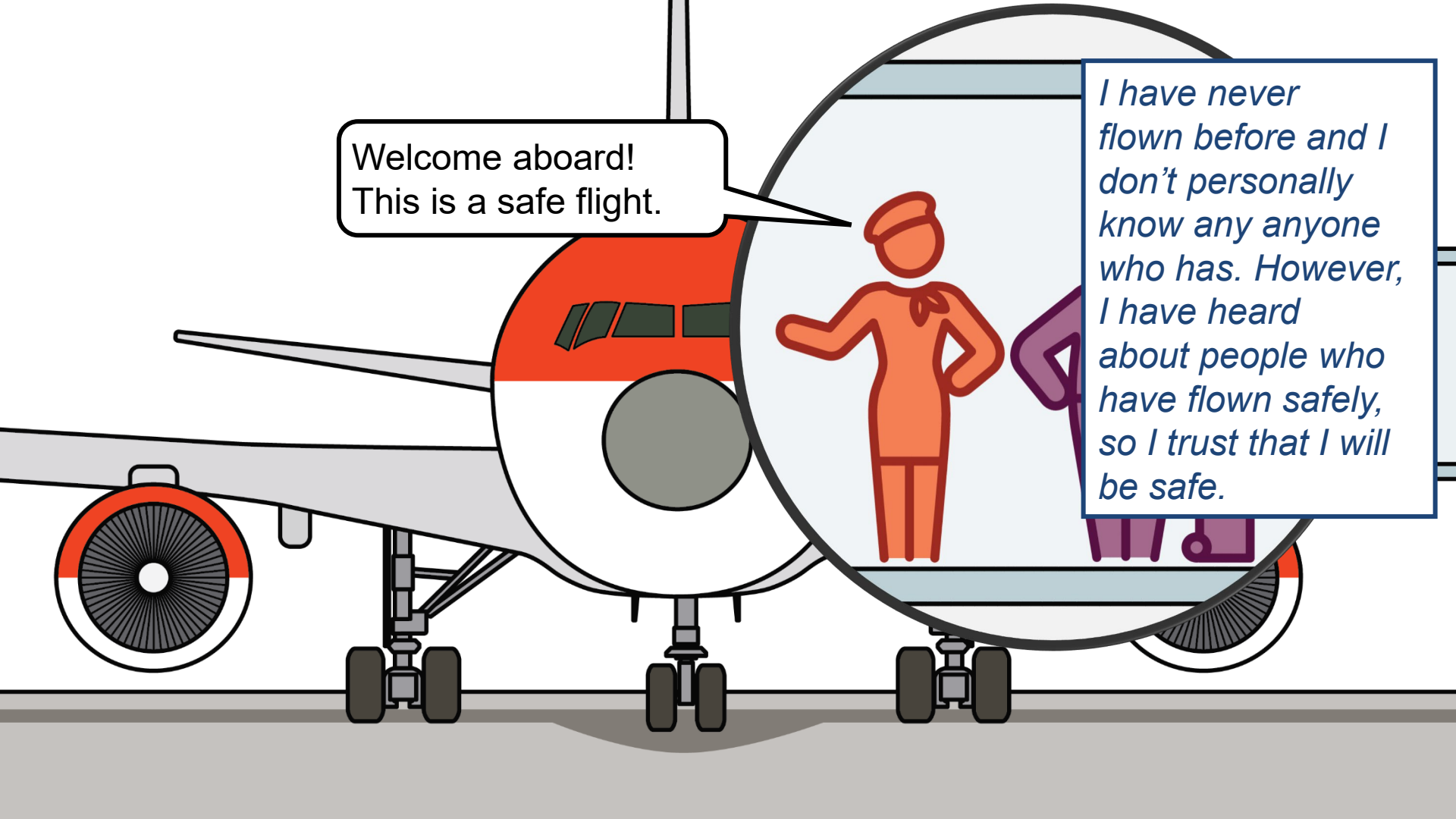
DM22-0896

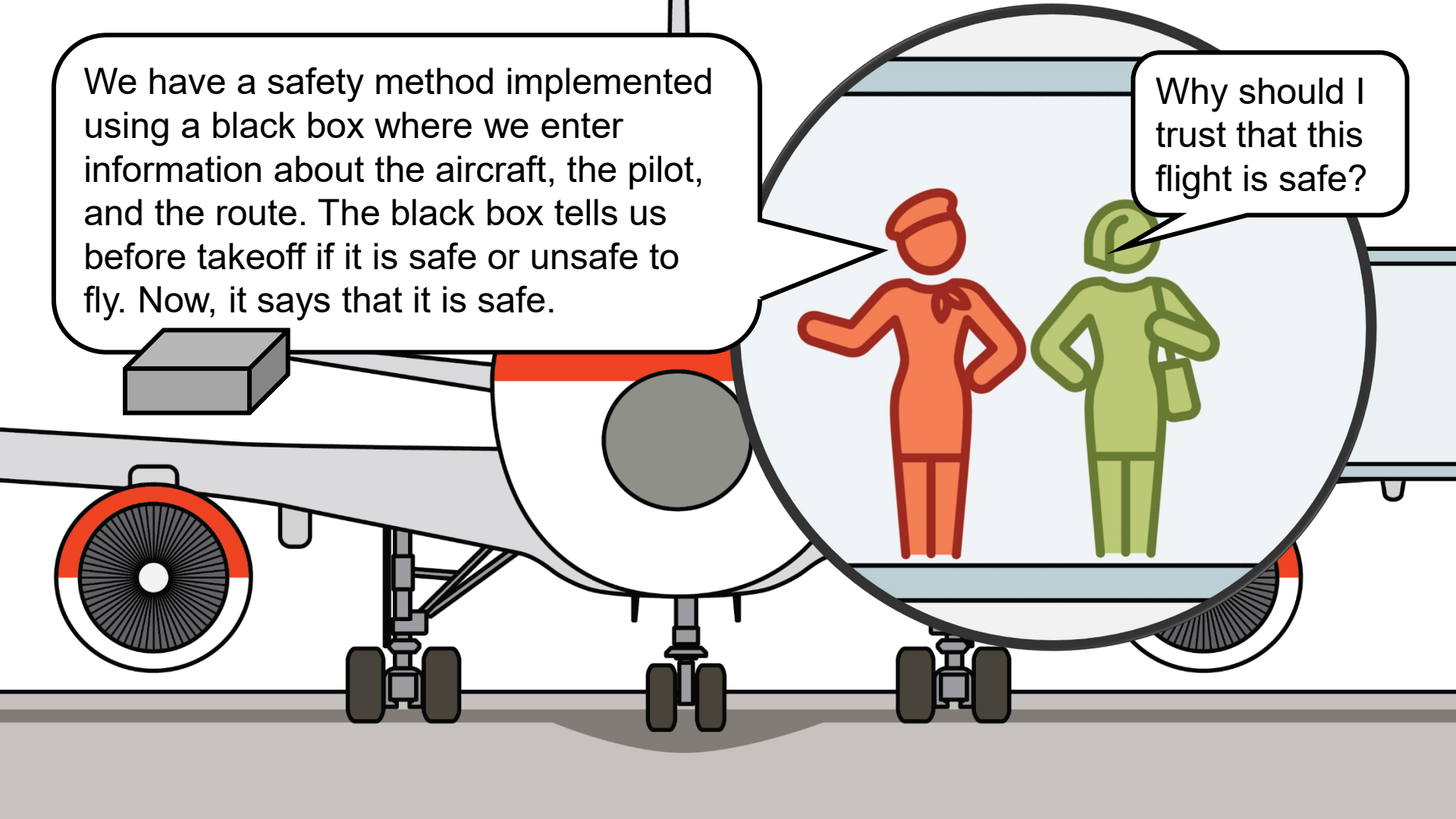
Motivation

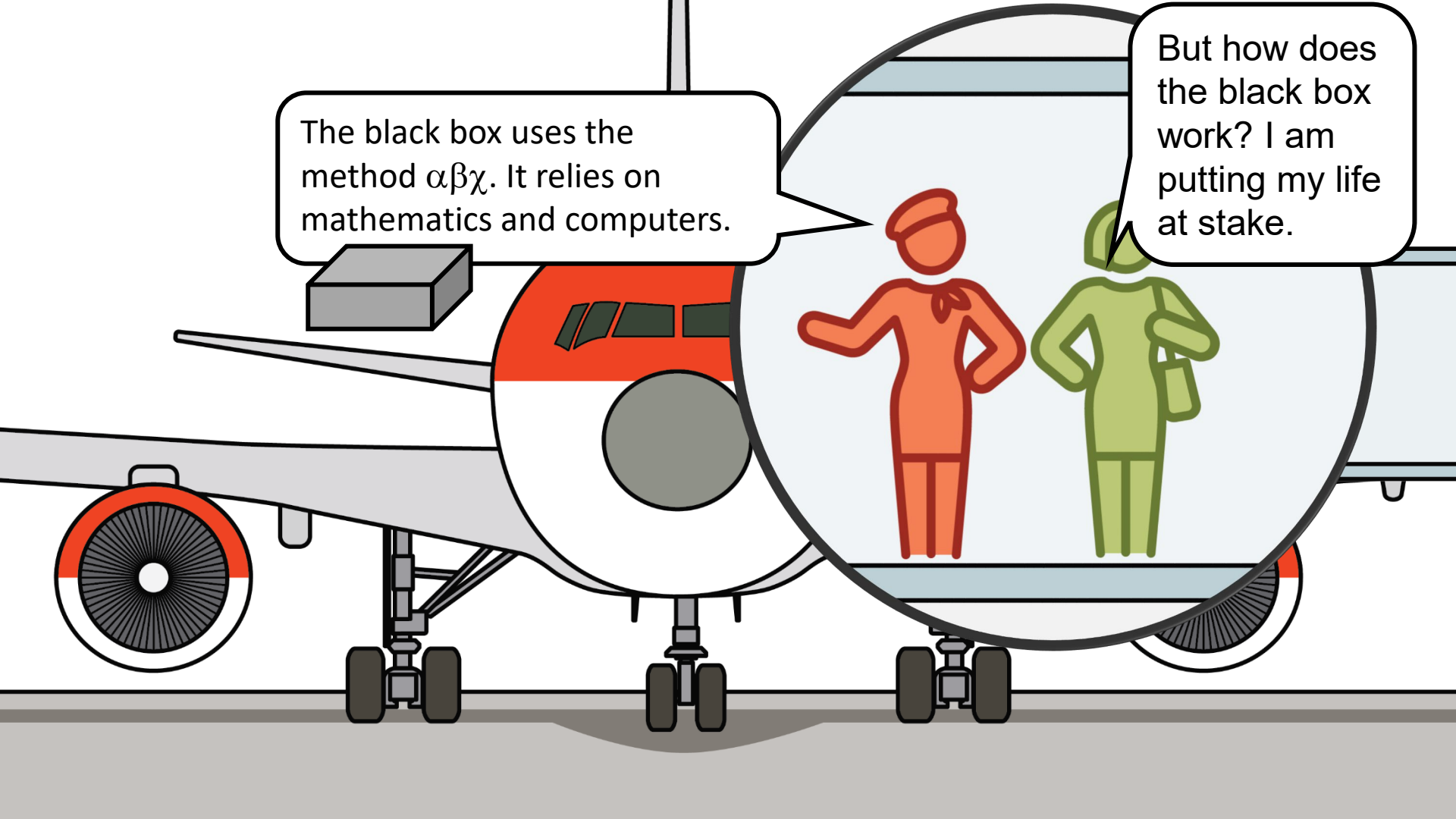

Welcome aboard!
This is a safe flight.
I have flown many times, so I trust that this flight will be safe. I board the airplane.

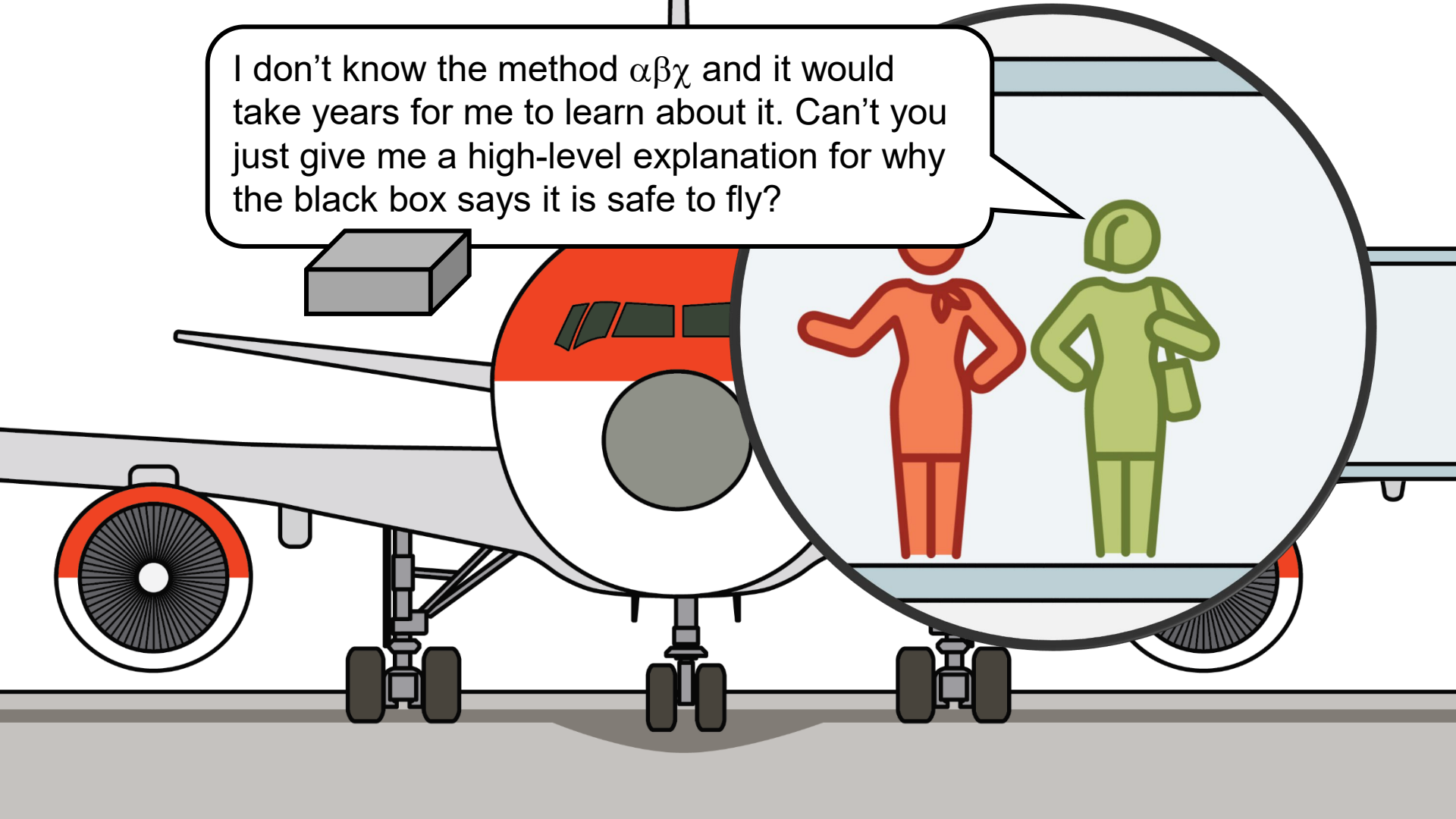

Welcome aboard!
This is a safe flight.
I have never flown before, but I know people who have, and they say it is safe. I trust the people I know, so I board the airplane.

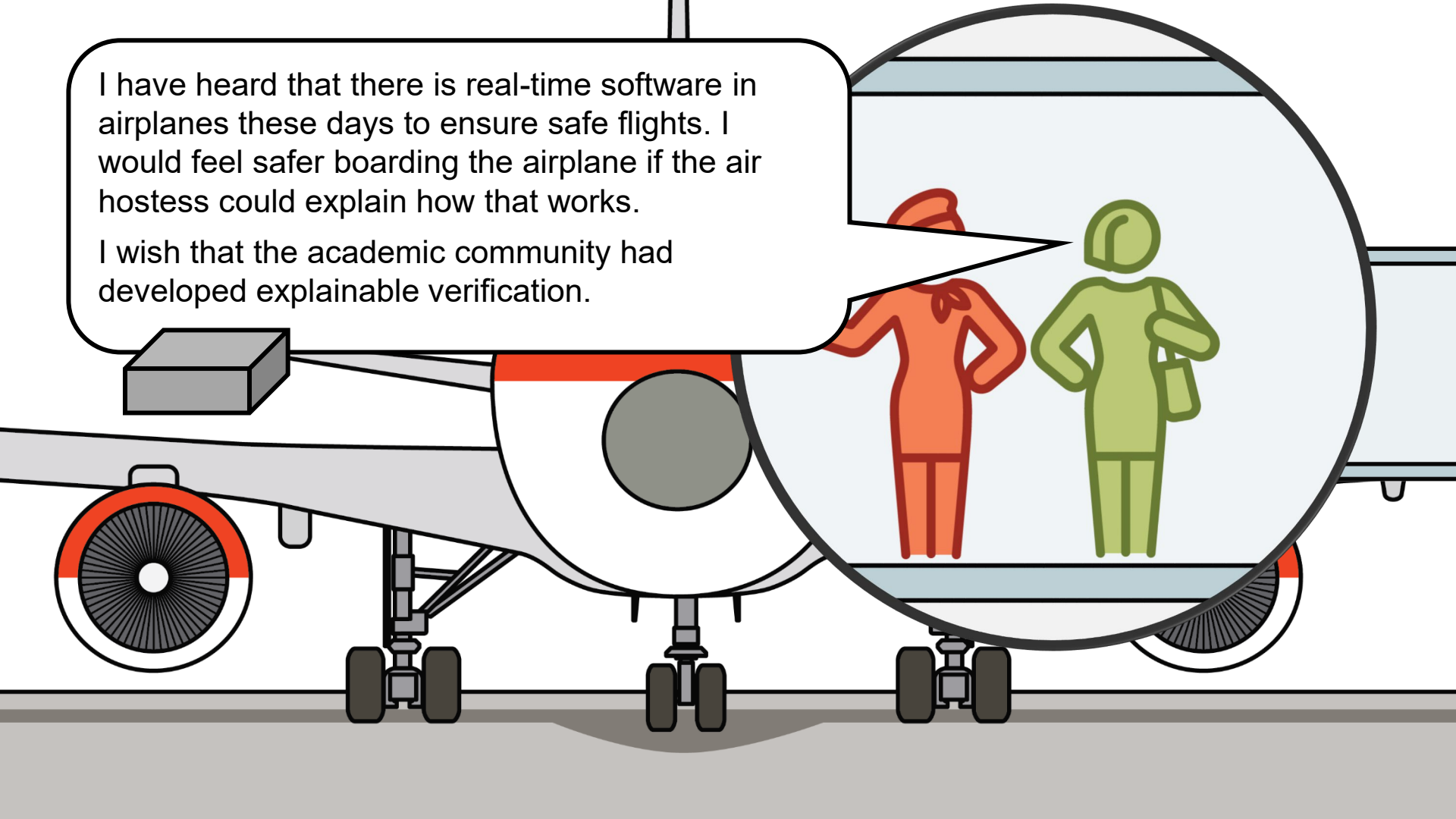

Welcome aboard!
This is a safe flight.
I have never flown before and I don't personally know any anyone who has. However, I have heard about people who have flown safely, so I trust that I will be safe.


Welcome aboard!
This is a safe flight.
I have never flown before, I don't know anyone who has flown before, and I've never heard of anyone who has flown safely


We have a safety method implemented using a black box where we enter information about the aircraft, the pilot, and the route. The black box tells us before takeoff if it is safe or unsafe to fly. Now, it says that it is safe.
Why should I trust that this flight is safe?


The black box uses the method $\alpha\beta\chi$. It relies on mathematics and computers.
But how does the black box work? I am putting my life at stake.


I don’t know the method $\alpha\beta\chi$ and it would take years for me to learn about it. Can’t you just give me a high-level explanation for why the black box says it is safe to fly?


I have heard that there is real-time software in airplanes these days to ensure safe flights. I would feel safer boarding the airplane if the air hostess could explain how that works.
I wish that the academic community had developed explainable verification.

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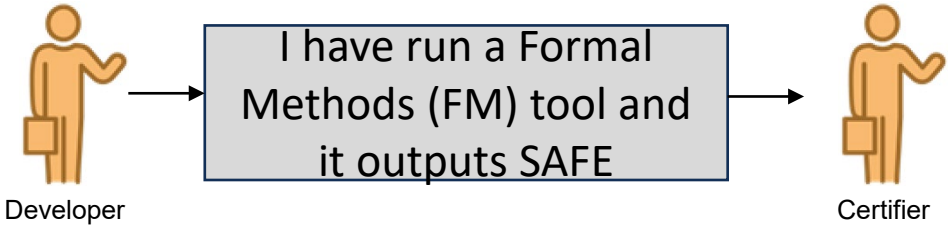

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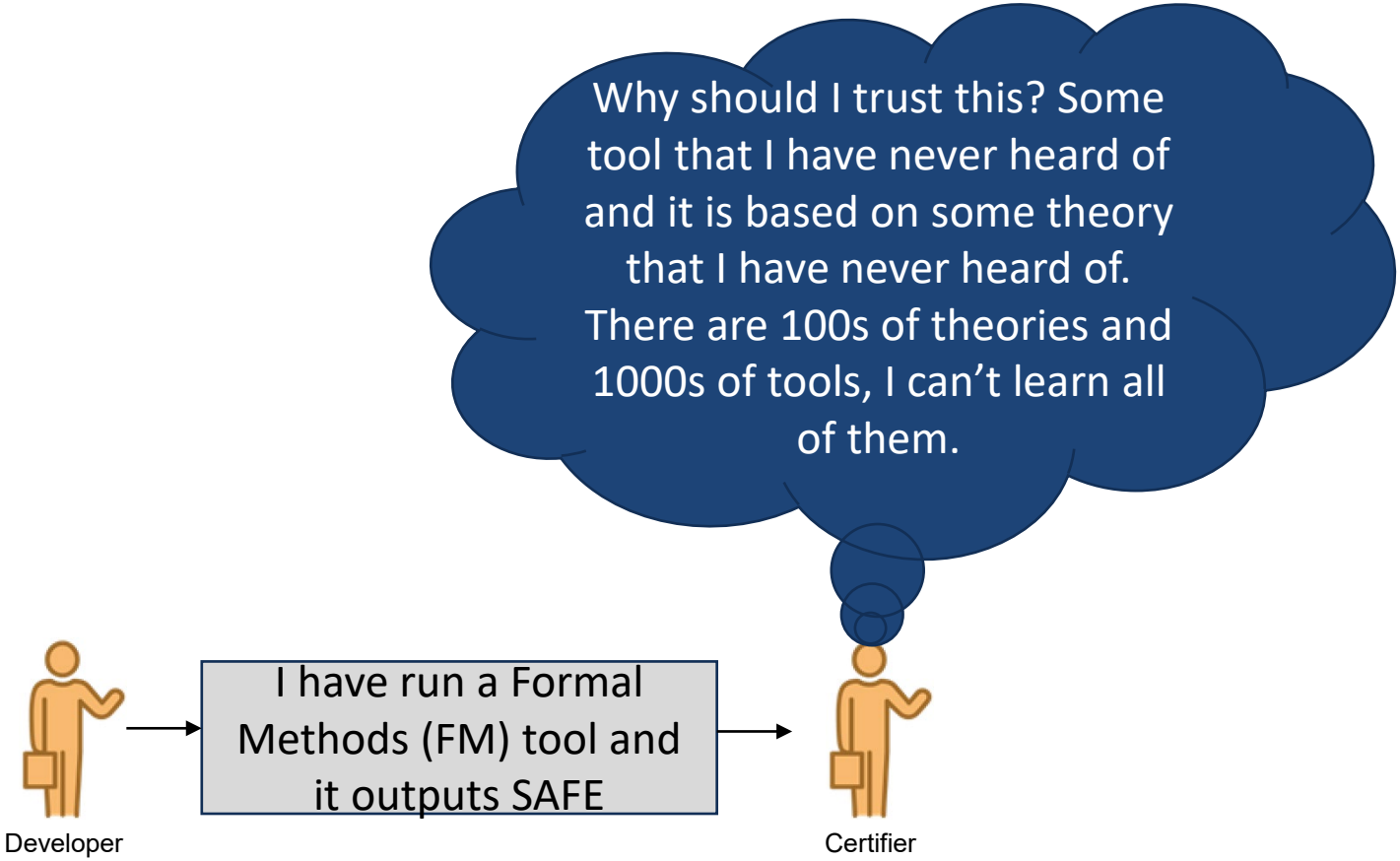

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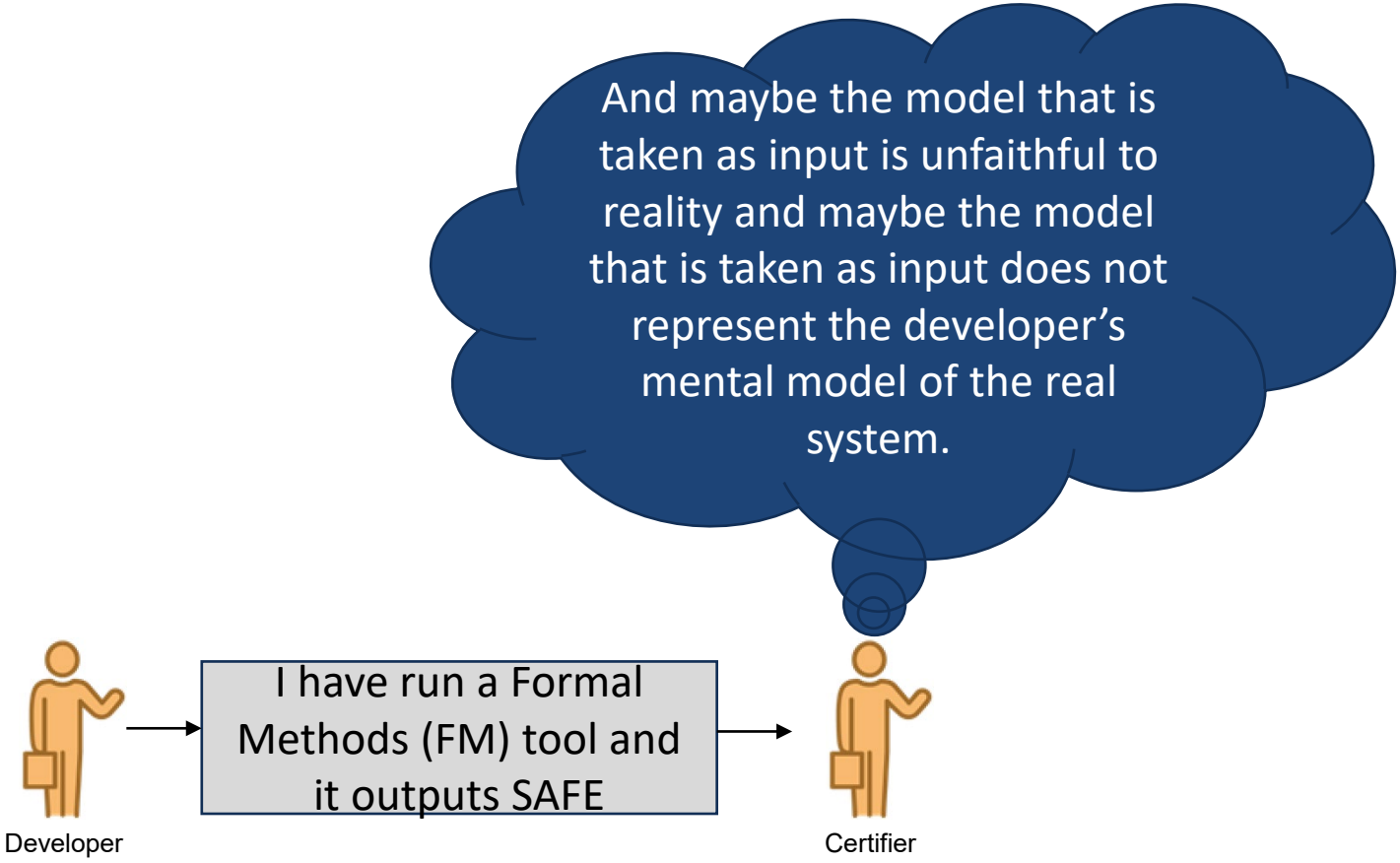
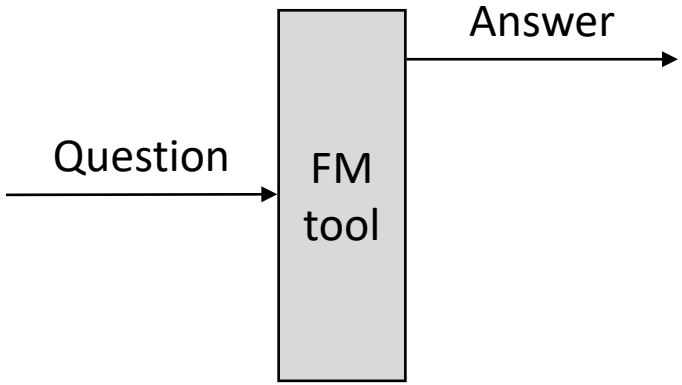


Explanation of Analysis

The Present: A World Without Explainability



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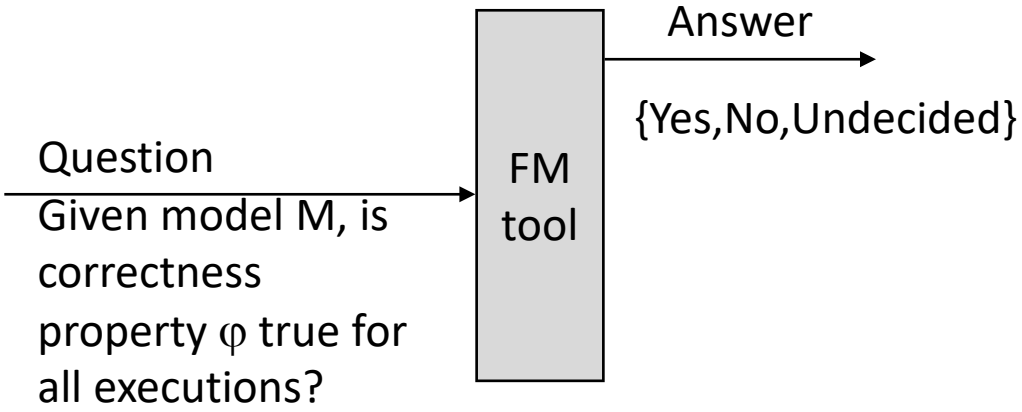

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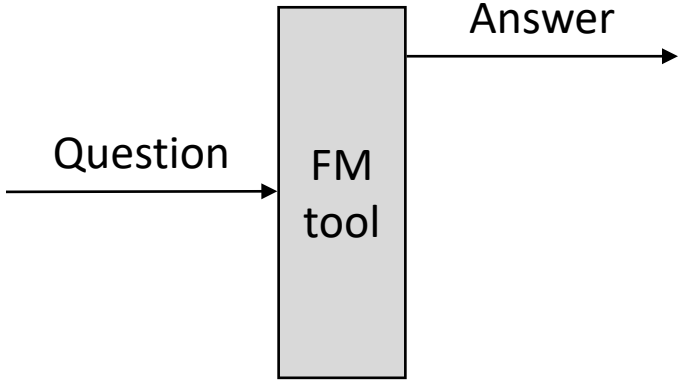

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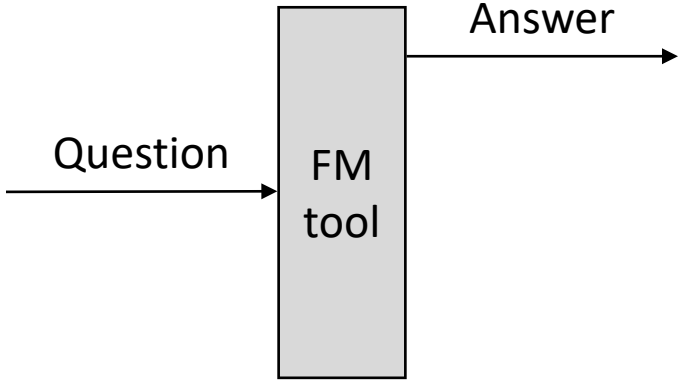

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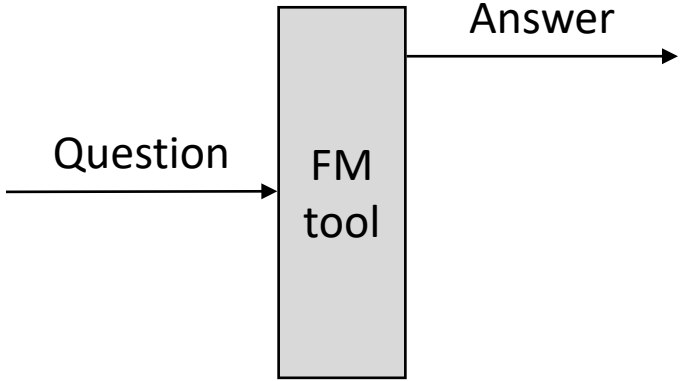

Answer: No

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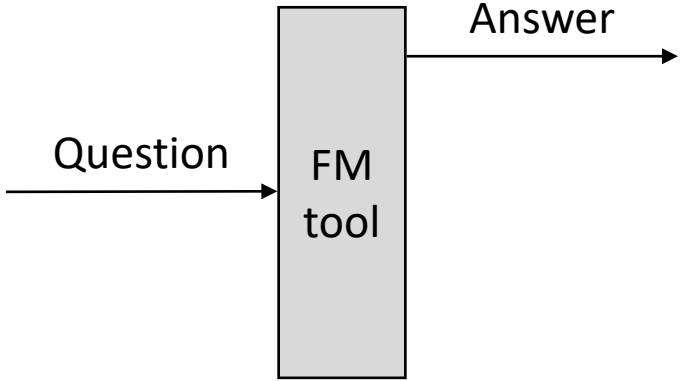
Question: Given program P and restriction R on inputs, what is its worst-case execution time?



Answer: 42 milliseconds

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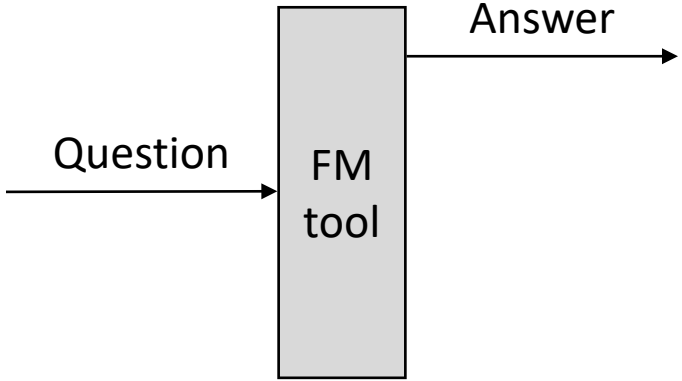

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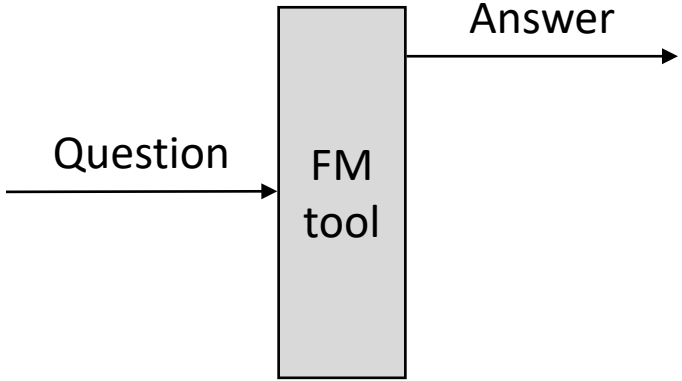

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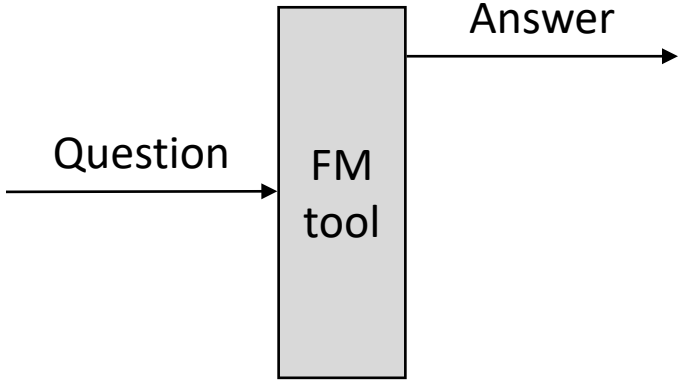

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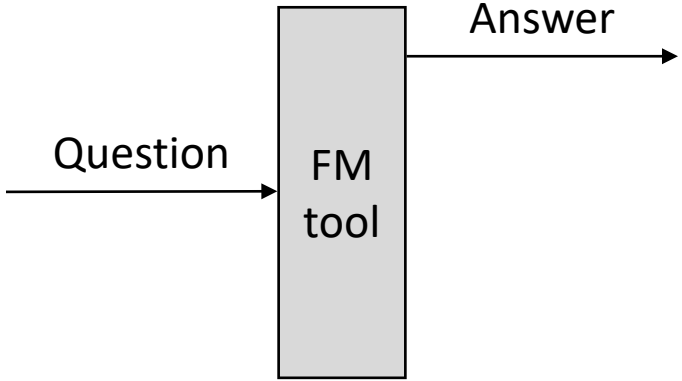

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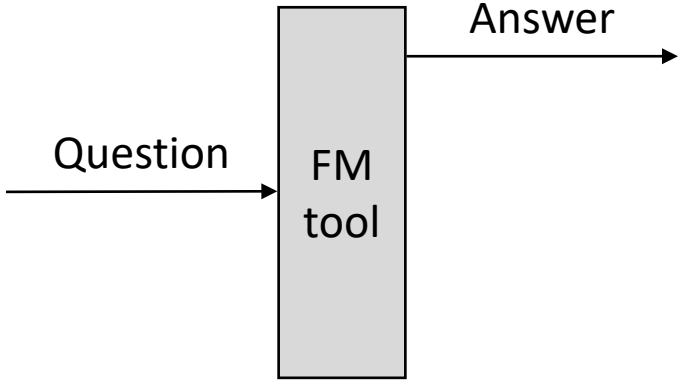

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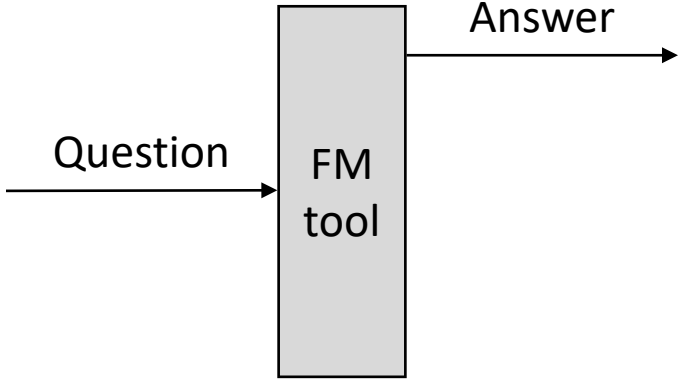

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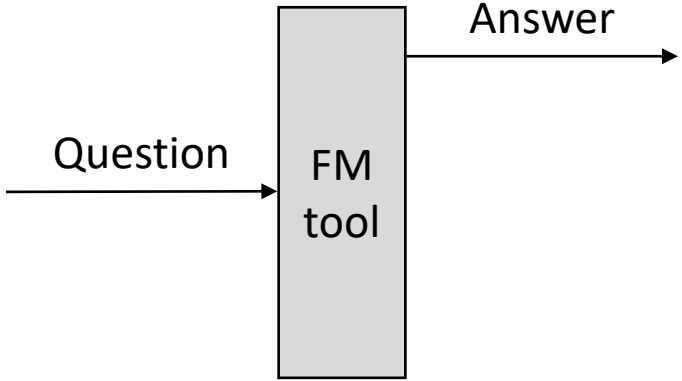

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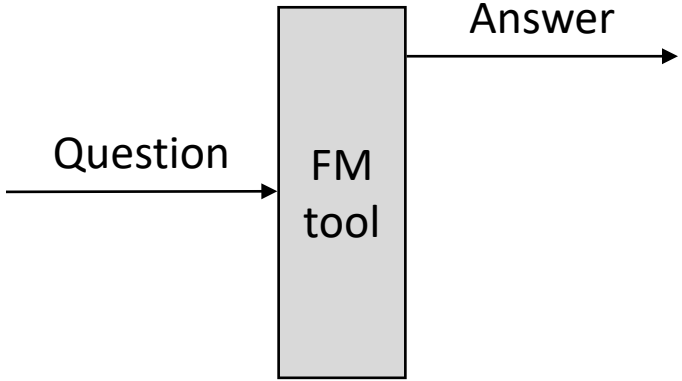

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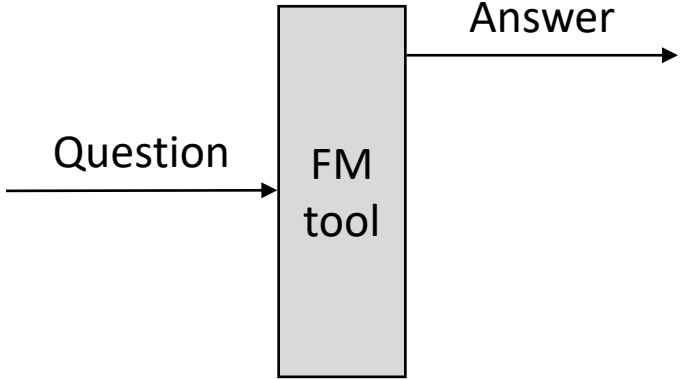

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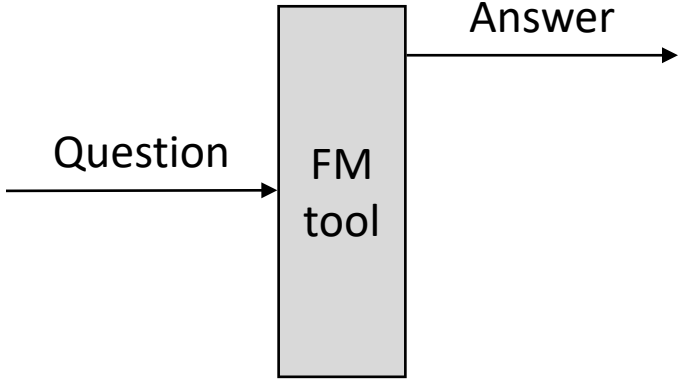

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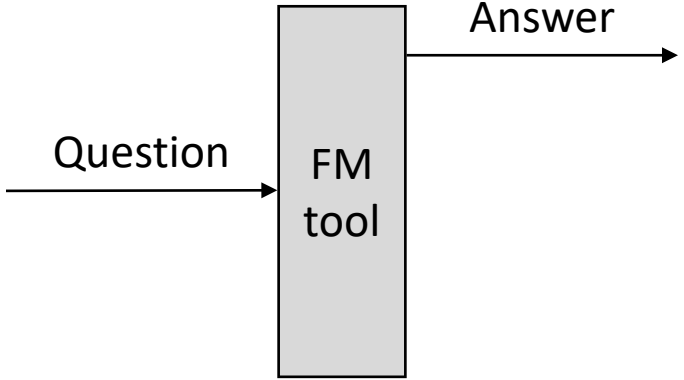

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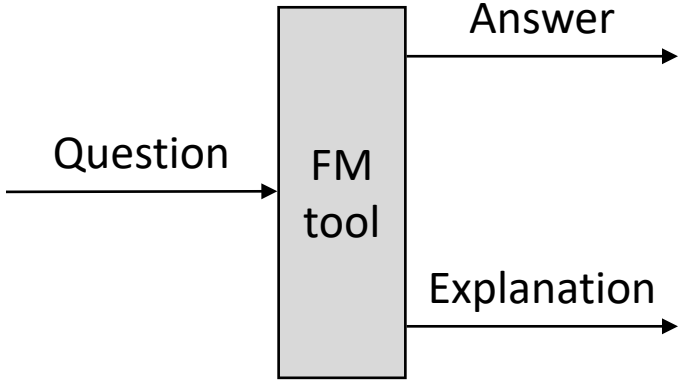
How much should one trust the FM tool that outputs answers?

If one does not know how the FM tool works, why should one trust it?

If practitioners do not trust the FM tools that researchers produce, how much value is there in the research?

A Potential Future: A World With Explainability

Question: Given taskset τ, the following run-time scheduler S, and assumptions A, is it schedulable?



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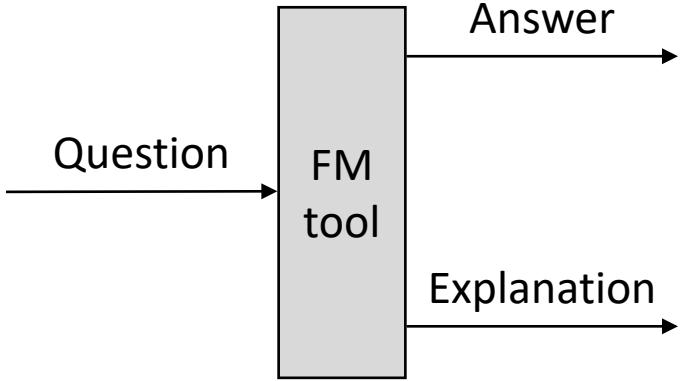

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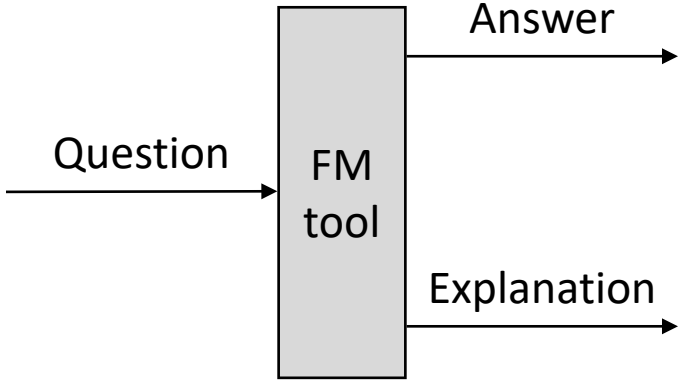

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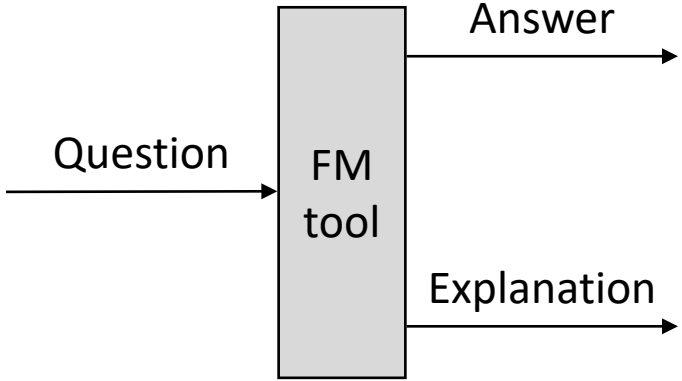

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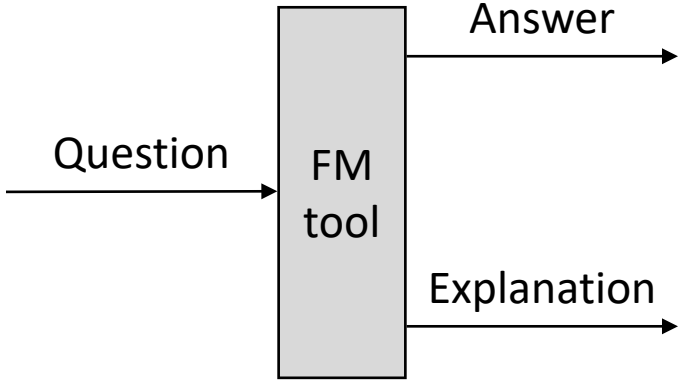

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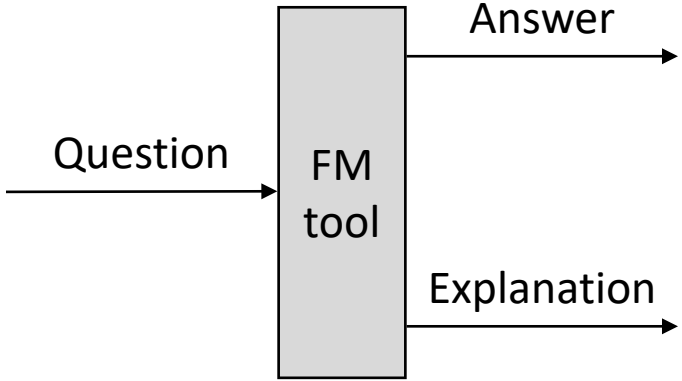

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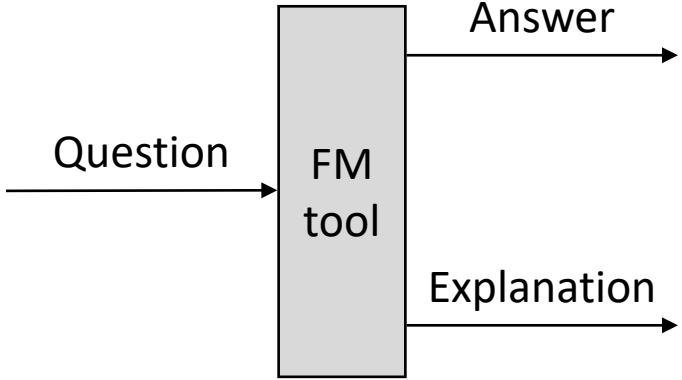

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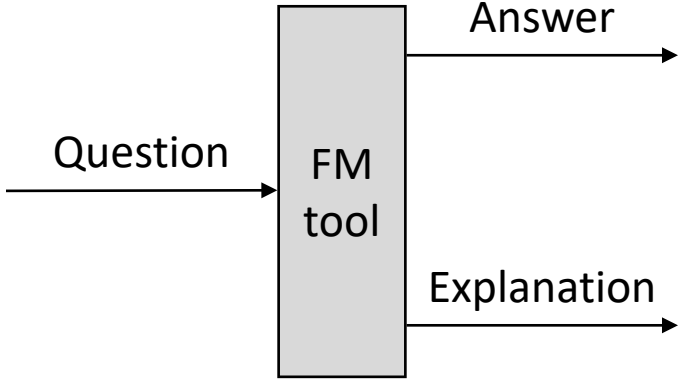

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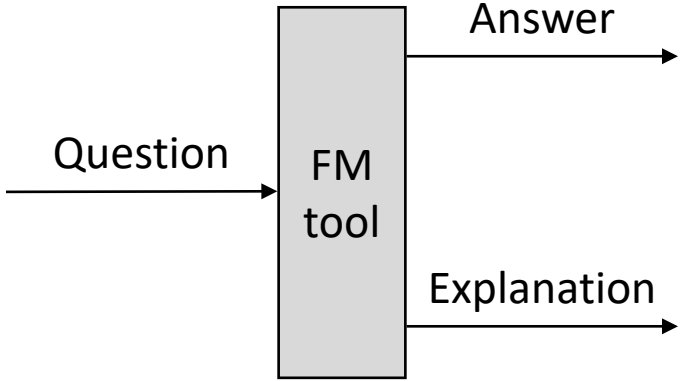

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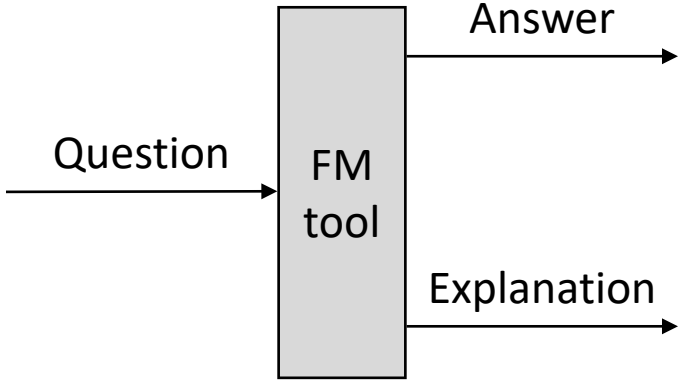

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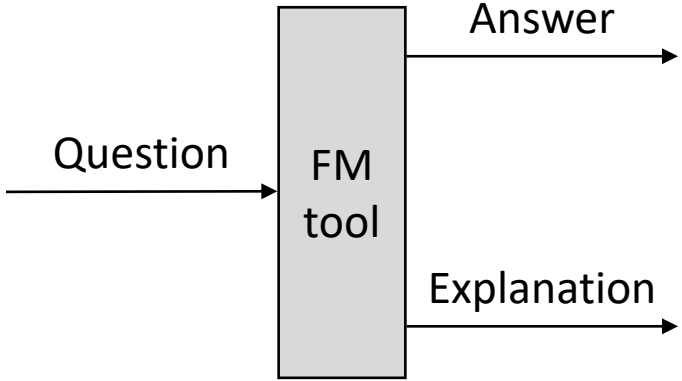
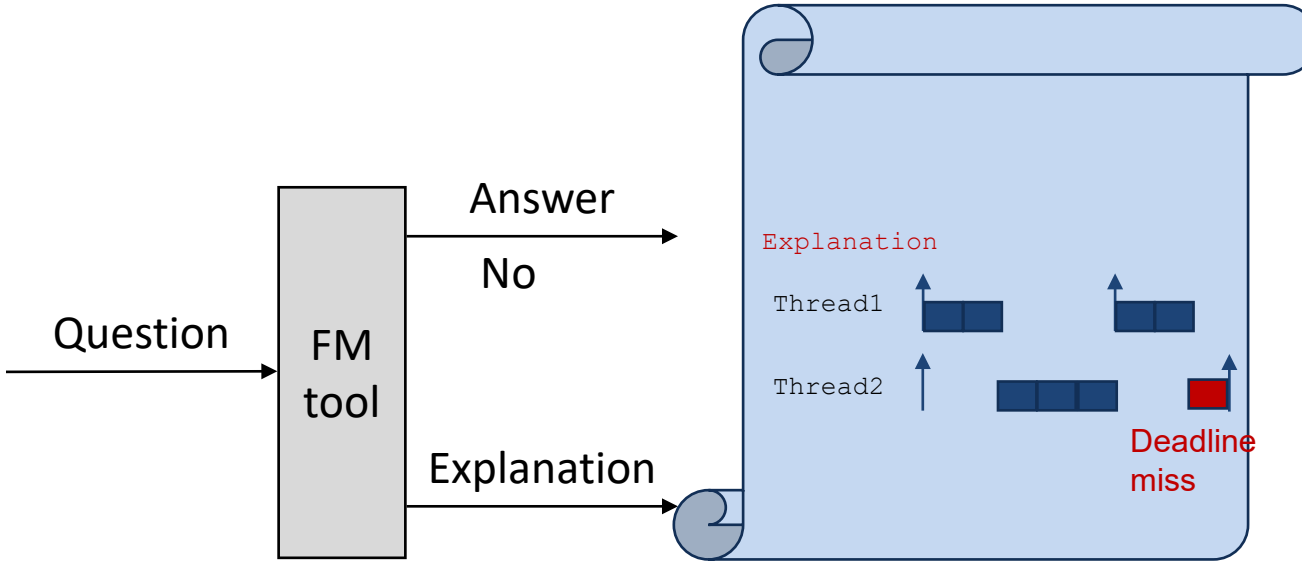


Answer: Yes

- Explanability has the potential to make practitioners trust FM tools
- Explanability has the potential to make sure our research becomes more valuable to practitioners

What Could an Explanation Look Like?

Question: Given taskset τ, the following run-time scheduler S, and assumptions A, is it schedulable?



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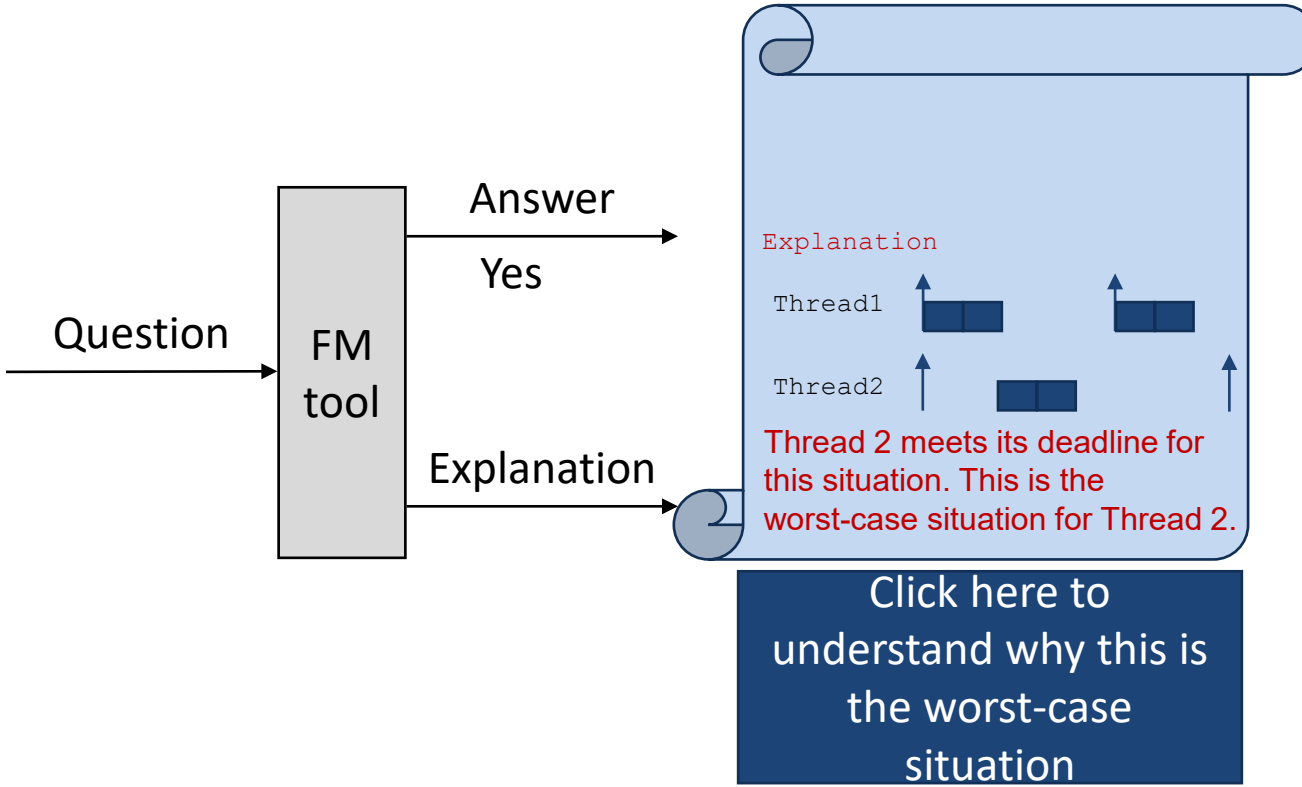

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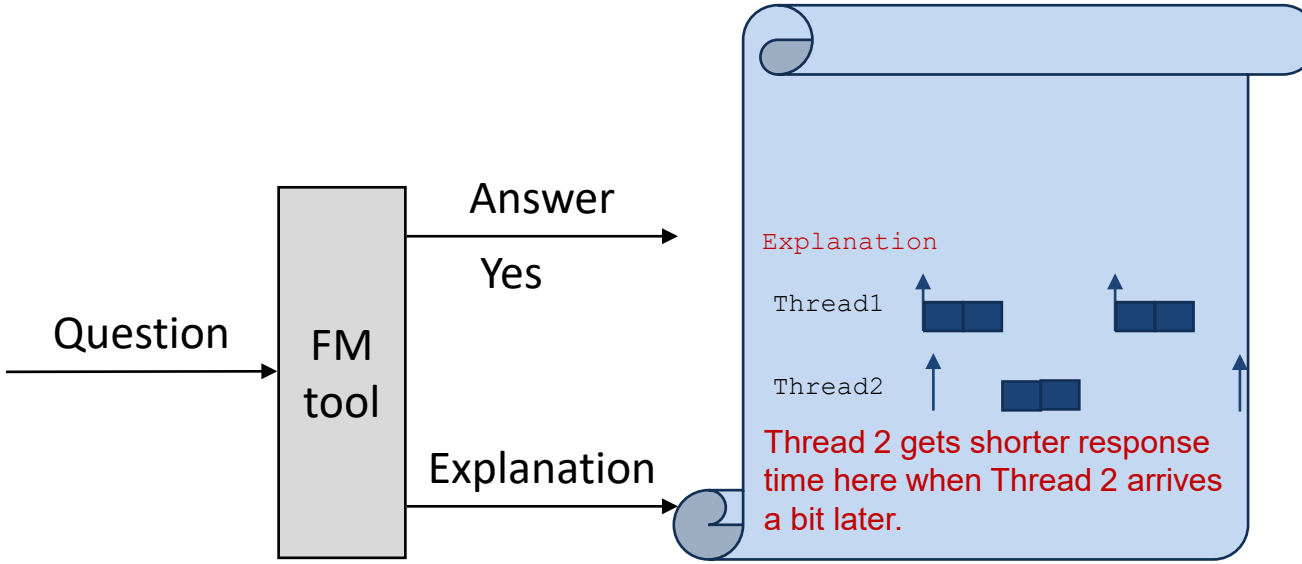

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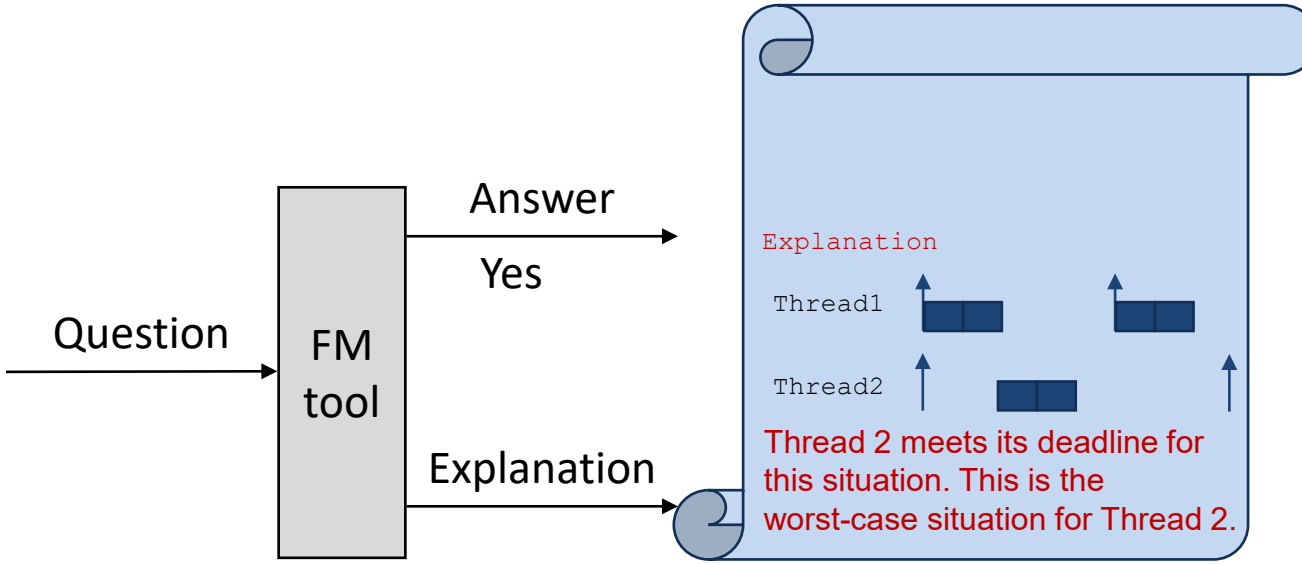

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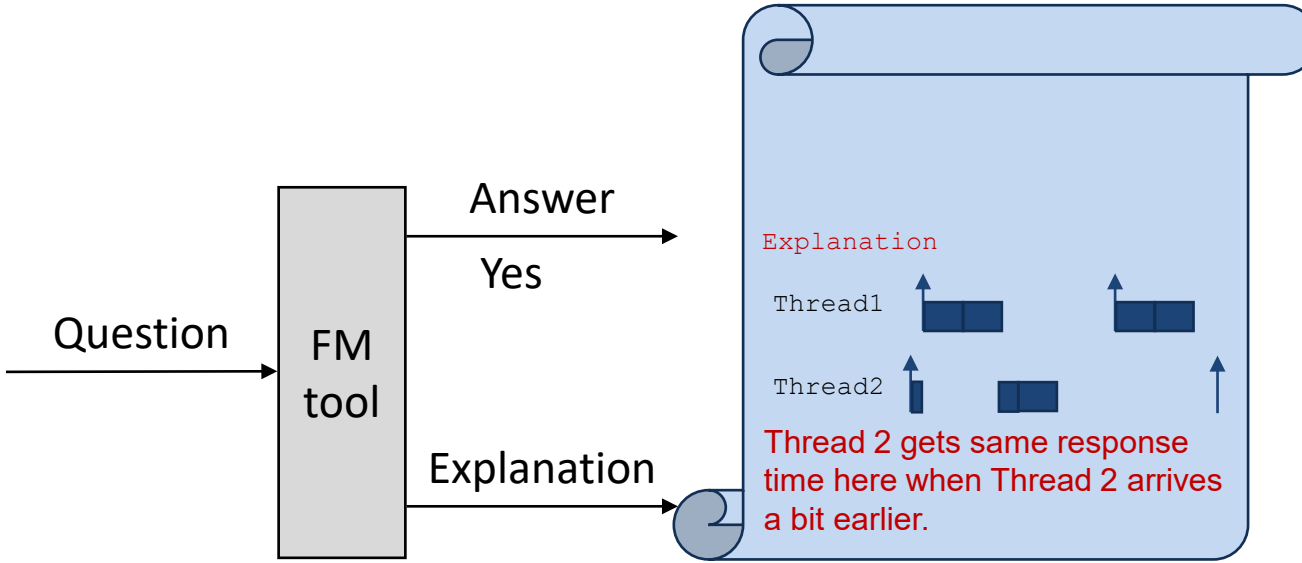

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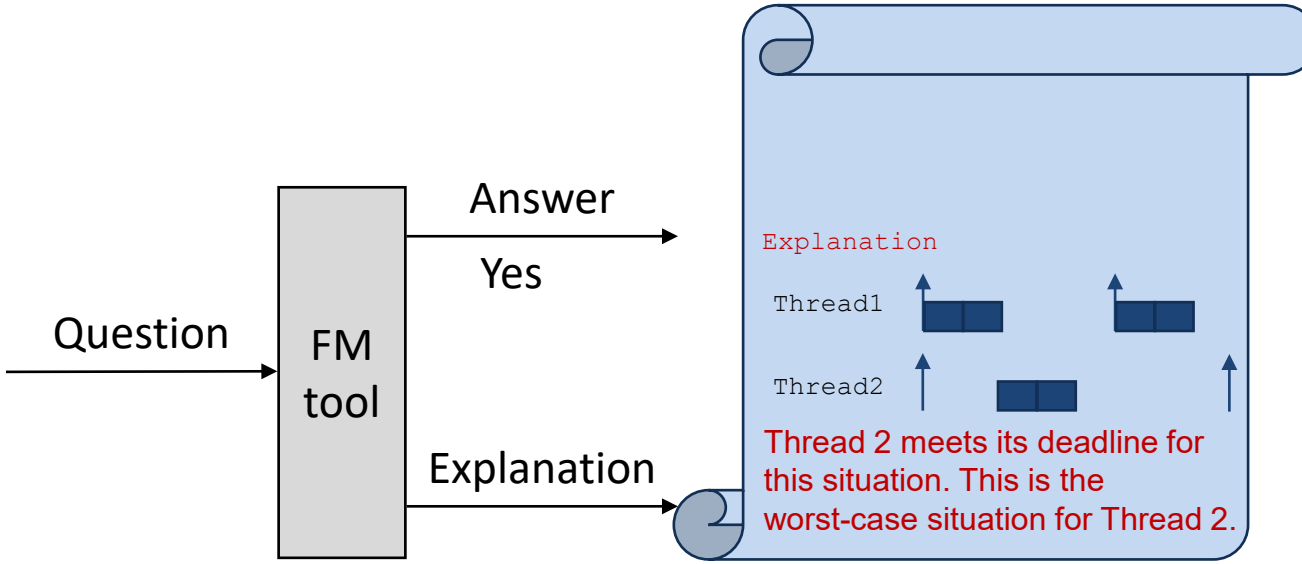

Explanation of Input

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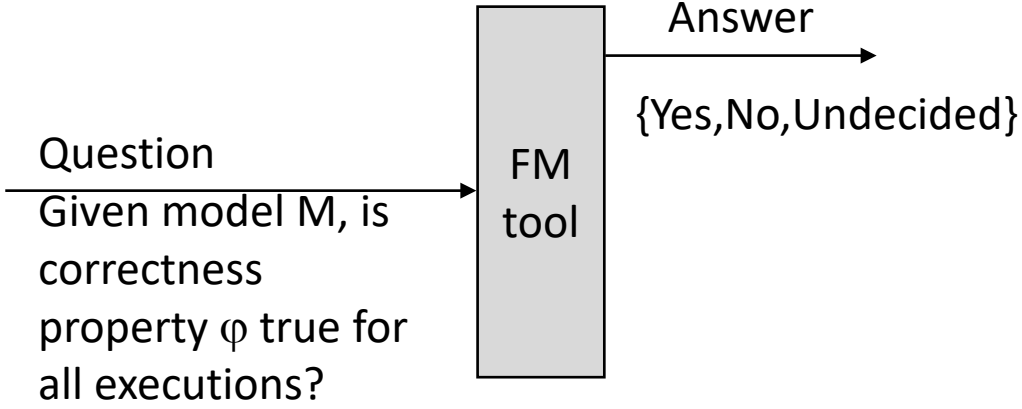

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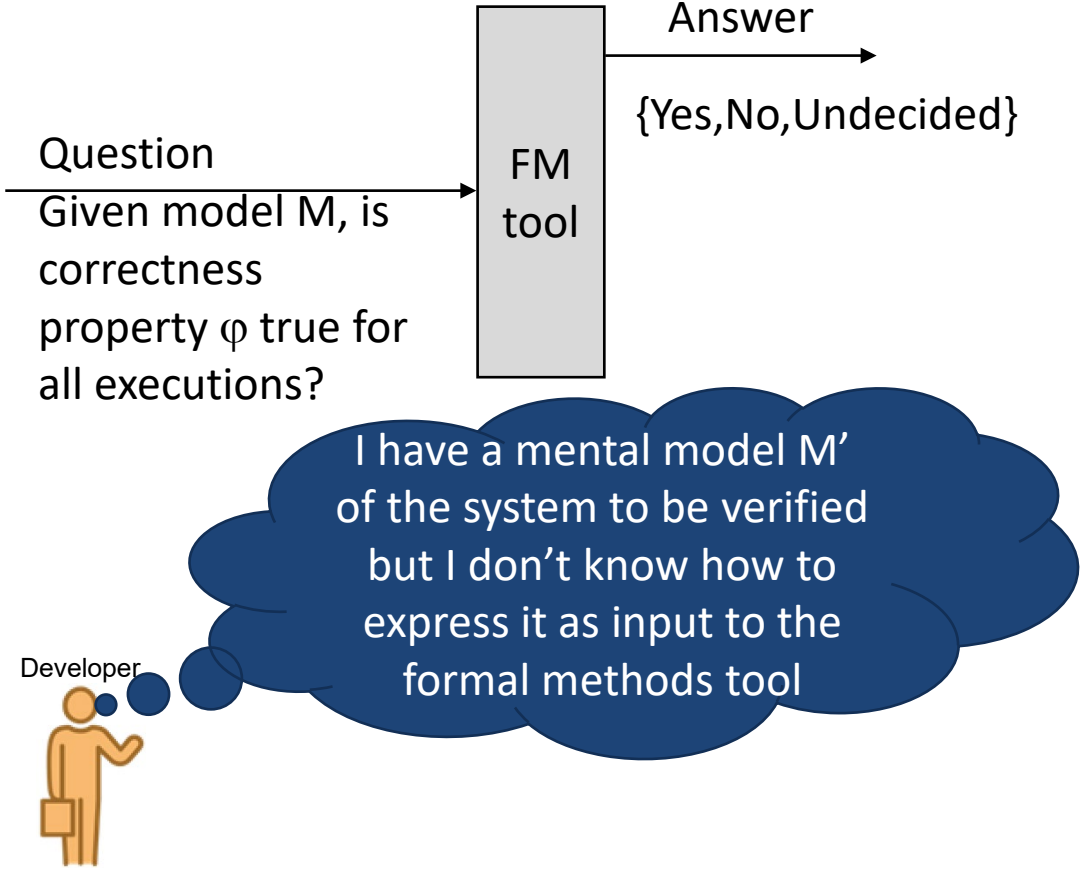
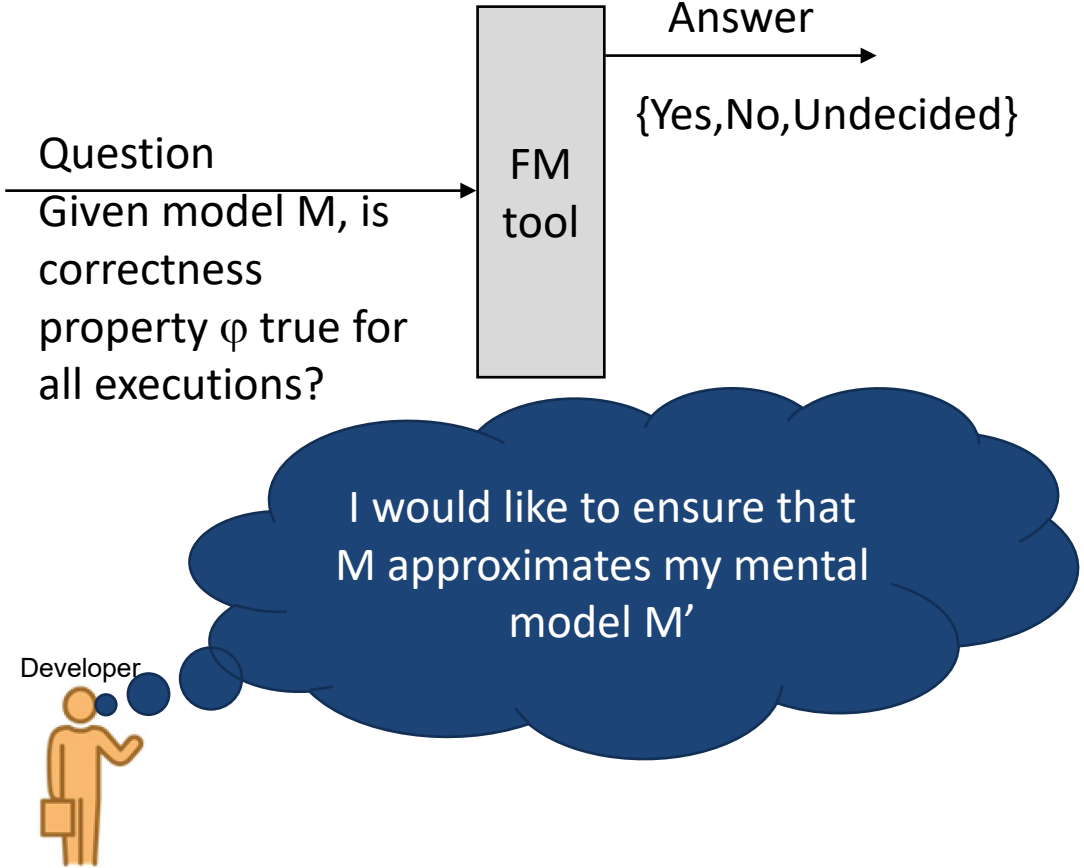
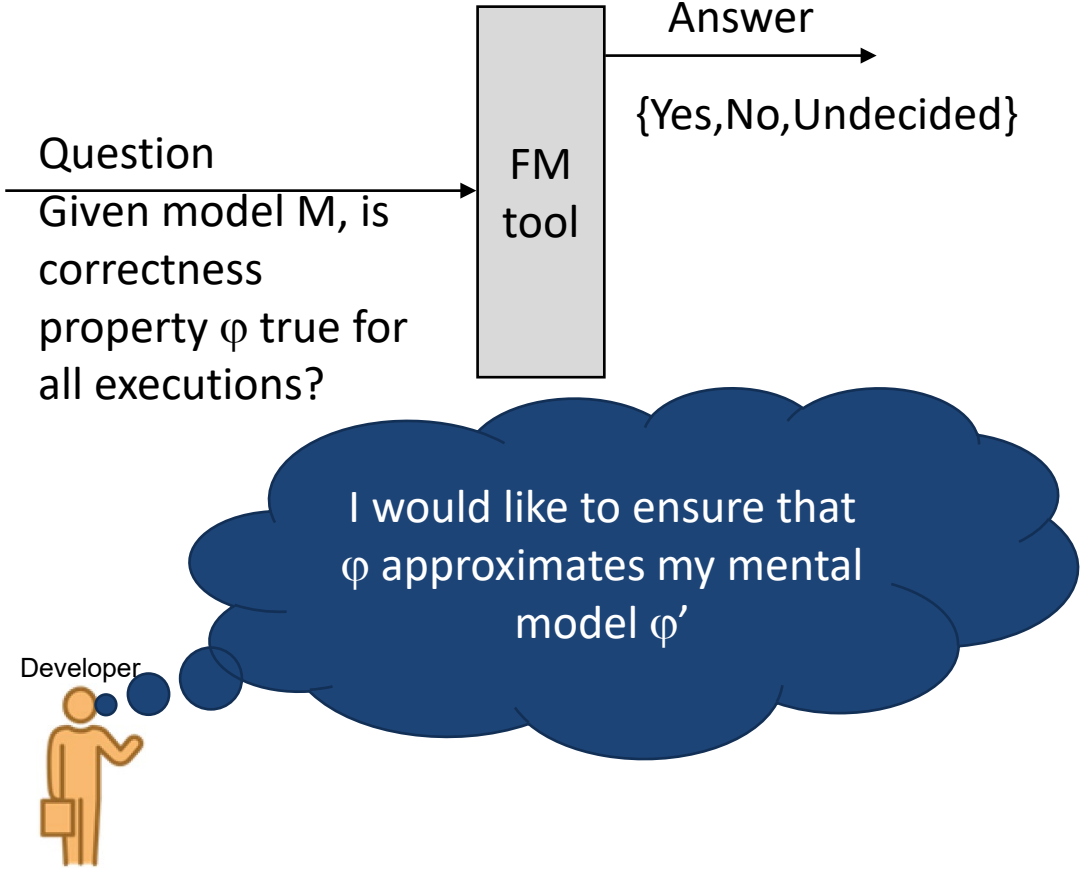


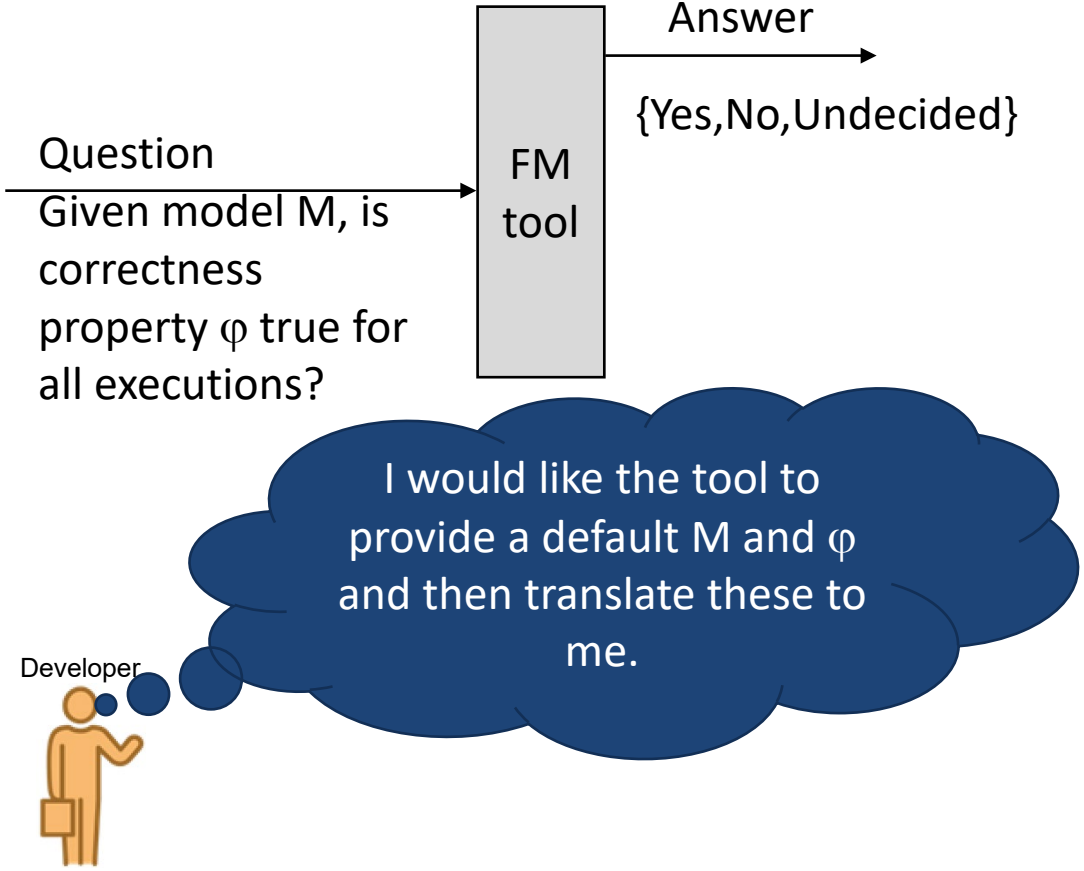
Explainability of Input



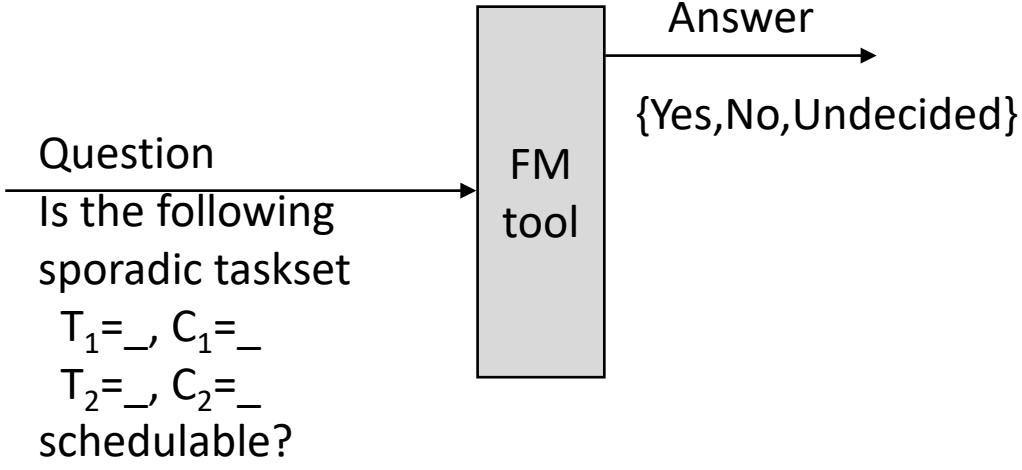
Explainability of Input



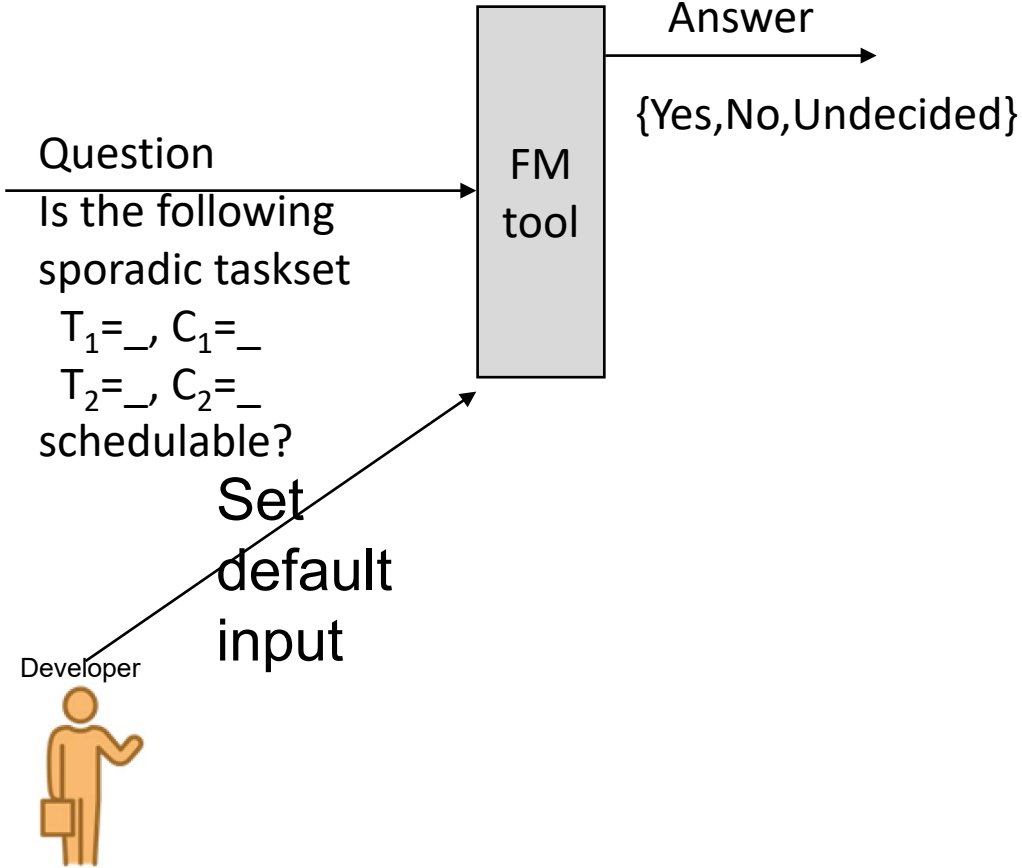
Explainability of Input



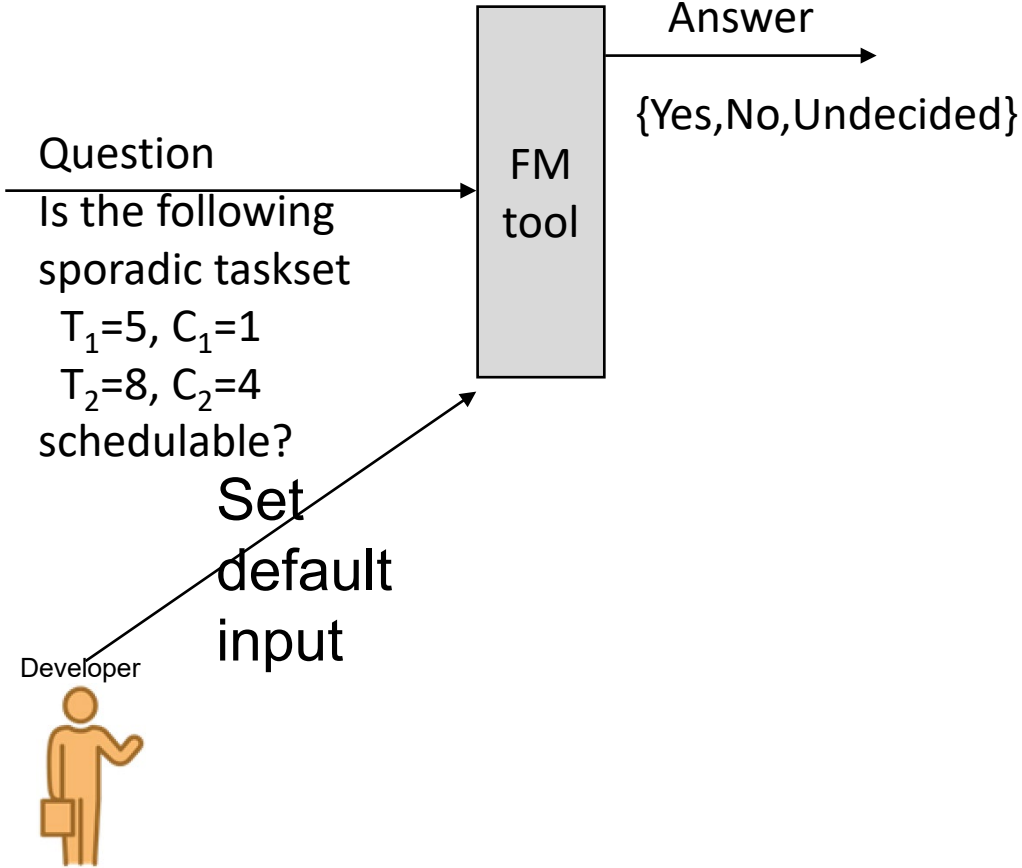
Explainability of Input: Real-Time Scheduling



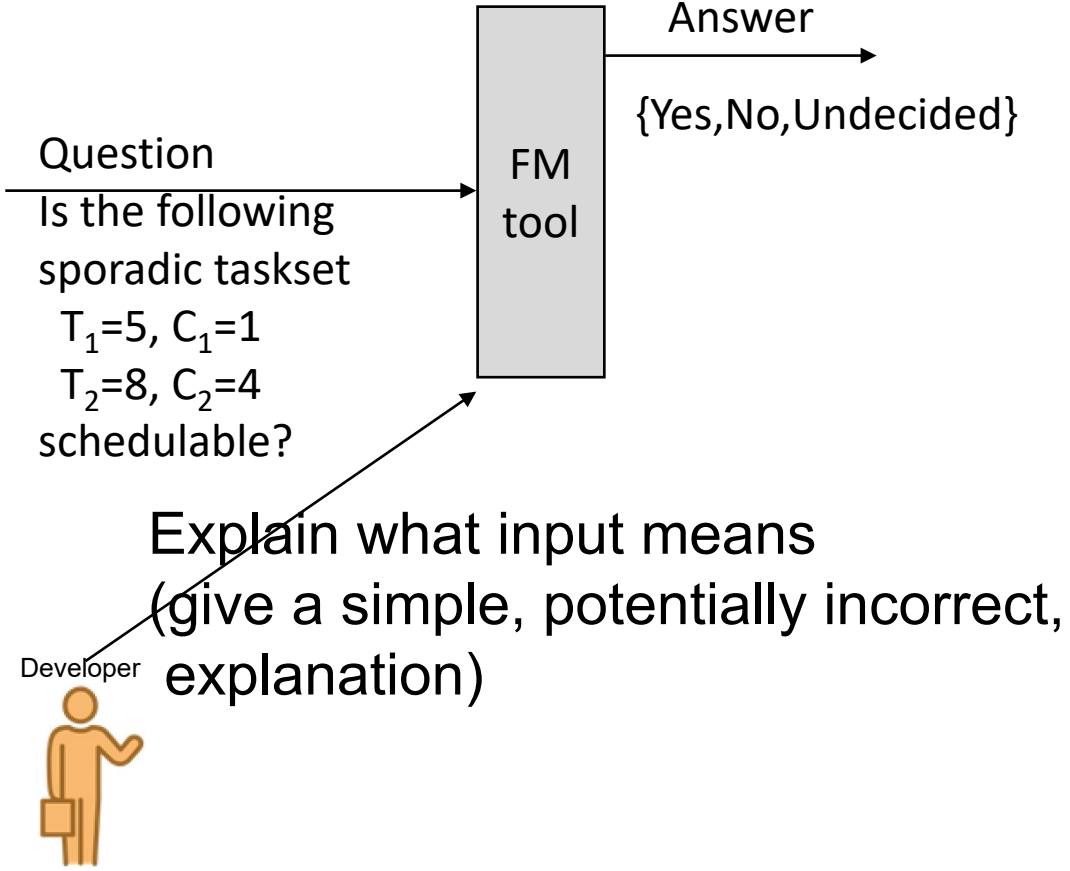
Explainability of Input: Real-Time Scheduling



Explainability of Input: Real-Time Scheduling



Explainability of Input: Real-Time Scheduling



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5
  while (1) {
    delay_until(t)
    do_work(1)
    t=t+5
  }

Thread 2
  t=now()+8
  while (1) {
    delay_until(t)
    do_work(4)
    t=t+8
  }
```

Question

Is the following sporadic taskset

T_1=5, C_1=1

T_2=8, C_2=4

schedulable?

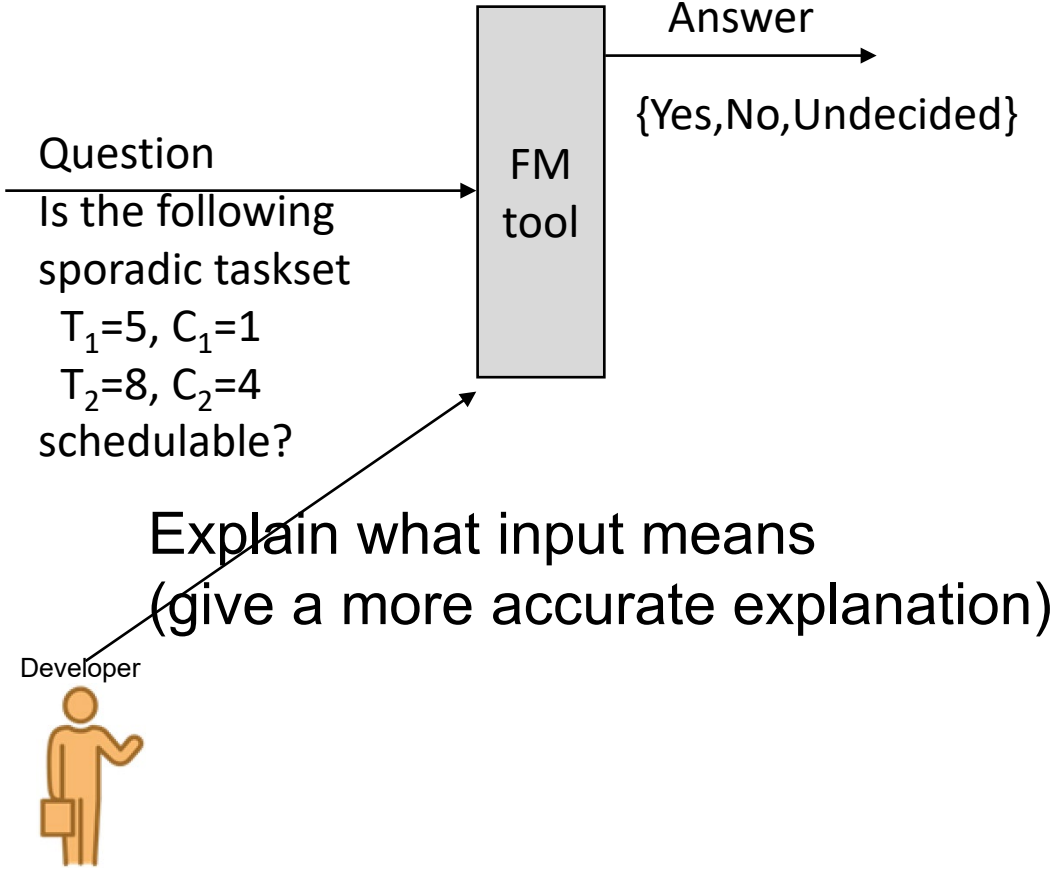
FM tool

Answer

{Yes,No,Undecided}

Developer

Explainability of Input: Real-Time Scheduling



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5
  }

Thread 2
  t=now()+8
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8
  }
```

Question

Is the following sporadic taskset

T_1=5, C_1=1

T_2=8, C_2=4

schedulable?

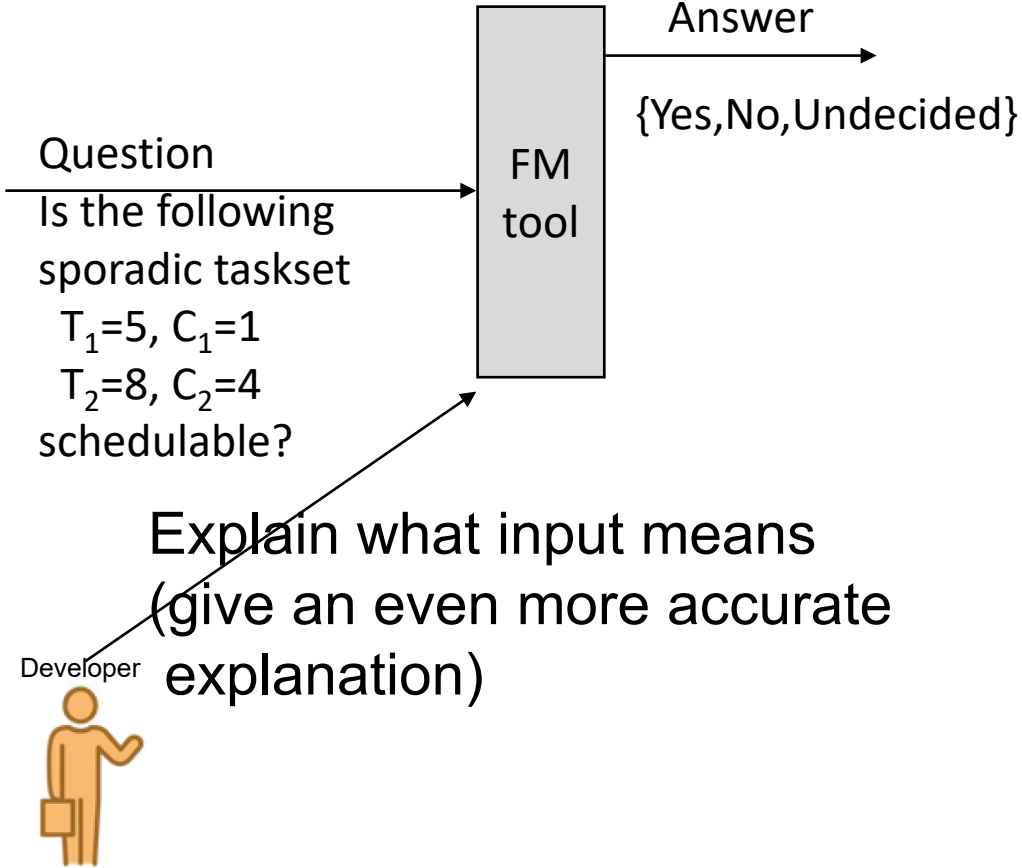
FM tool

Answer

{Yes,No,Undecided}

Developer

Explainability of Input: Real-Time Scheduling



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```

Question

Is the following sporadic taskset

T_1=5, C_1=1

T_2=8, C_2=4

schedulable?

FM tool

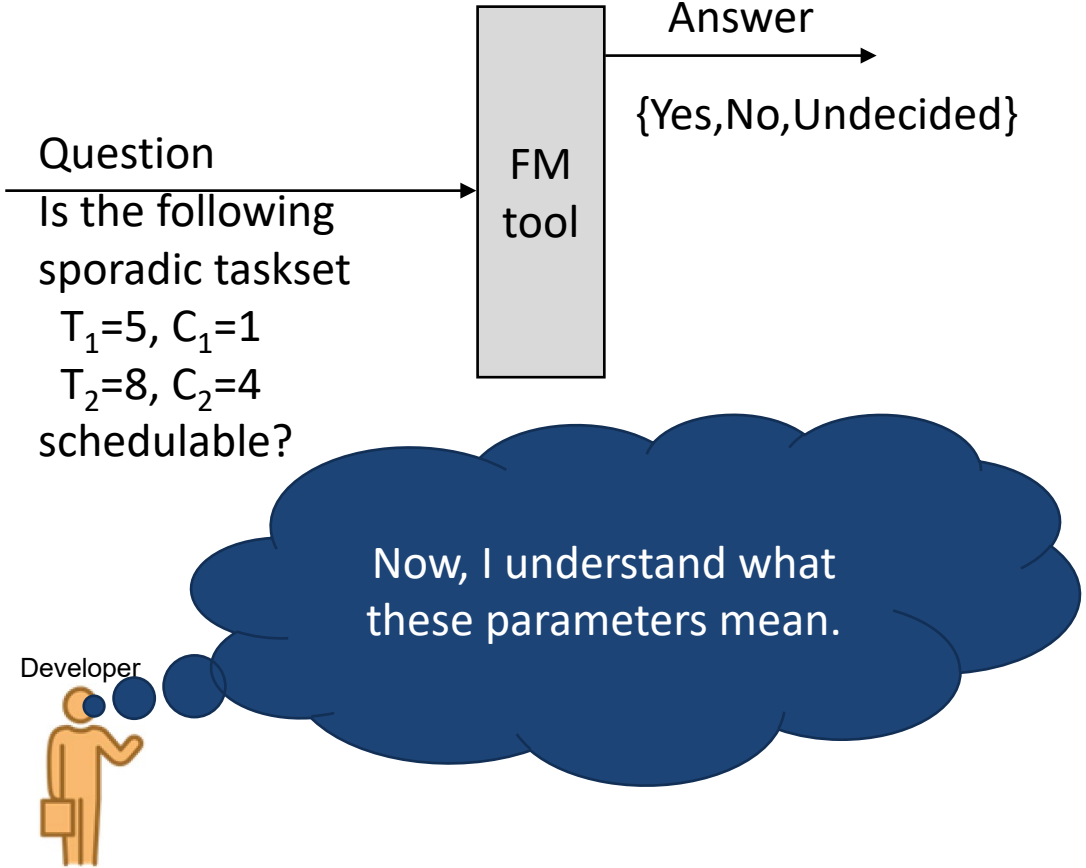
Answer

{Yes,No,Undecided}

Developer

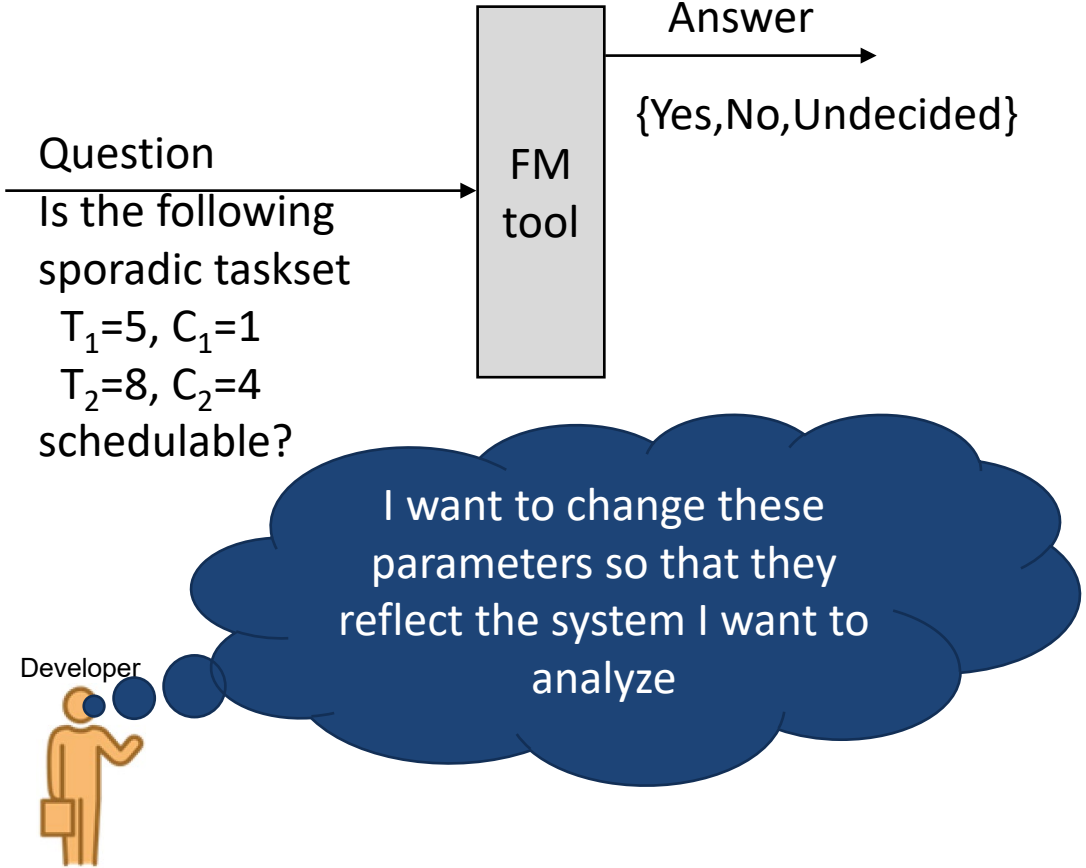
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```



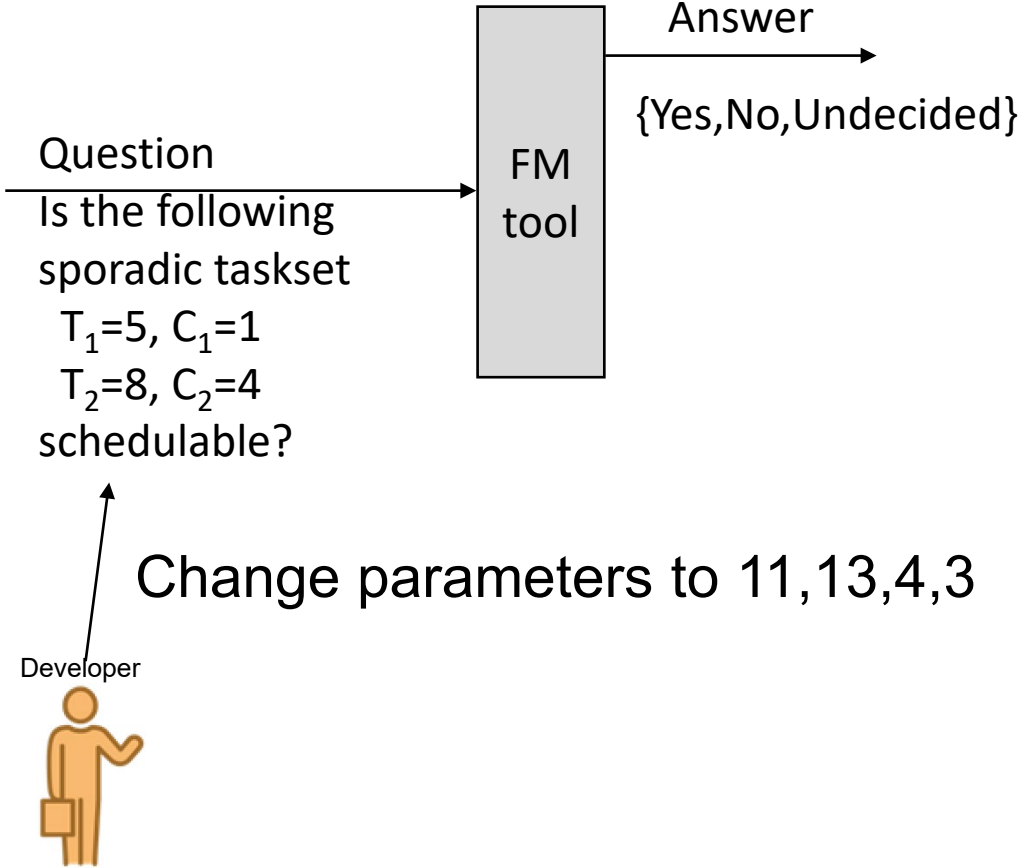
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```

Question

Is the following sporadic taskset

T_1=11, C_1=4

T_2=13, C_2=3

schedulable?

FM tool

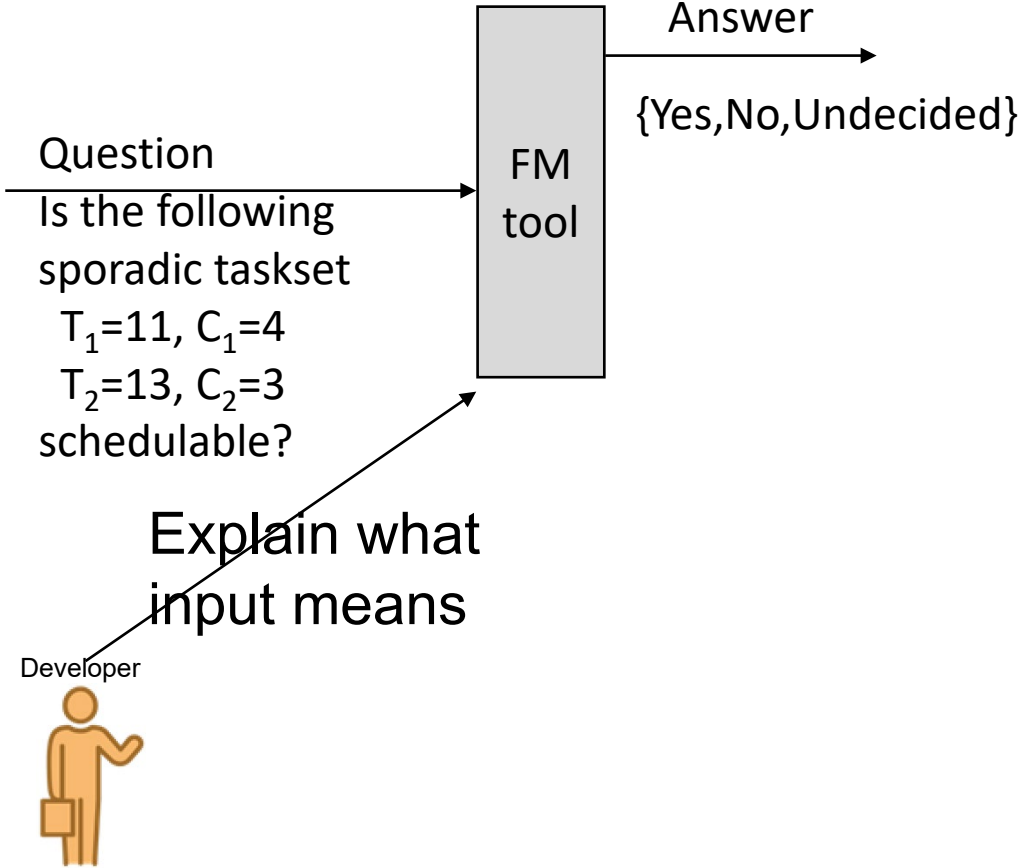
Answer

{Yes,No,Undecided}

Developer

Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+5+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,1))
    t=t+5+rand(0,inf)
  }

Thread 2
  t=now()+8+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+8+rand(0,inf)
  }
```



Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+11+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
  }

Thread 2
  t=now()+13+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
  }
```

Question

Is the following sporadic taskset

T_1=11, C_1=4

T_2=13, C_2=3

schedulable?

FM tool

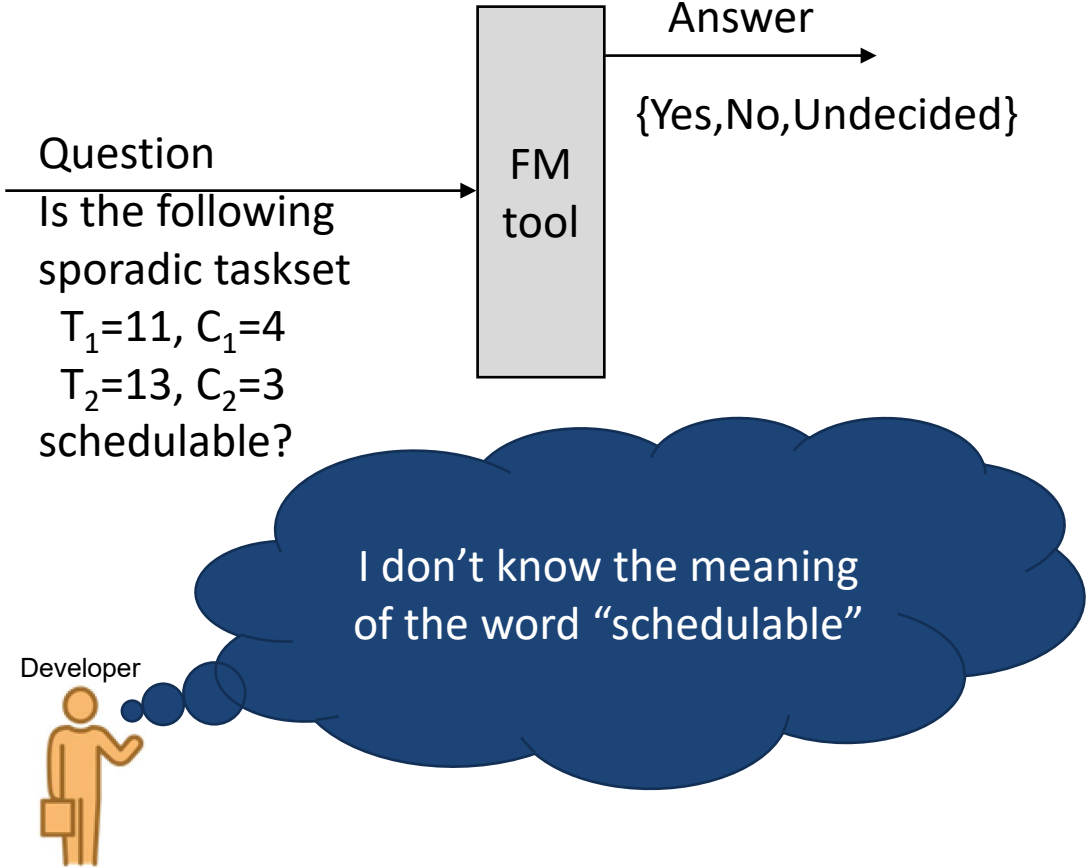
Answer

{Yes,No,Undecided}

Developer

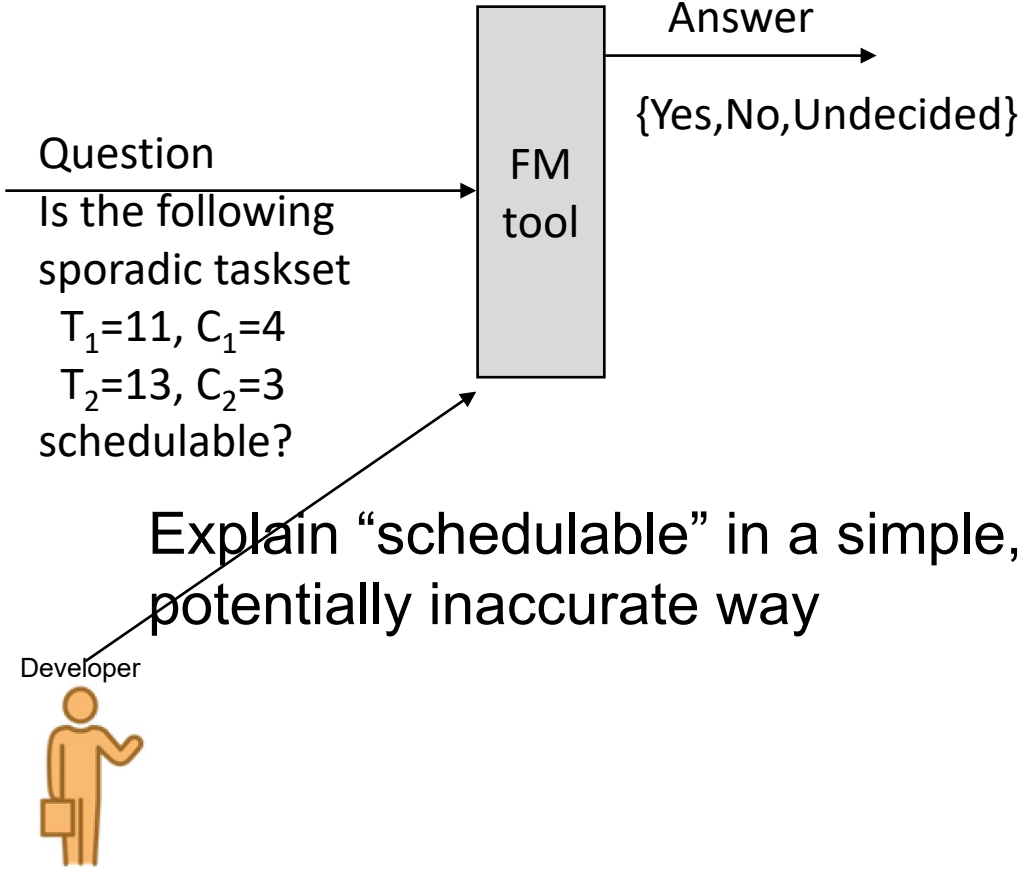
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+11+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
  }

Thread 2
  t=now()+13+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
  }
```



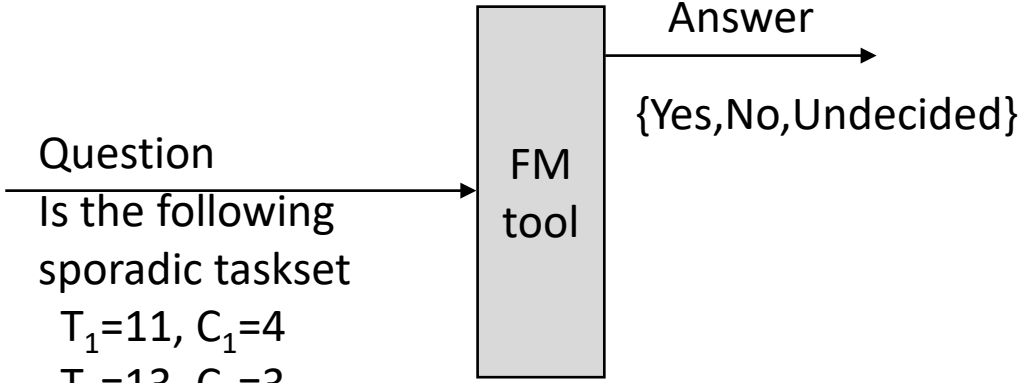
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+11+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
  }

Thread 2
  t=now()+13+rand(0,inf)
  while (1) {
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
  }
```



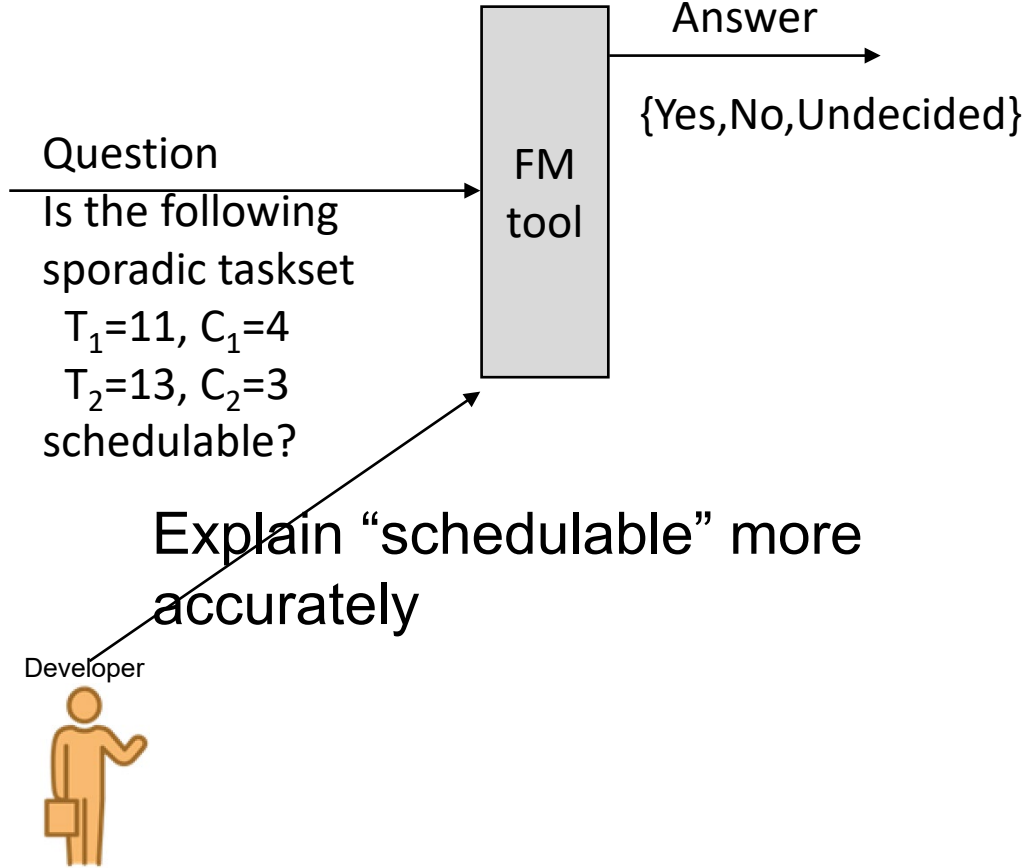
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+11+rand(0,inf)
  while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
  }

Thread 2
  t=now()+13+rand(0,inf)
  while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
  }
```



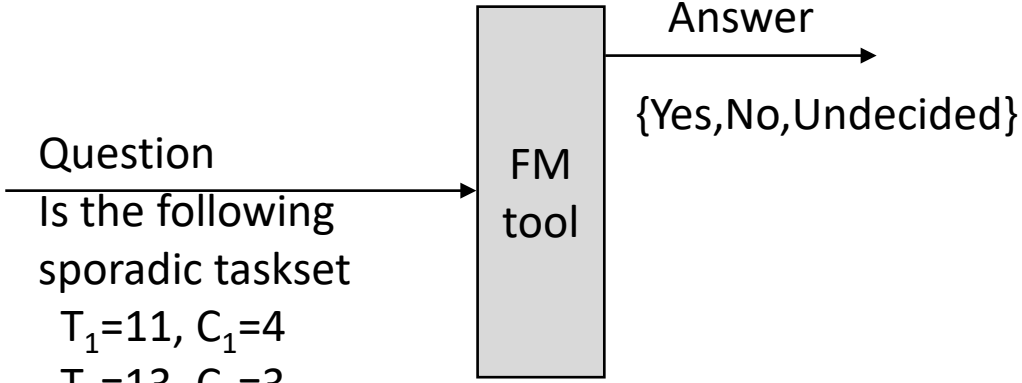
Explainability of Input: Real-Time Scheduling

```
Thread 1
  t=now()+11+rand(0,inf)
  while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,4))
    t=t+11+rand(0,inf)
  }

Thread 2
  t=now()+13+rand(0,inf)
  while (1) {
    assert(now()<=t)
    delay_until(t)
    do_work(rand(0,3))
    t=t+13+rand(0,inf)
  }
```



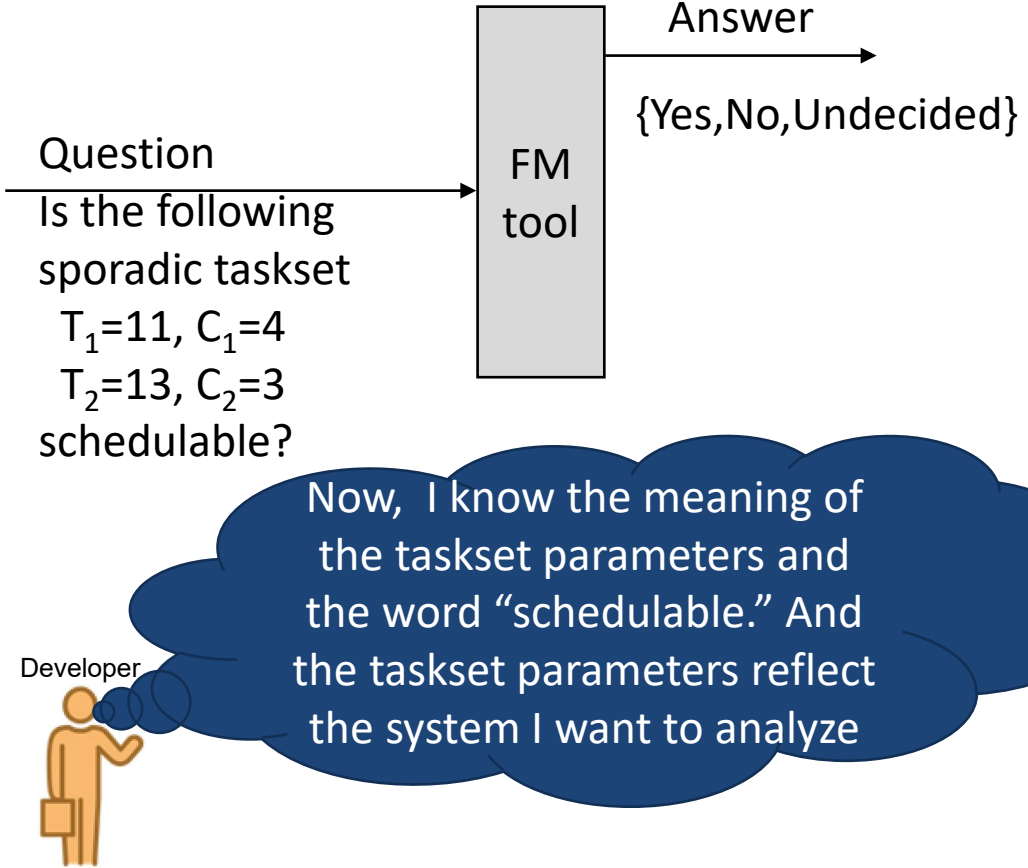
Explainability of Input: Real-Time Scheduling

```
Thread 1
  gap=rand(0,inf)
  t=now()+11+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
  }

Thread 2
  gap=rand(0,inf)
  t=now()+13+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
  }
```



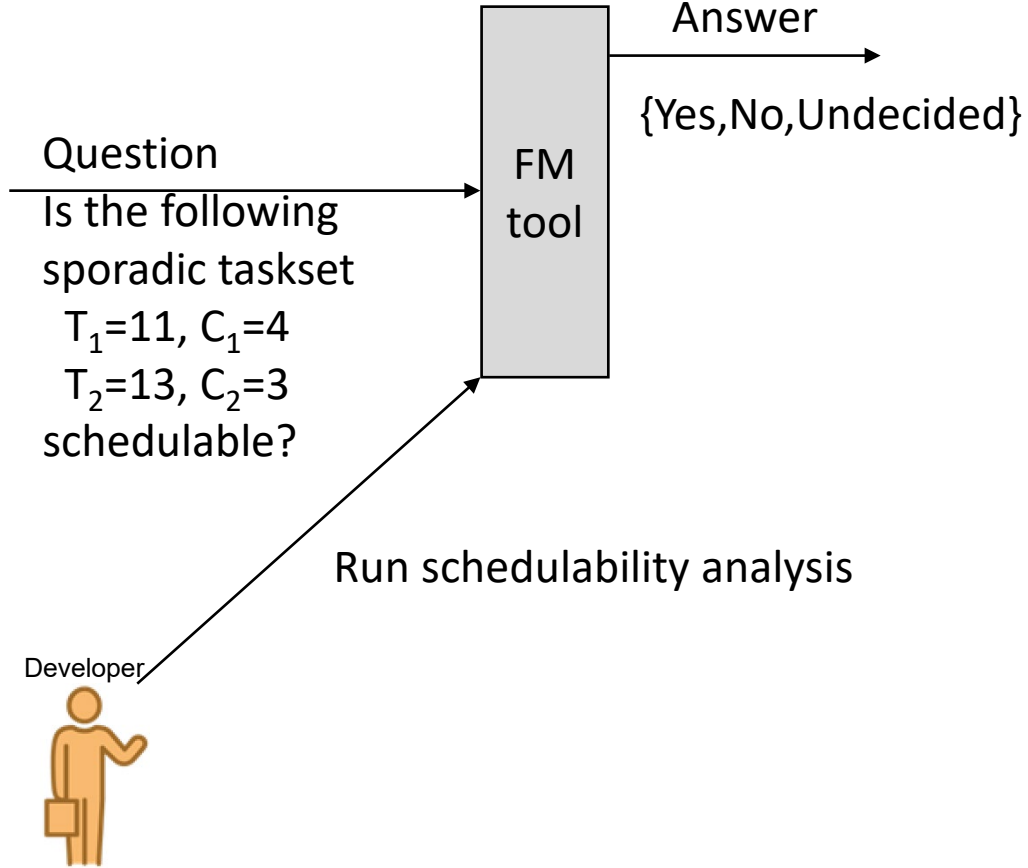
Explainability of Input: Real-Time Scheduling

```
Thread 1
  gap=rand(0,inf)
  t=now()+11+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
  }

Thread 2
  gap=rand(0,inf)
  t=now()+13+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
  }
```



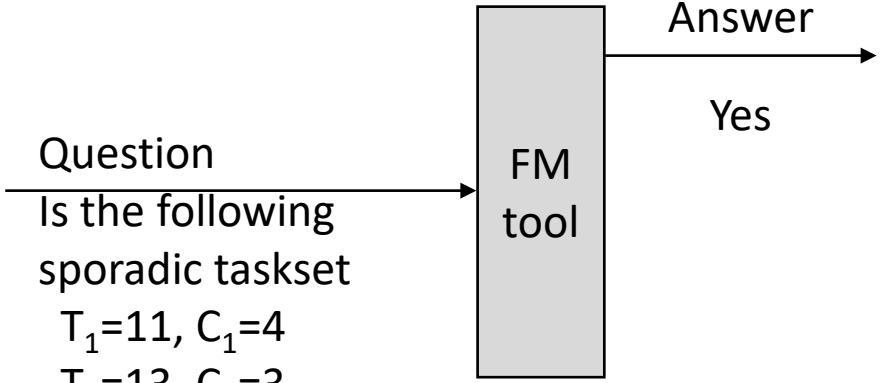
Explainability of Input: Real-Time Scheduling

```
Thread 1
  gap=rand(0,inf)
  t=now()+11+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
  }

Thread 2
  gap=rand(0,inf)
  t=now()+13+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
  }
```



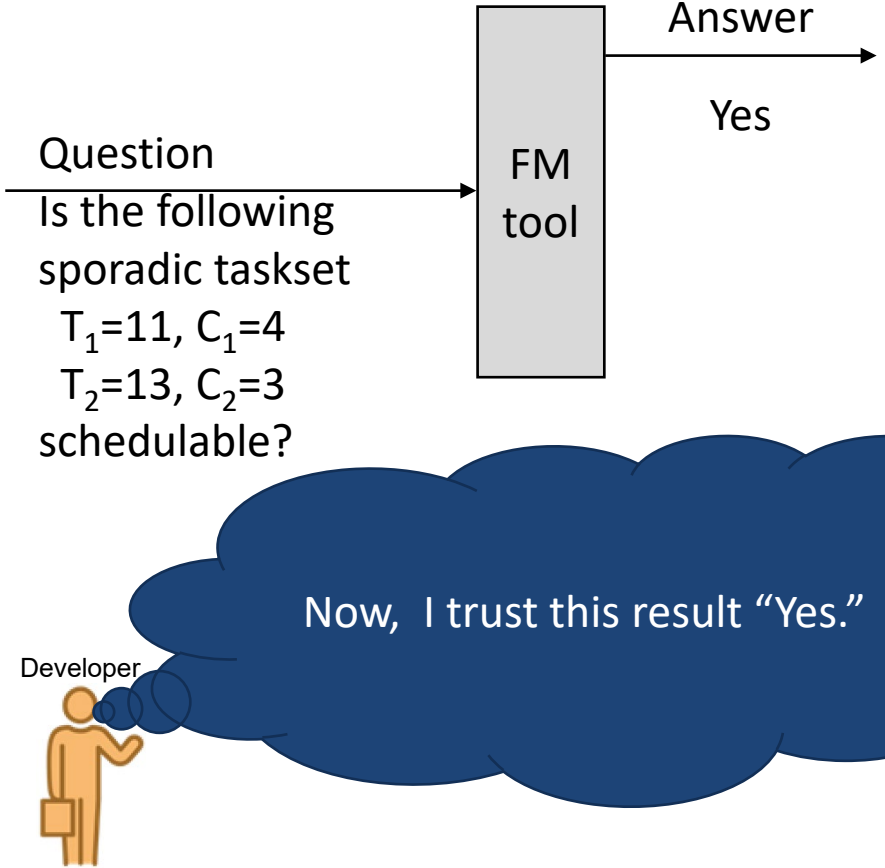
Explainability of Input: Real-Time Scheduling

```
Thread 1
  gap=rand(0,inf)
  t=now()+11+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
  }

Thread 2
  gap=rand(0,inf)
  t=now()+13+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
  }
```

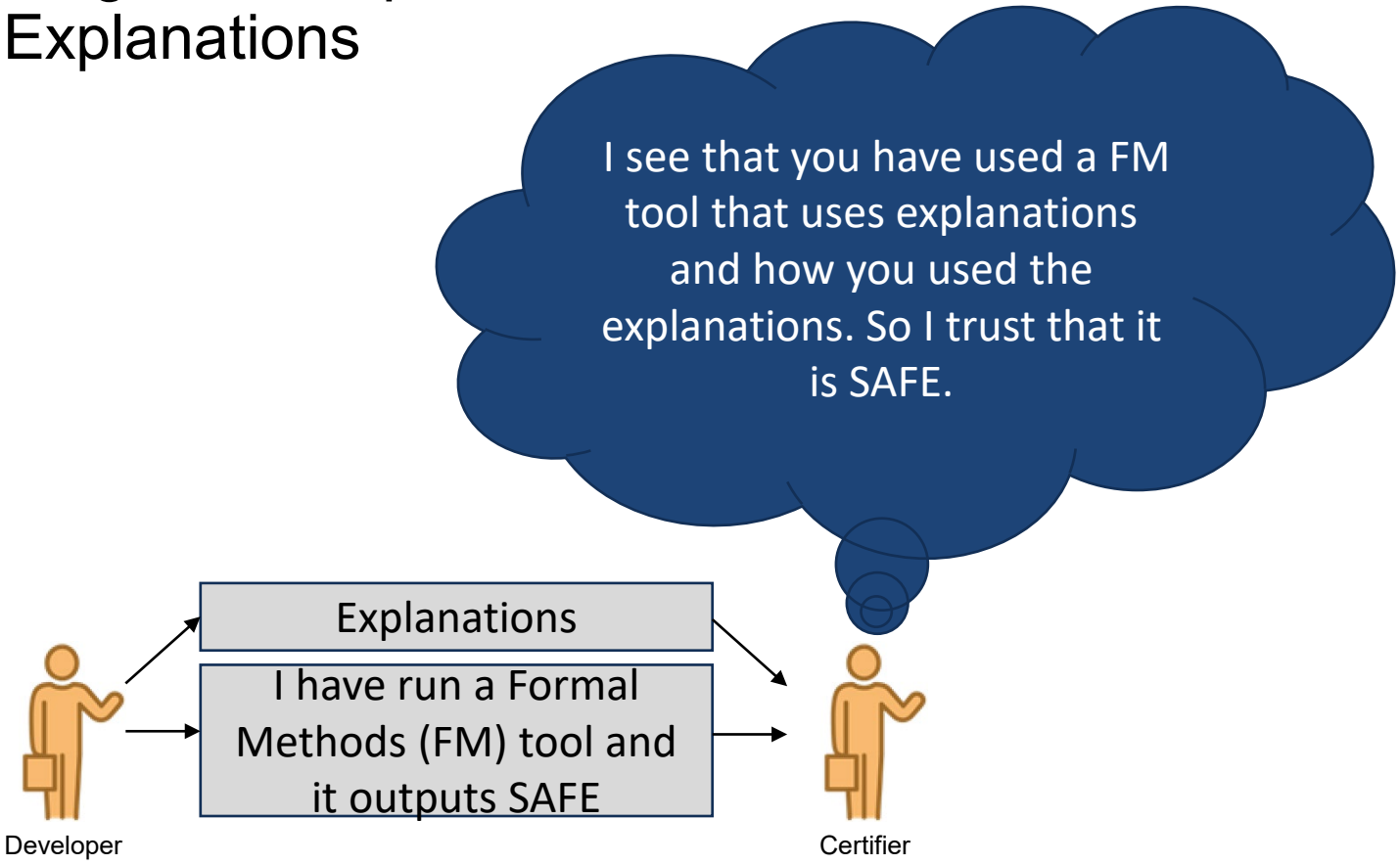


Explainability of Input: Real-Time Scheduling

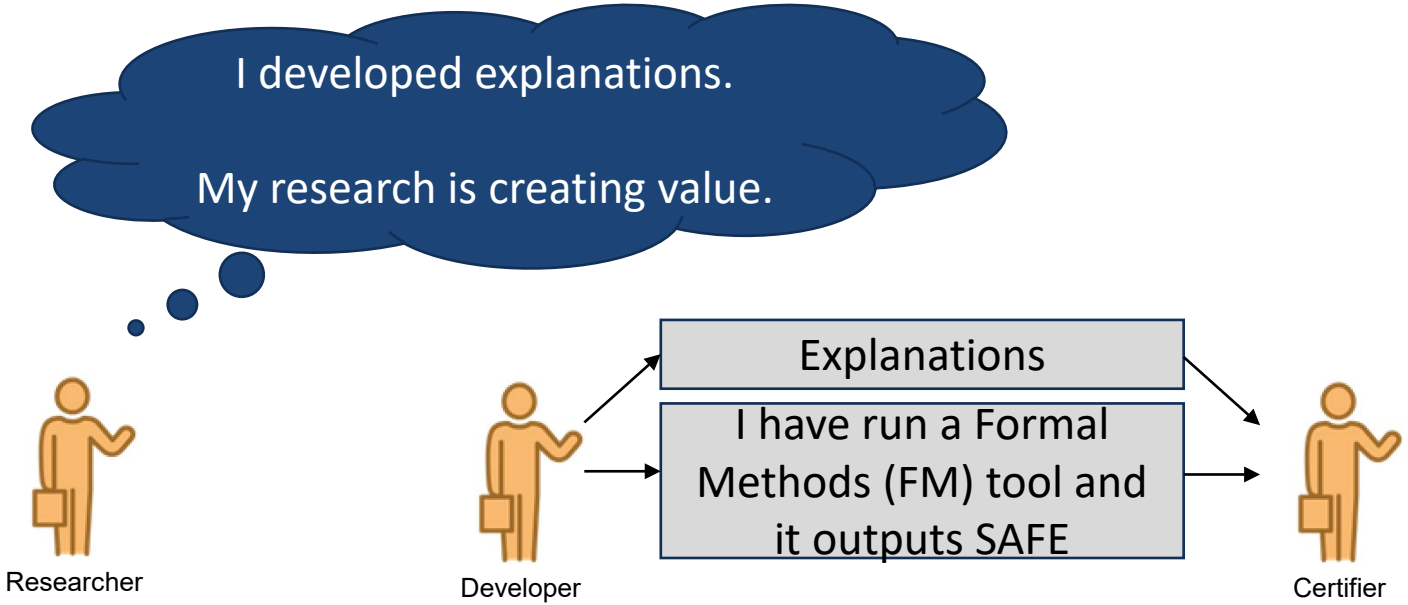
```
Thread 1
  gap=rand(0,inf)
  t=now()+11+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,4))
    gap=rand(0,inf)
    t=t+11+gap
  }

Thread 2
  gap=rand(0,inf)
  t=now()+13+gap
  while (1) {
    assert(now()<=t-gap)
    delay_until(t)
    do_work(rand(0,3))
    gap=rand(0,inf)
    t=t+13+gap
  }
```



Problem: How to get developers and certifiers to trust FMs
Solution: Use Explanations



Problem: How to get developers and certifiers to trust FMs
Solution: Use Explanations



Conclusion

Explainability of real-time systems and their analysis:

has the potential to make research results valuable to practitioners

is a new, rich, area of research