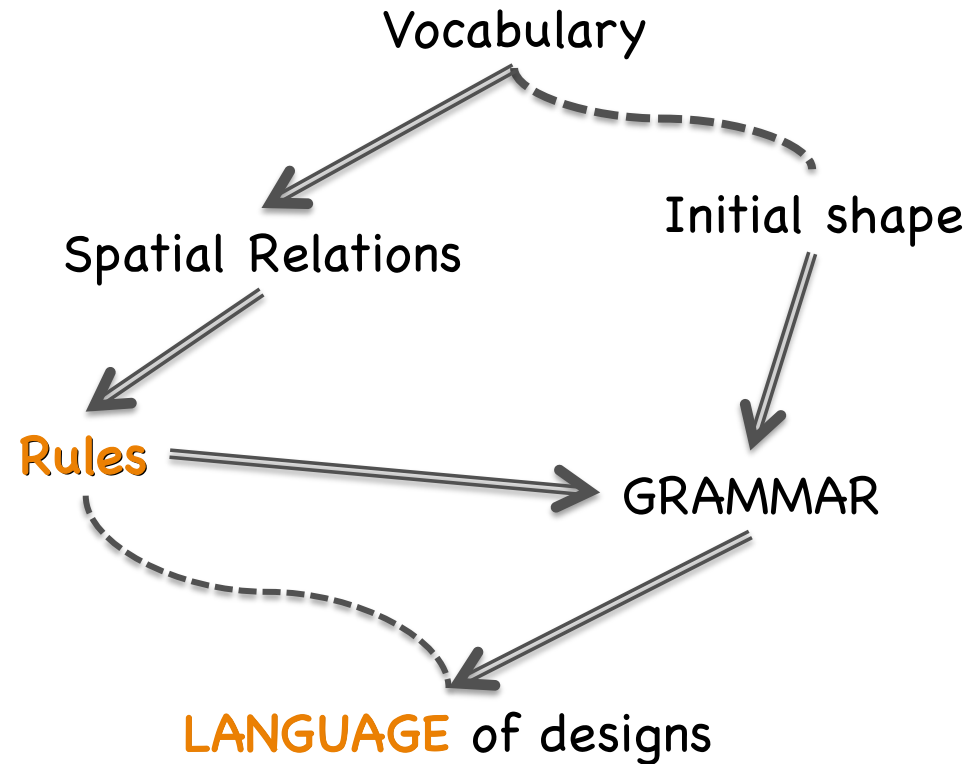
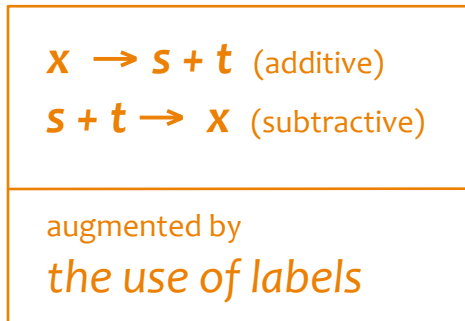


48-747 Shape Grammars

Forming New Languages from Old

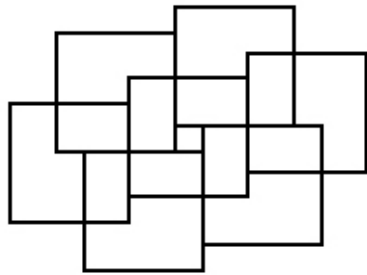
Spatial Metathesis



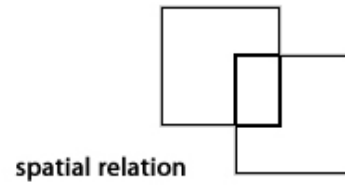
Gothic
Classical
Islamic
Eastern styles

*Encapsulates “theme”
Explored for design “styles”
Perhaps motivated by revival
or restoration*

RECAP grammar paradigm



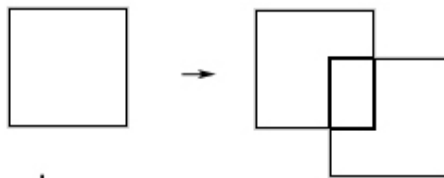
design



spatial relation



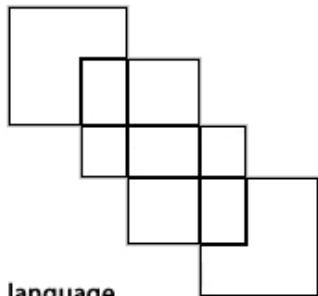
vocabulary



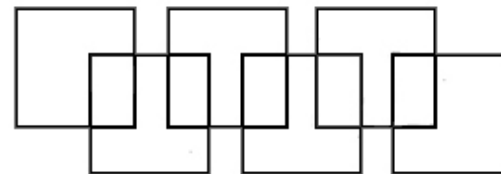
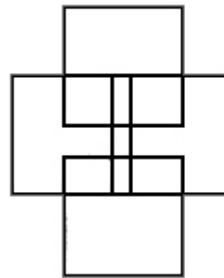
rule



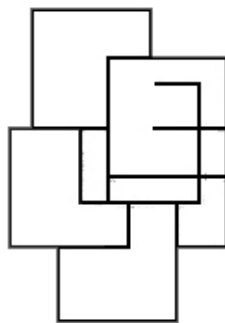
initial shape



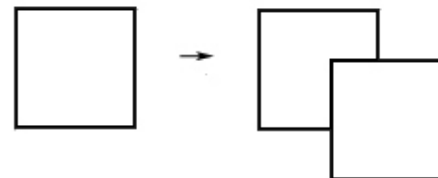
language



Languages give birth to new languages



design



new rule

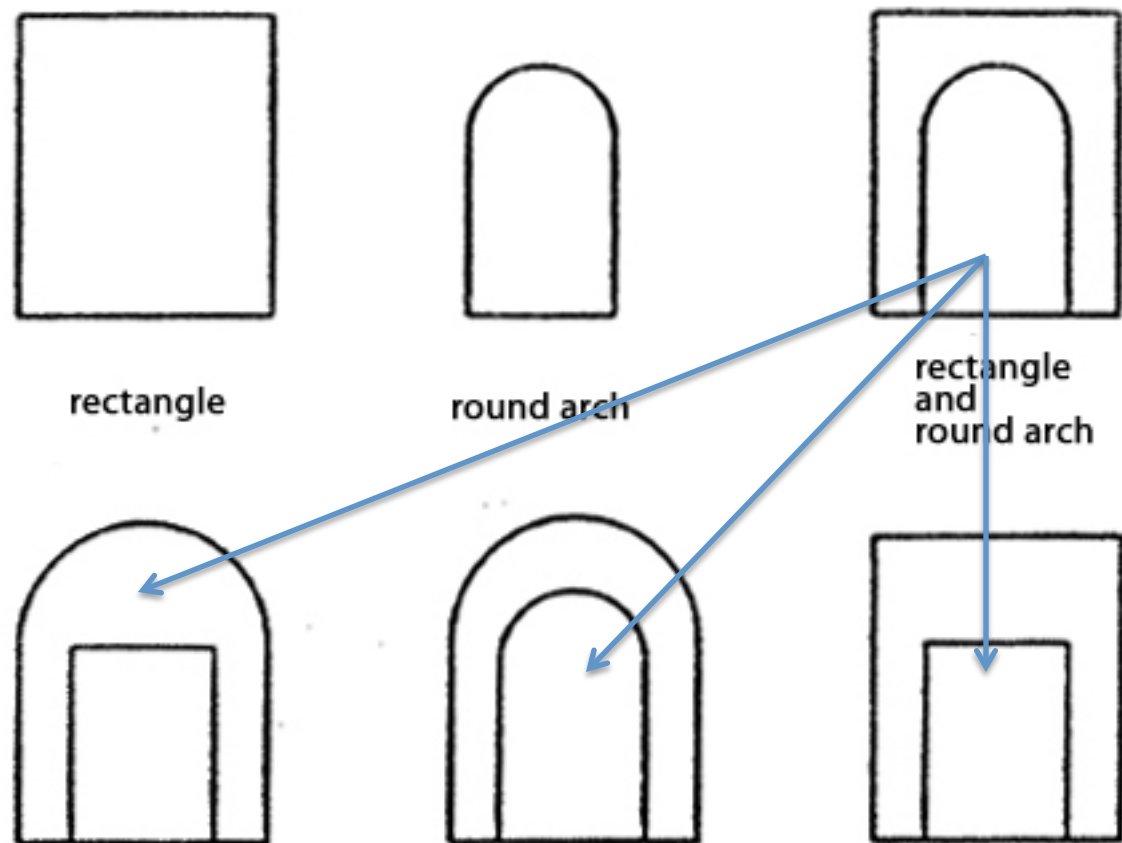
transposition of letters, words, sounds, syllables

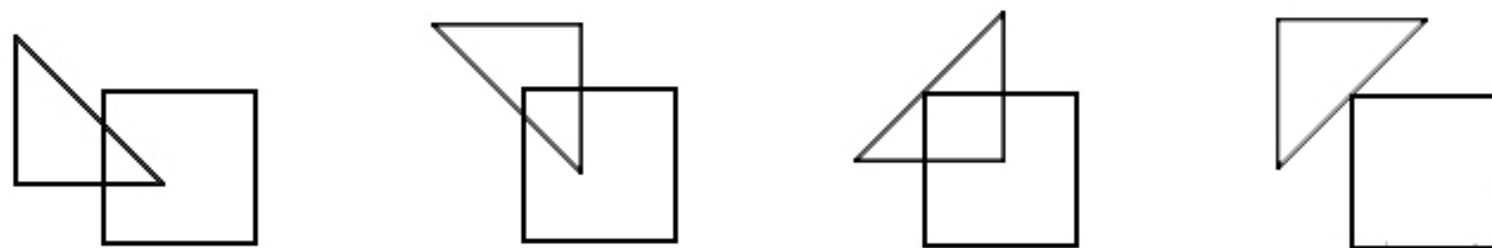
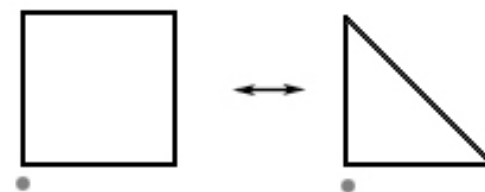
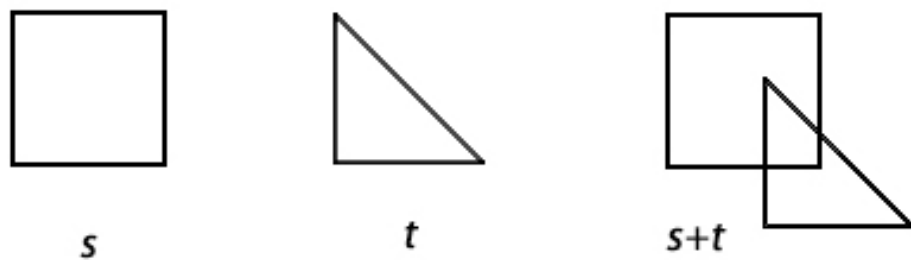
bird ← brid

evelate ← elevate (spoonerism)

why not transposition of shapes or images by shape replacement ?

metathesis





possible spatial metathesis

metathesis

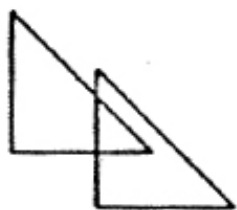
Is of the form $\mathbf{a} \leftrightarrow \mathbf{b}$ where neither \mathbf{a} nor \mathbf{b} is empty

Apply the rule to a spatial relation R , a set of shapes, to produce a new spatial relation N provided R contains a shape s and there is a geometrical transformation f such that either $s = f(\mathbf{a})$ or $s = f(\mathbf{b})$

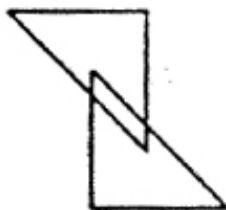
$$N = S - f(\mathbf{a}) + f(\mathbf{b}) \text{ if } s = f(\mathbf{a})$$

$$N = S - f(\mathbf{b}) + f(\mathbf{a}) \text{ if } s = f(\mathbf{b})$$

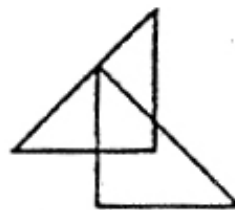
shape equivalence rule



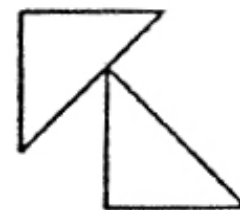
1



2

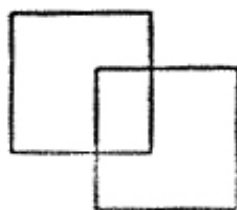


3



4

(→)



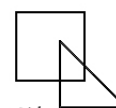
(←)



s



t



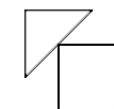
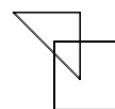
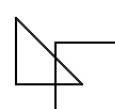
s+t



→



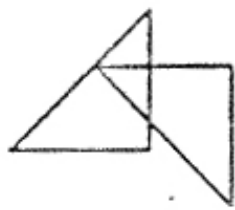
t



possible spatial metathesis



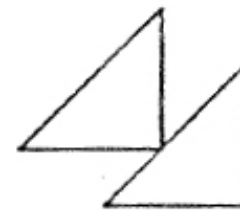
1



2



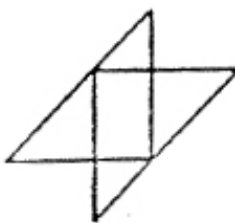
3



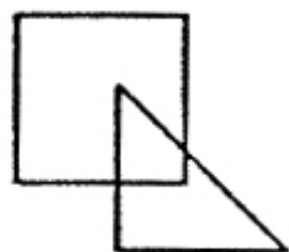
4



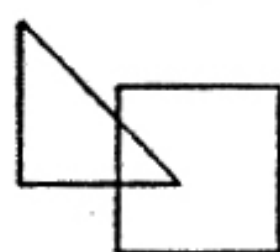
5



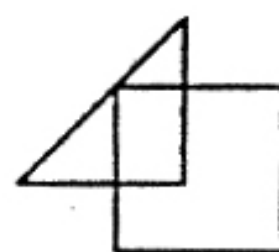
6



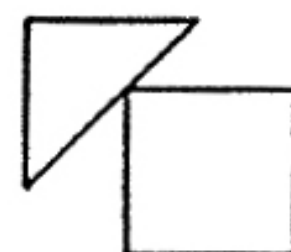
1



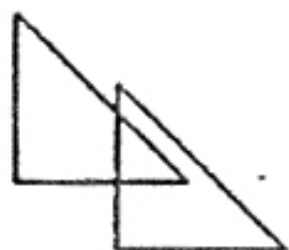
2



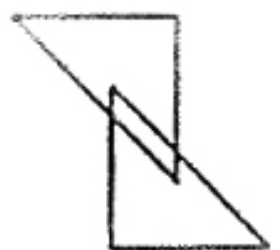
3



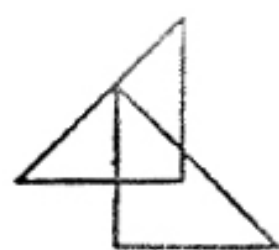
4



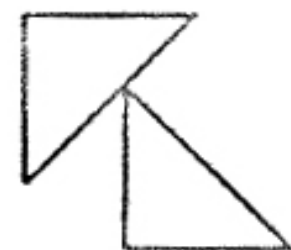
5



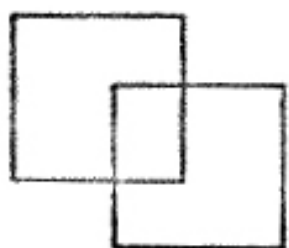
6



7



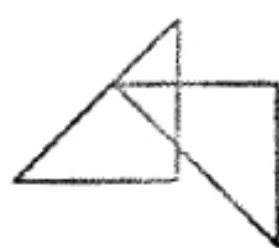
8



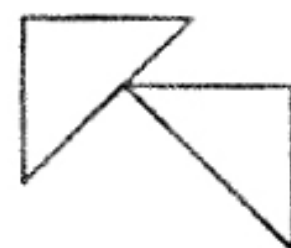
9



10



11



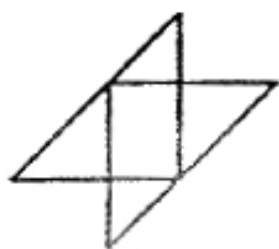
12



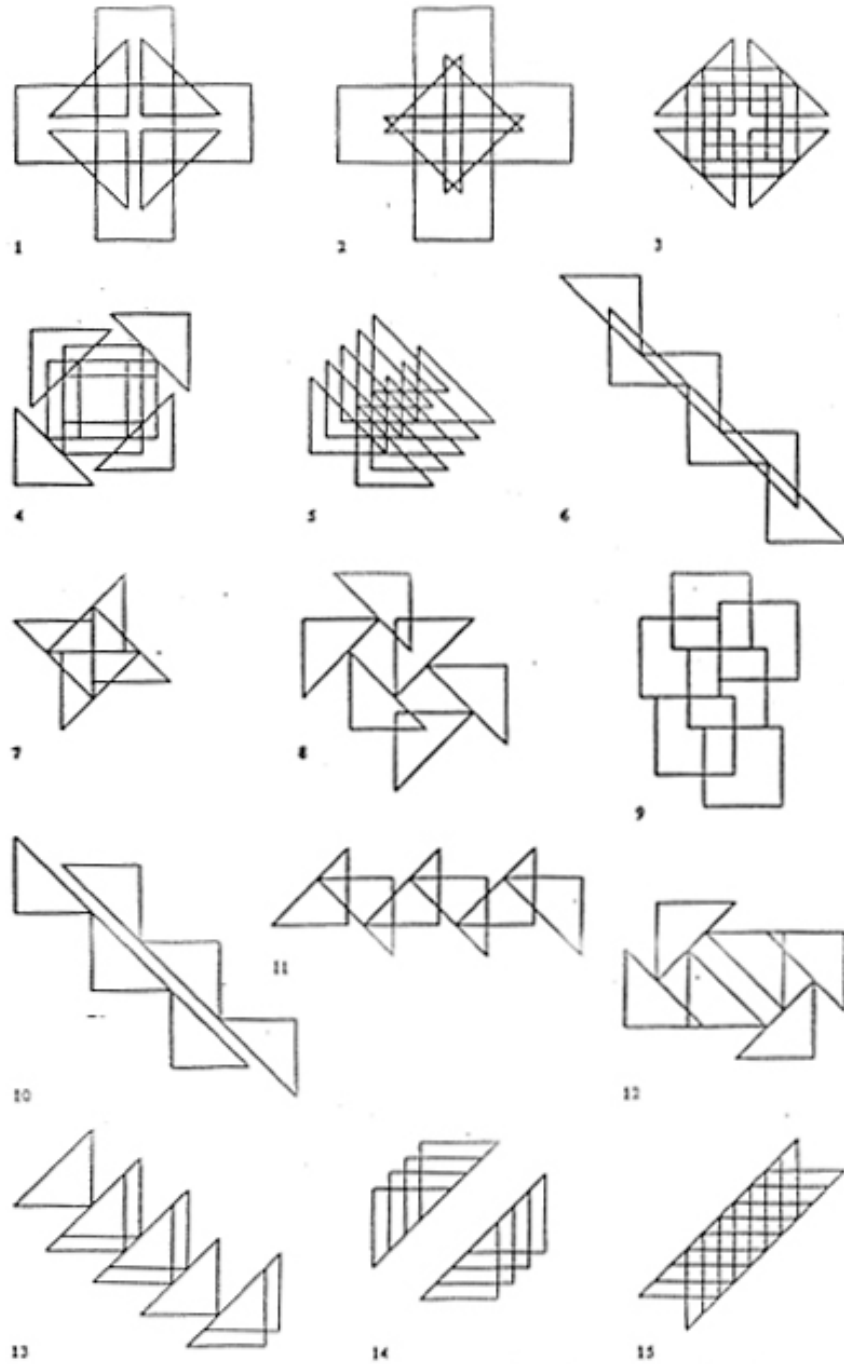
13



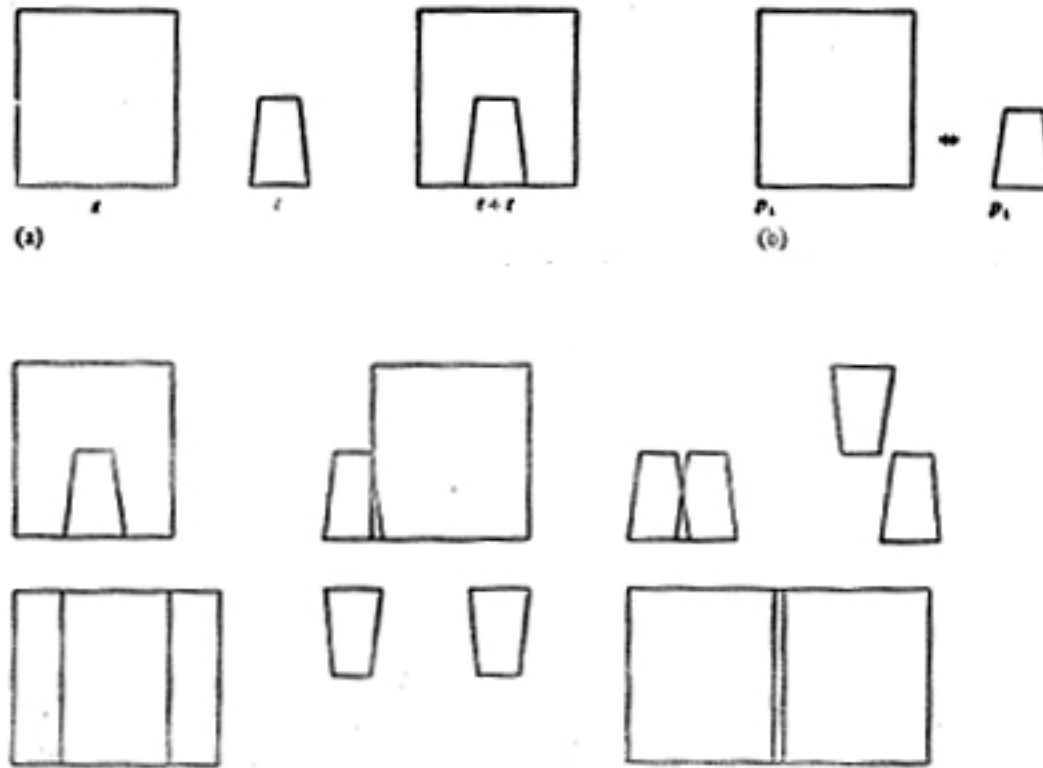
14



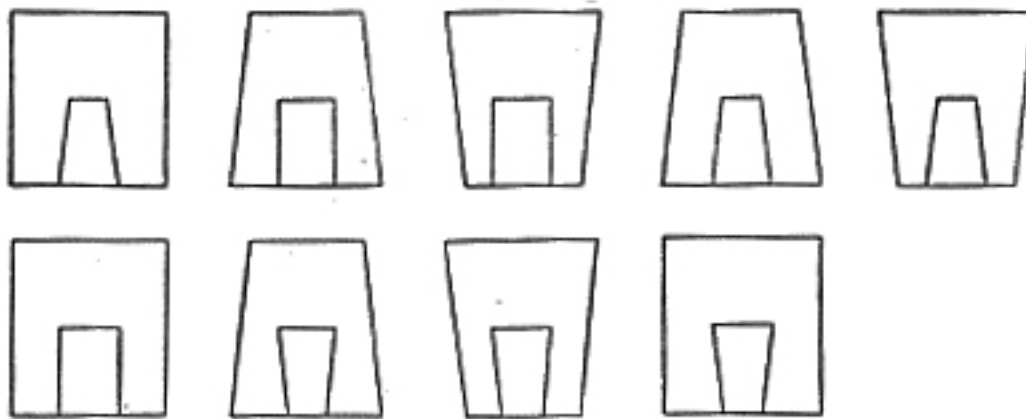
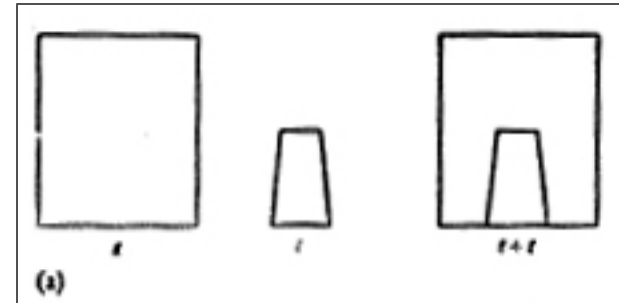
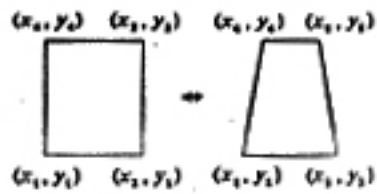
15



example designs



if we can have shape equivalence rules why not **shape equivalence schemas?**



shape equivalence schema

Is a schema of the form $\mathbf{a} \leftrightarrow \mathbf{b}$ where neither \mathbf{a} nor \mathbf{b} is empty, \mathbf{a} and \mathbf{b} have open terms

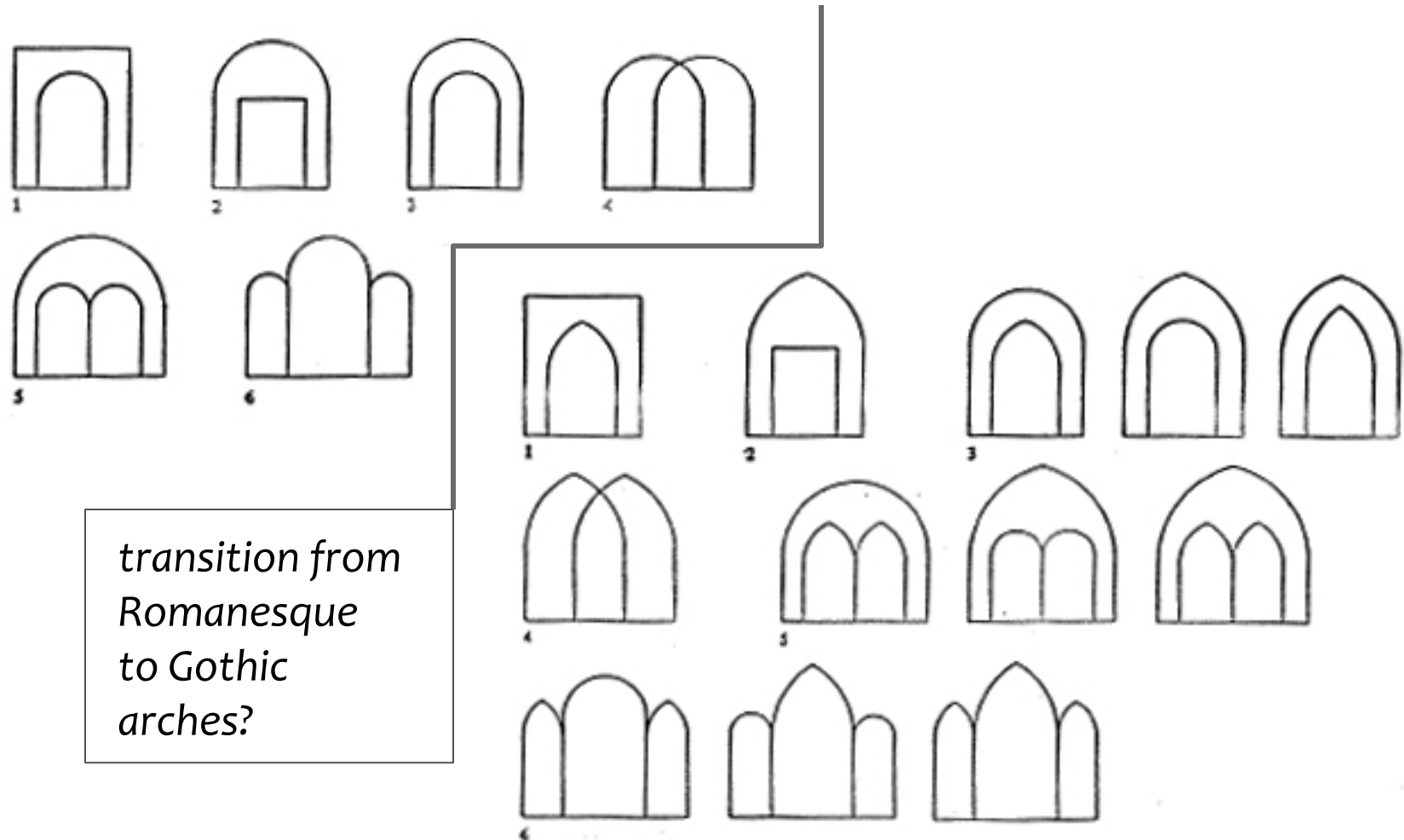
Apply the schema to a spatial relation R , a set of shapes, to produce a new spatial relation N provided R contains a shape s , there is an assignment g to all open variables in \mathbf{a} and \mathbf{b} , and there is a geometrical transformation f such that either $s = f(\mathbf{a})$ or $s = f(\mathbf{b})$

$$N = S - f(g[\mathbf{a}]) + f(\mathbf{b}) \text{ if } s = f(g[\mathbf{a}])$$

$$N = S - f(\mathbf{b}) + f(\mathbf{a}) \text{ if } s = f(g[\mathbf{b}])$$

shape equivalence schema

What we have seen so far is a FLIP-FLOP between shapes/schemas with the implicit PROVISIO that no new shapes are introduced into the relation



what about **introducing new shapes** into the **equivalence rule**



1



2



3



4



5



6



7



8



9



10



11



12



13



14



15



16



17



18



19



20



21



22



23



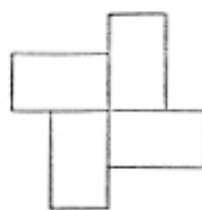
24



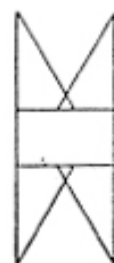
25



1



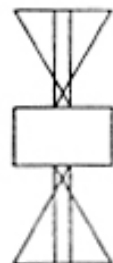
2



3



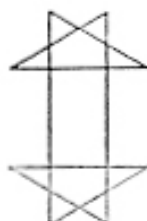
4



5



6



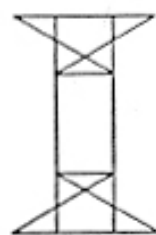
7



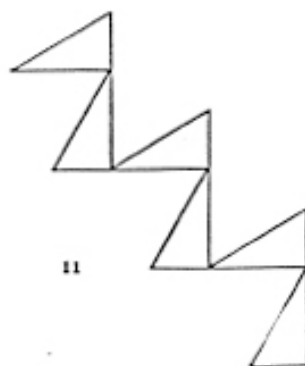
8



9



10



11



12



13



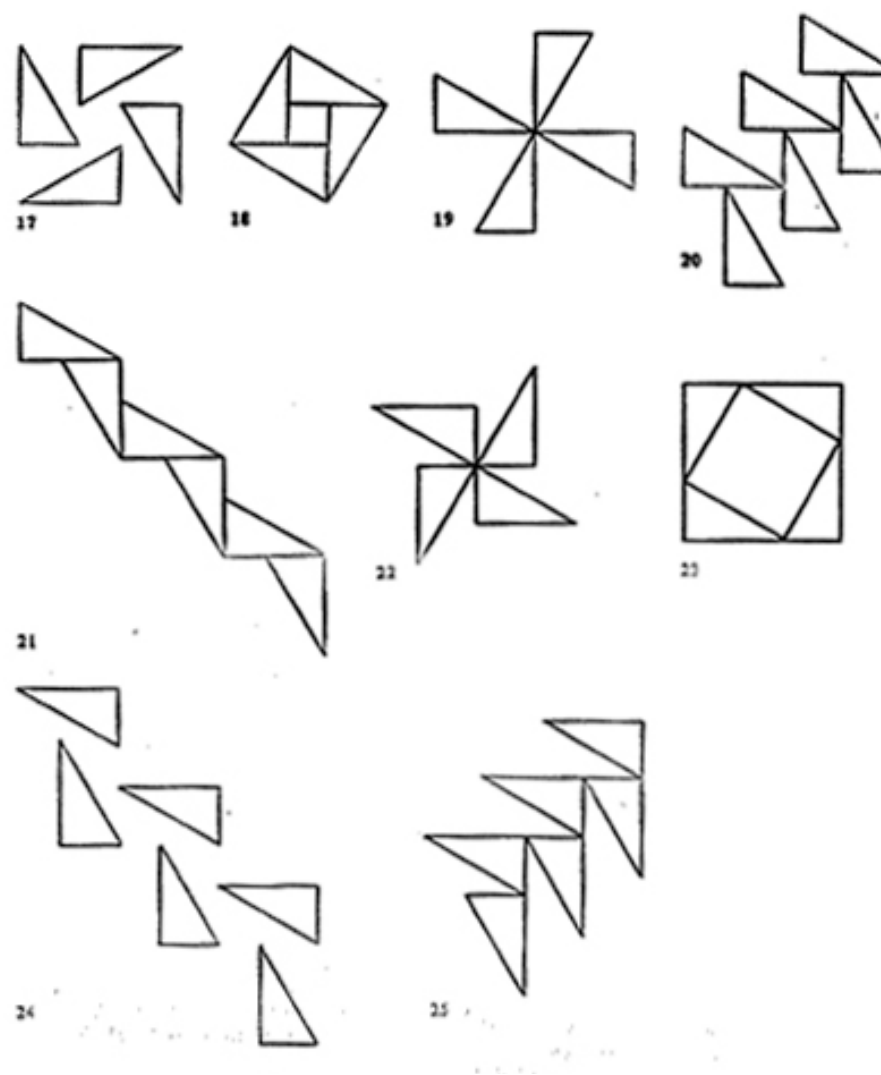
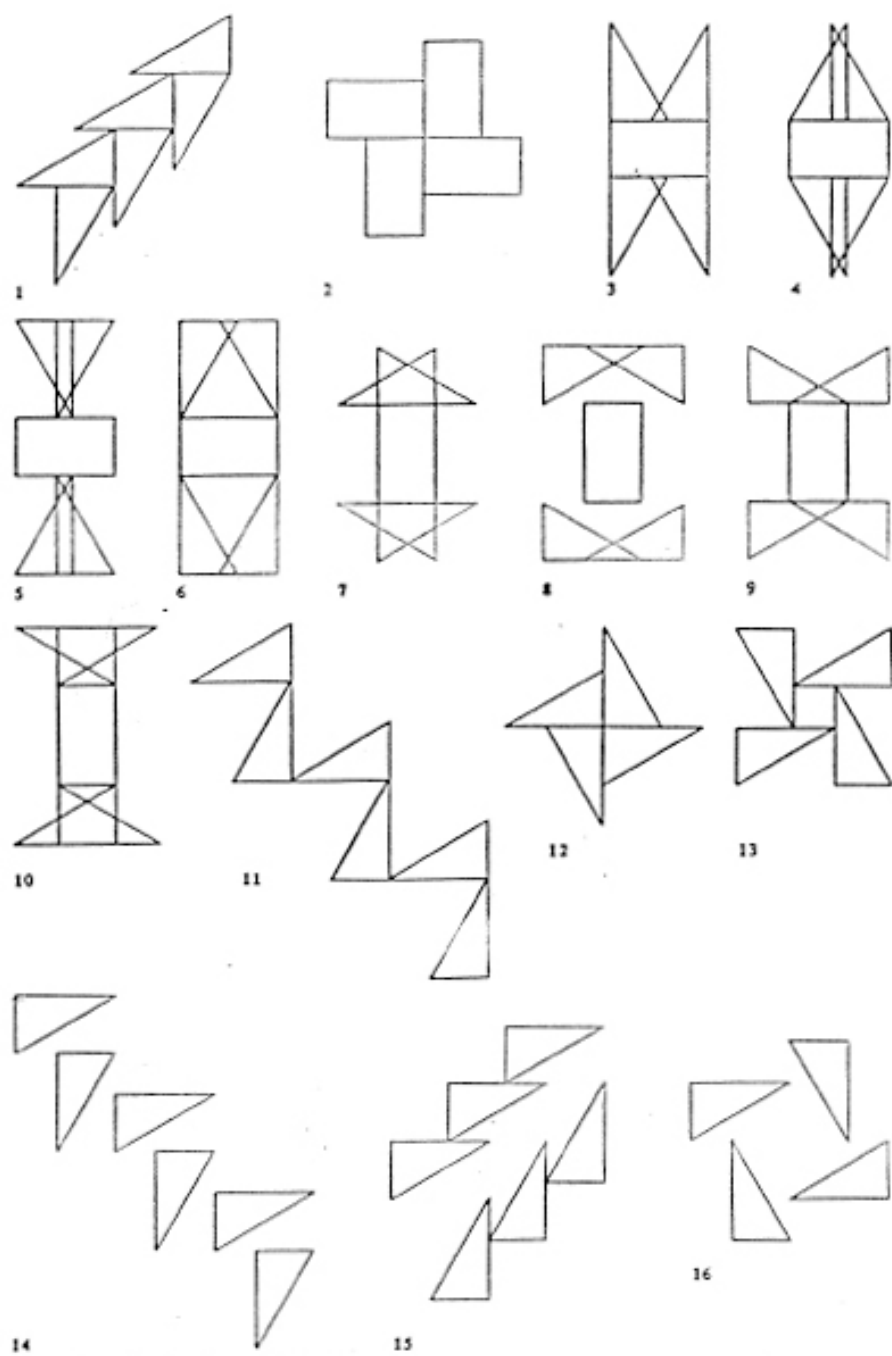
14



15



16



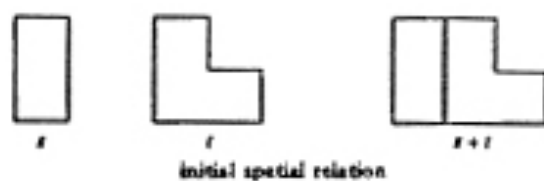
We have $\mathbf{a} \leftrightarrow \mathbf{b}$

We can construct classes of spatial relations by looking at $f[h(\mathbf{a})]$ and $g[j(\mathbf{b})]$ so that

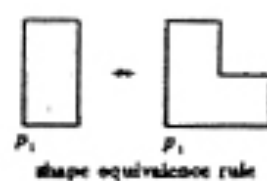
$$N = S - f(\mathbf{a}) + g(\mathbf{b})$$

$$N = S - f(h(\mathbf{b})) + g(j(\mathbf{a}))$$

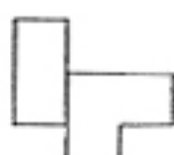
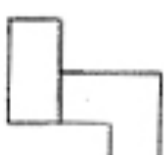
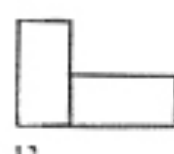
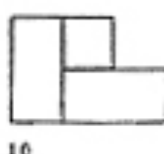
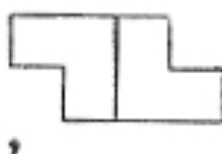
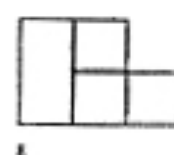
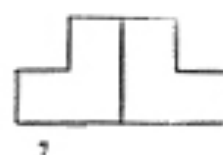
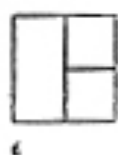
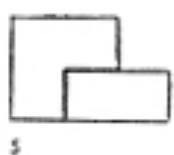
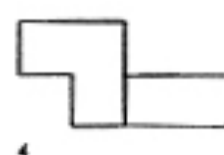
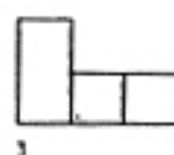
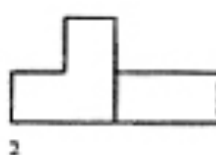
any more variations?

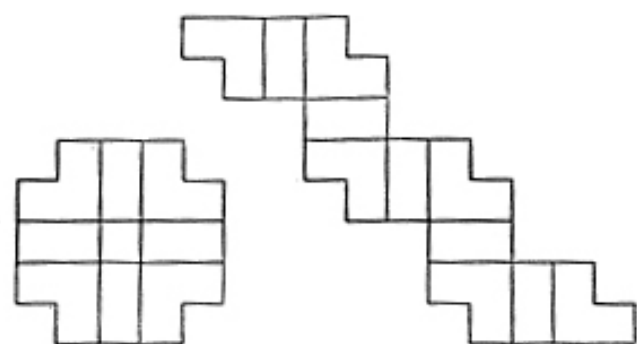


initial spatial relation

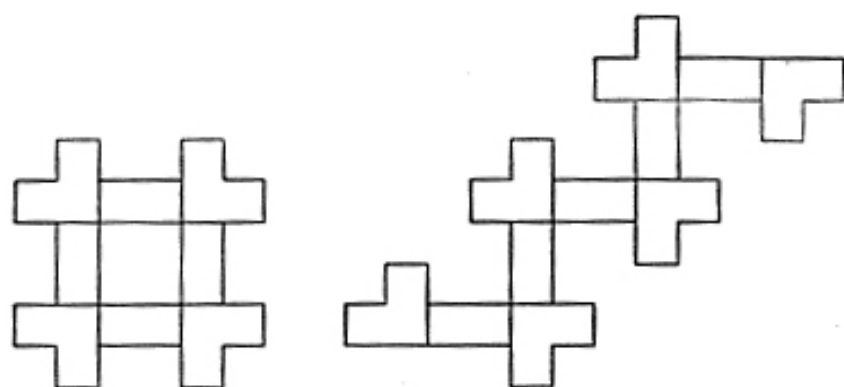


shape equivalence rule





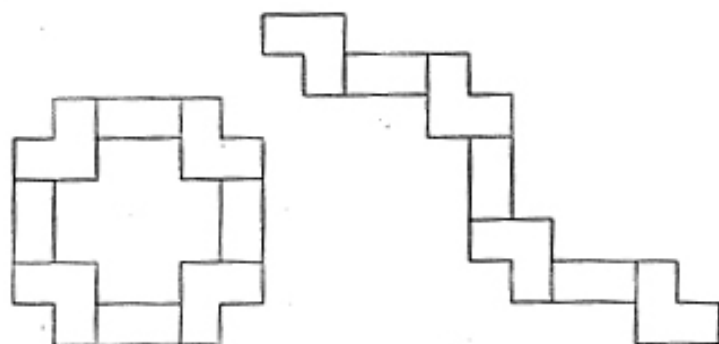
1



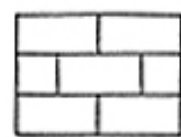
2



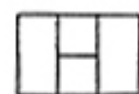
3



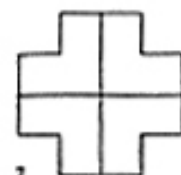
4



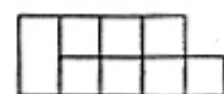
5



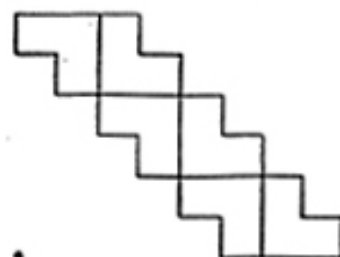
6



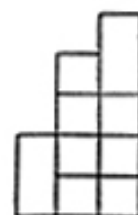
7



8



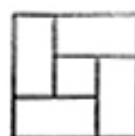
9



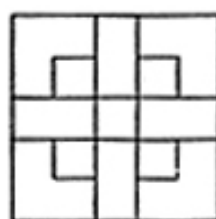
10



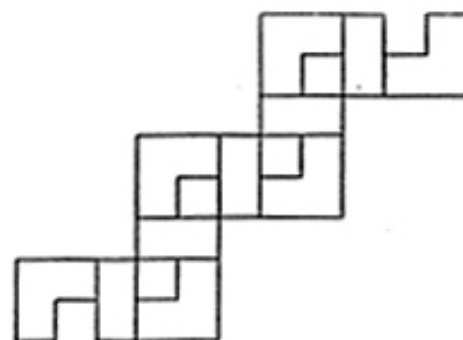
11



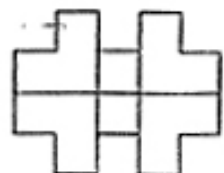
12



13



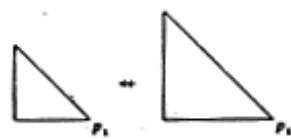
14



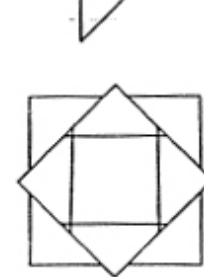
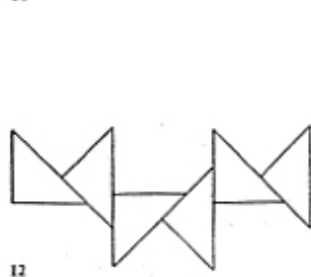
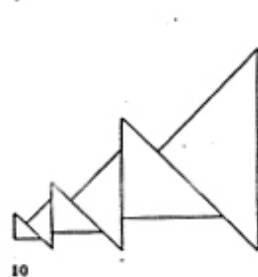
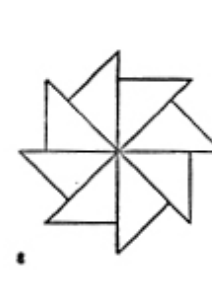
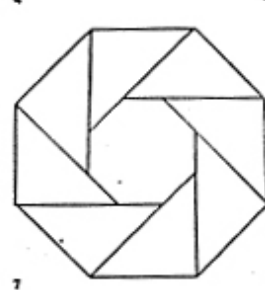
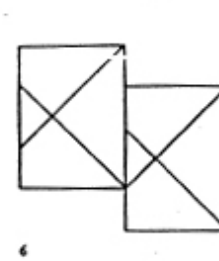
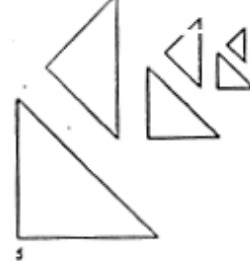
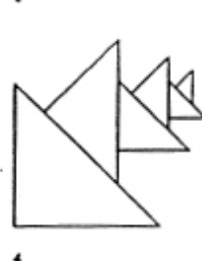
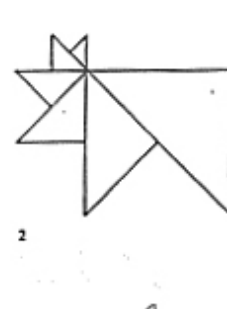
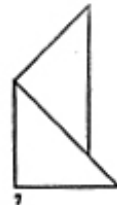
15



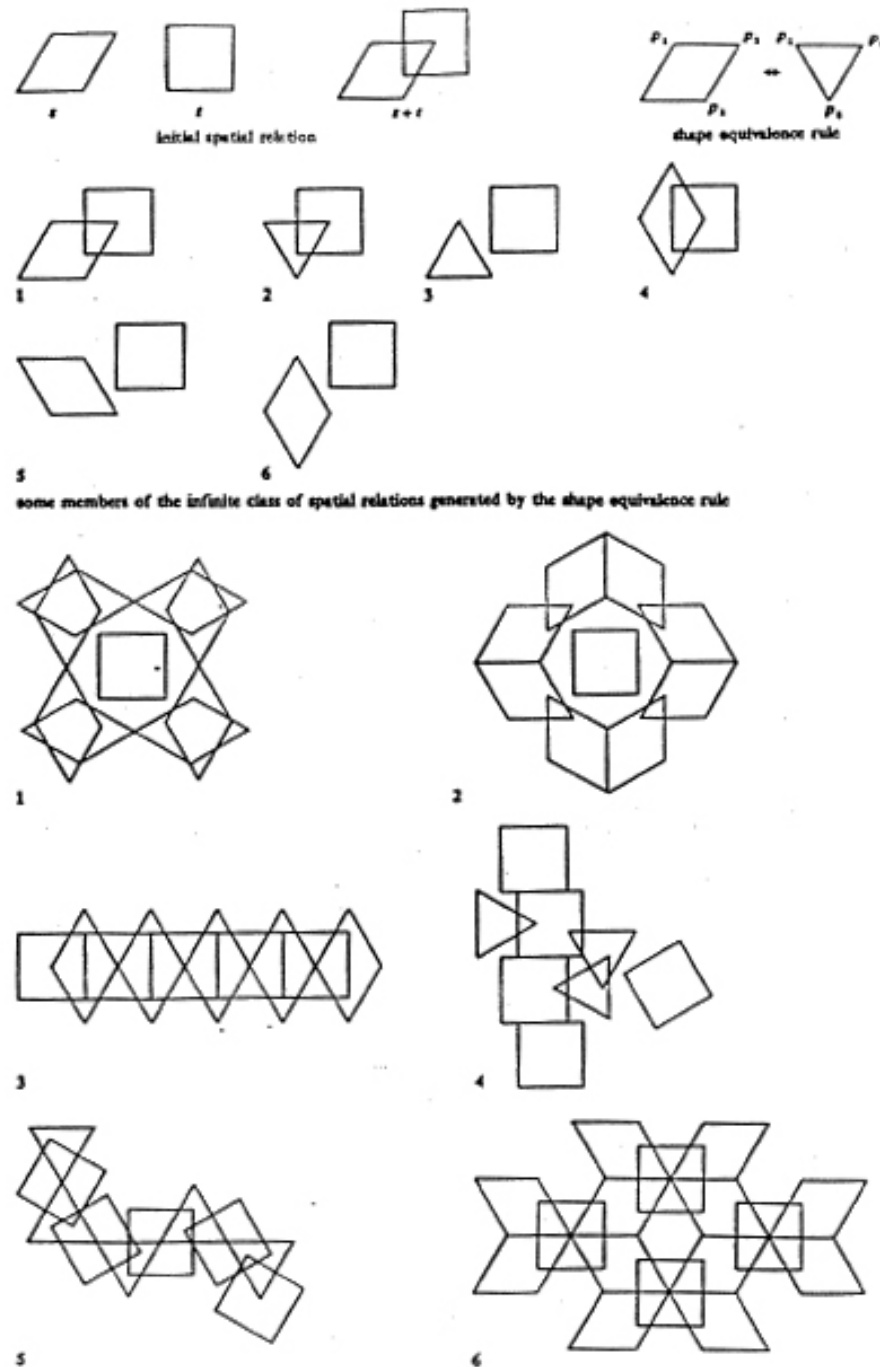
Initial special relation



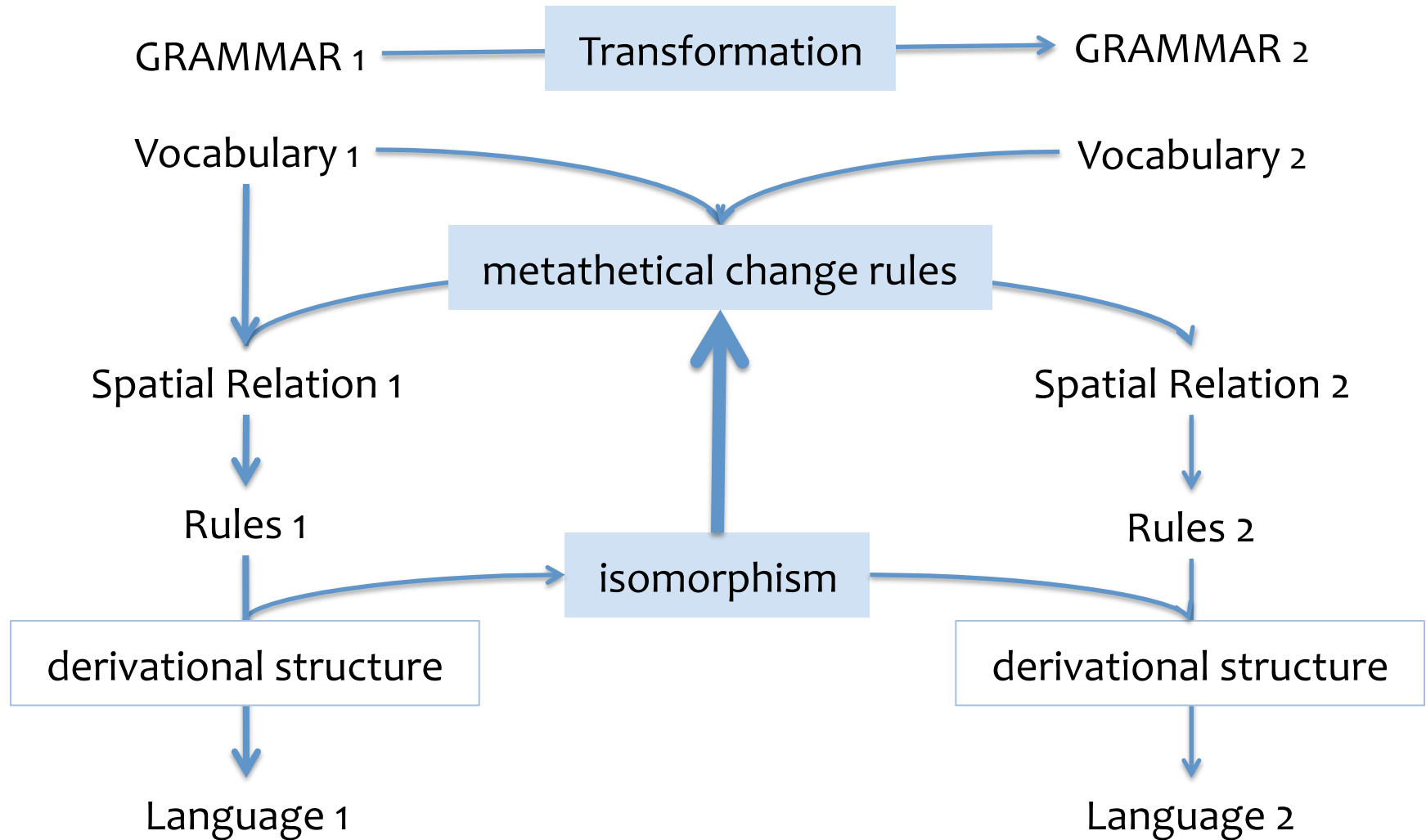
shape equivalence rule



to make this transition idea work one must consider heuristics in how the shape equivalence rules are applied.



Transformation of Grammars



the basic idea

we need to ensure that
grammars are specified in an
normalized fashion –
i.e., *in the same sort of way
every time*

hence, ***grammars*** in ***normal
form***

to compare languages

Vocabulary

Purely Additive rules

Purely Subtractive rules

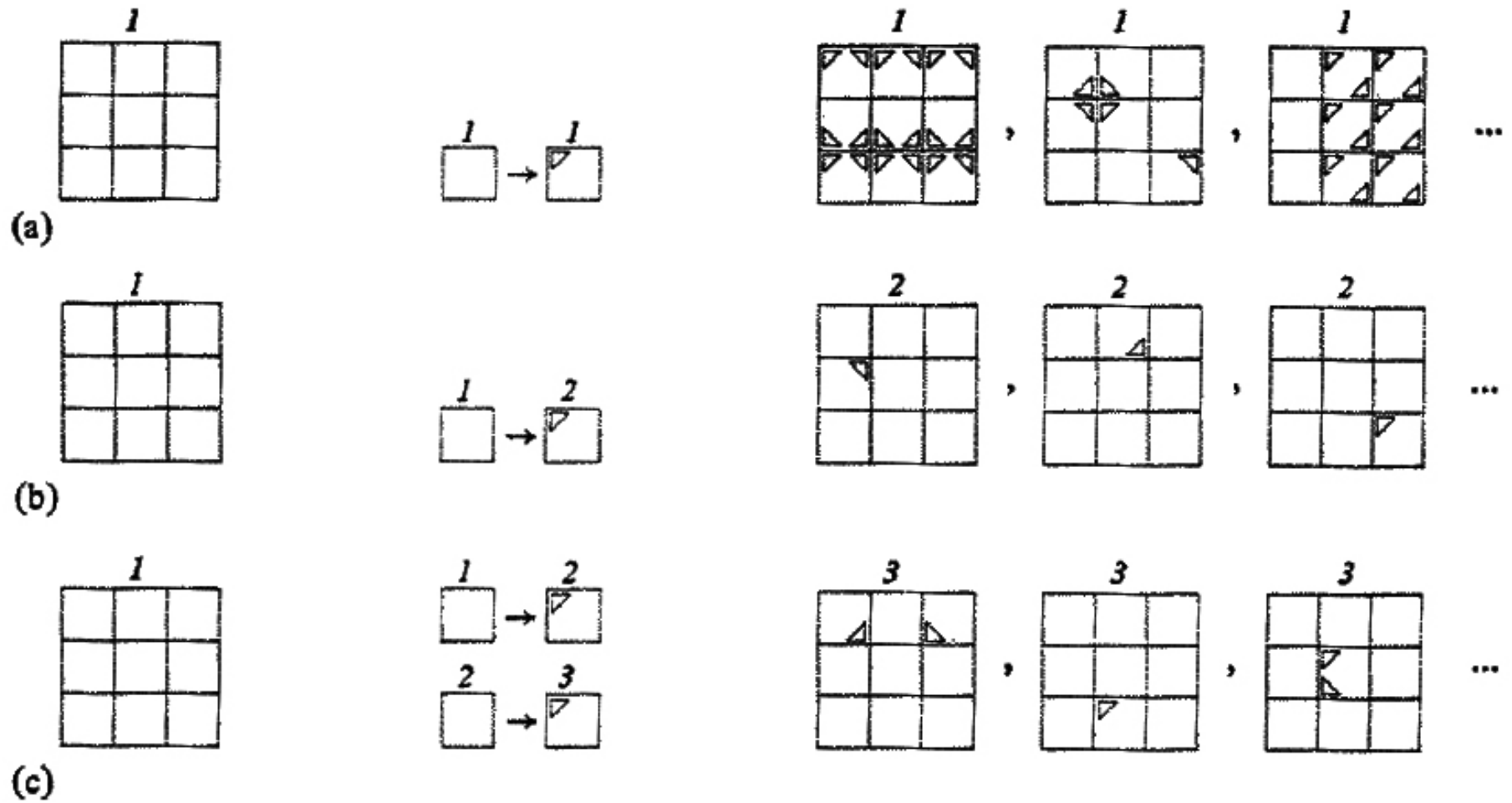
Labels are spatial

– *how*

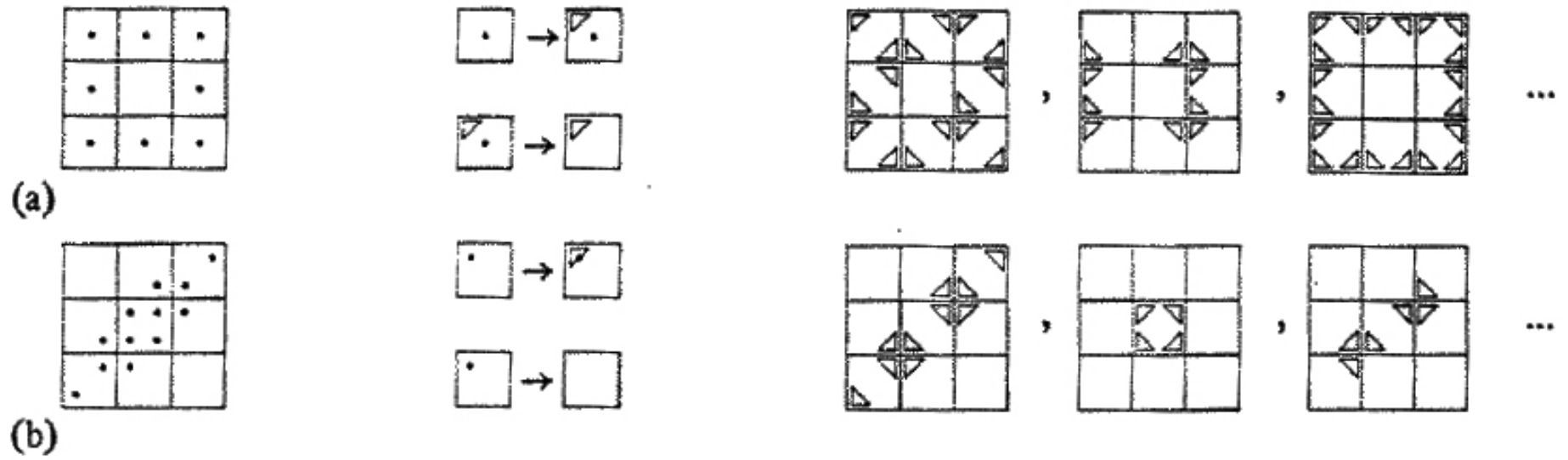
– *where*

States are nonspatial

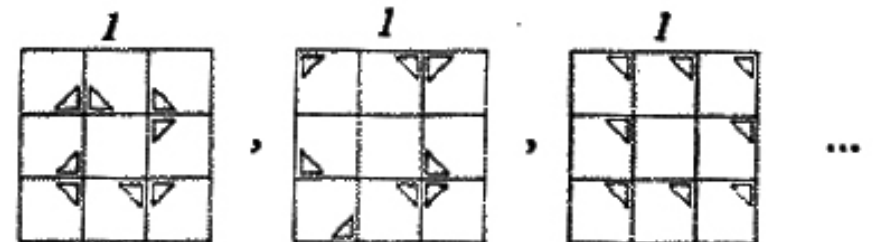
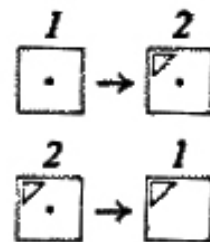
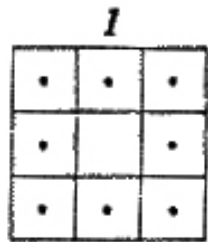
– *when*



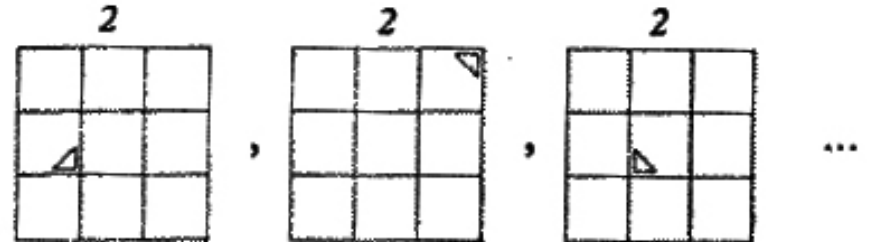
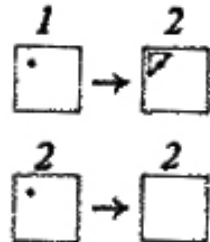
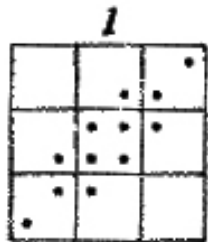
nonspatial or state labels



spatial – *where* and *how* labels



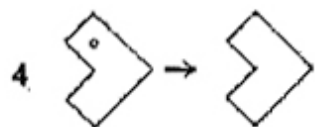
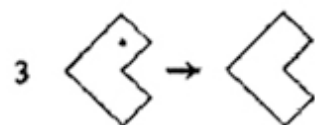
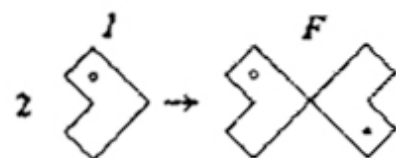
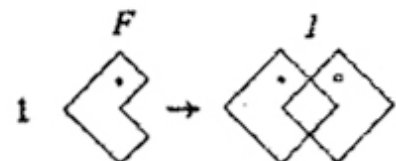
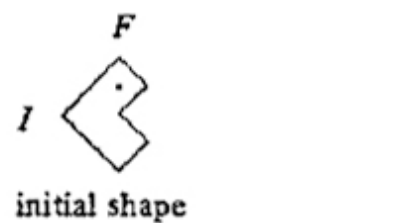
(a)



(b)

state and **spatial labels**

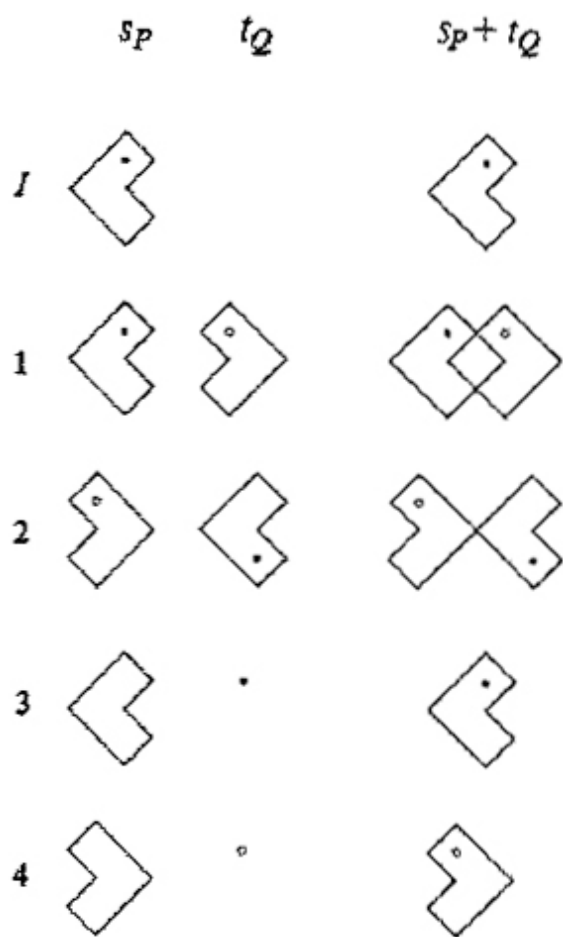
G :



rules

final state: F

(a)



$I \neq F$

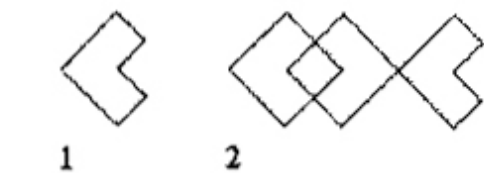
1 (F, I)

2 (I, F)

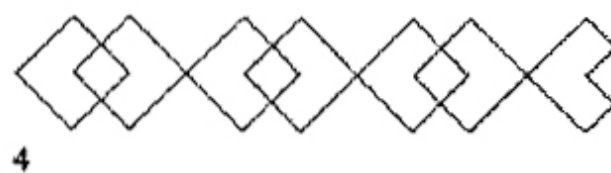
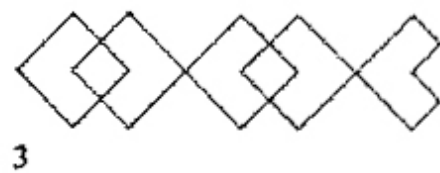
3 $(\#, \#)$

4 $(\#, \#)$

(b) grammar in normal form (c)



(d)



Is a basic property of grammars

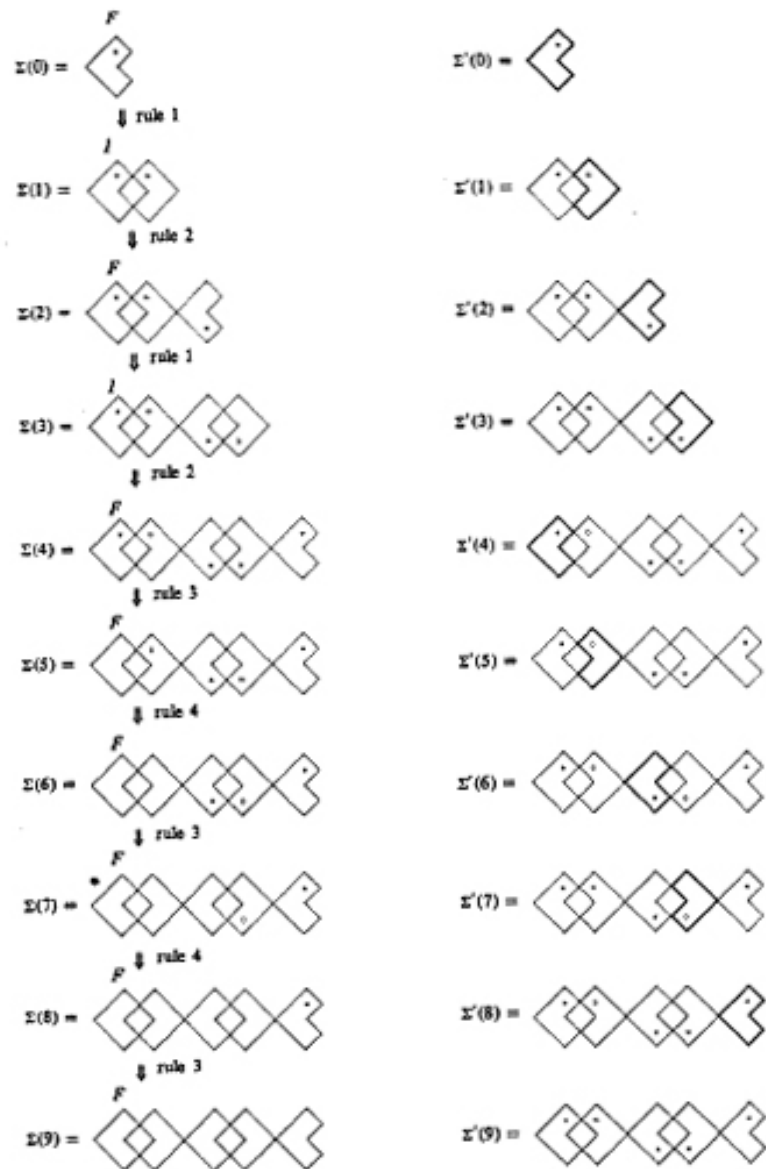
Expresses a relationship on rules, the initial shape and selected typical derivations of designs in the grammar

$R(G) = \{(\text{rule } x, \text{rule } y) \dots \}$ where $(\text{rule } x, \text{rule } y)$ is a member of $R(G)$ whenever

- Rule x is additive or is the initial shape
- Rule y is purely additive or purely subtractive and rule y is applied to that part of the design that includes a subshape of a labeled shape which was added by a previous application of rule x

i.e., rule x makes rule y possible

recursive structure $R(G)$



(c)

$$\begin{aligned}
 R(G) &= R_1(G) \cup R_2(G) \cup R_3(G) \cup R_4(G) \\
 &= \{ (I, \text{rule 1}), (\text{rule 1}, \text{rule 2}), (\text{rule 2}, \text{rule 1}), (I, \text{rule 3}), (\text{rule 1}, \text{rule 4}), (\text{rule 2}, \text{rule 3}) \}
 \end{aligned}$$

(d)

$$R_1(G) = \emptyset$$

$$\begin{aligned} R_2(G) &= \emptyset \cup \{(I, \text{rule } 1)\} \\ &= \{(I, \text{rule } 1)\} \end{aligned}$$

$$\begin{aligned} R_3(G) &= \{(I, \text{rule } 1)\} \cup \{(\text{rule } 1, \text{rule } 2)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2)\} \end{aligned}$$

$$\begin{aligned} R_4(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2)\} \cup \{(\text{rule } 2, \text{rule } 1)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1)\} \end{aligned}$$

$$\begin{aligned} R_5(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1)\} \cup \{(\text{rule } 1, \text{rule } 3)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1)\} \end{aligned}$$

$$\begin{aligned} R_6(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1)\} \cup \{(I, \text{rule } 3)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3)\} \end{aligned}$$

$$\begin{aligned} R_7(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3)\} \cup \{(\text{rule } 1, \text{rule } 4)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4)\} \end{aligned}$$

$$\begin{aligned} R_8(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4)\} \cup \{(\text{rule } 2, \text{rule } 3)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4), (\text{rule } 2, \text{rule } 3)\} \end{aligned}$$

$$\begin{aligned} R_9(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4), (\text{rule } 2, \text{rule } 3)\} \cup \{(\text{rule } 1, \text{rule } 4)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4), (\text{rule } 2, \text{rule } 3)\} \end{aligned}$$

$$\begin{aligned} R_{10}(G) &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4), (\text{rule } 2, \text{rule } 3)\} \cup \{(\text{rule } 2, \text{rule } 3)\} \\ &= \{(I, \text{rule } 1), (\text{rule } 1, \text{rule } 2), (\text{rule } 2, \text{rule } 1), (I, \text{rule } 3), (\text{rule } 1, \text{rule } 4), (\text{rule } 2, \text{rule } 3)\} \end{aligned}$$

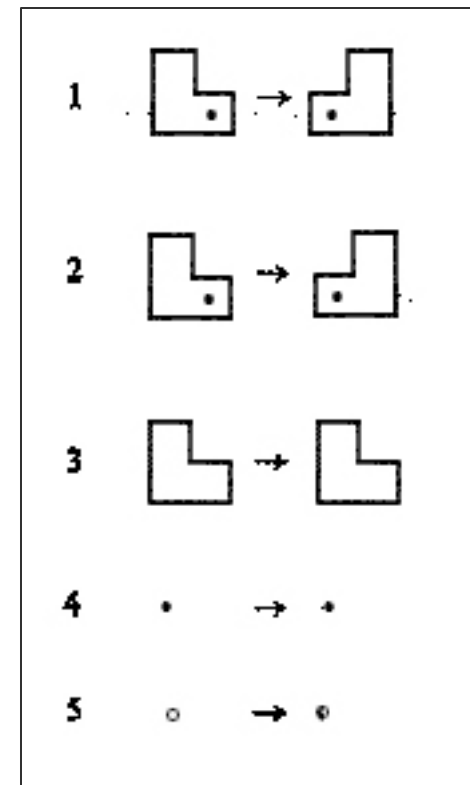
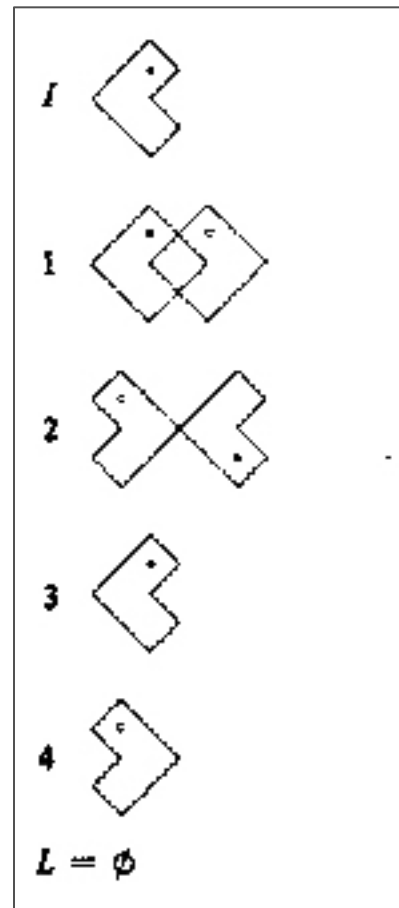
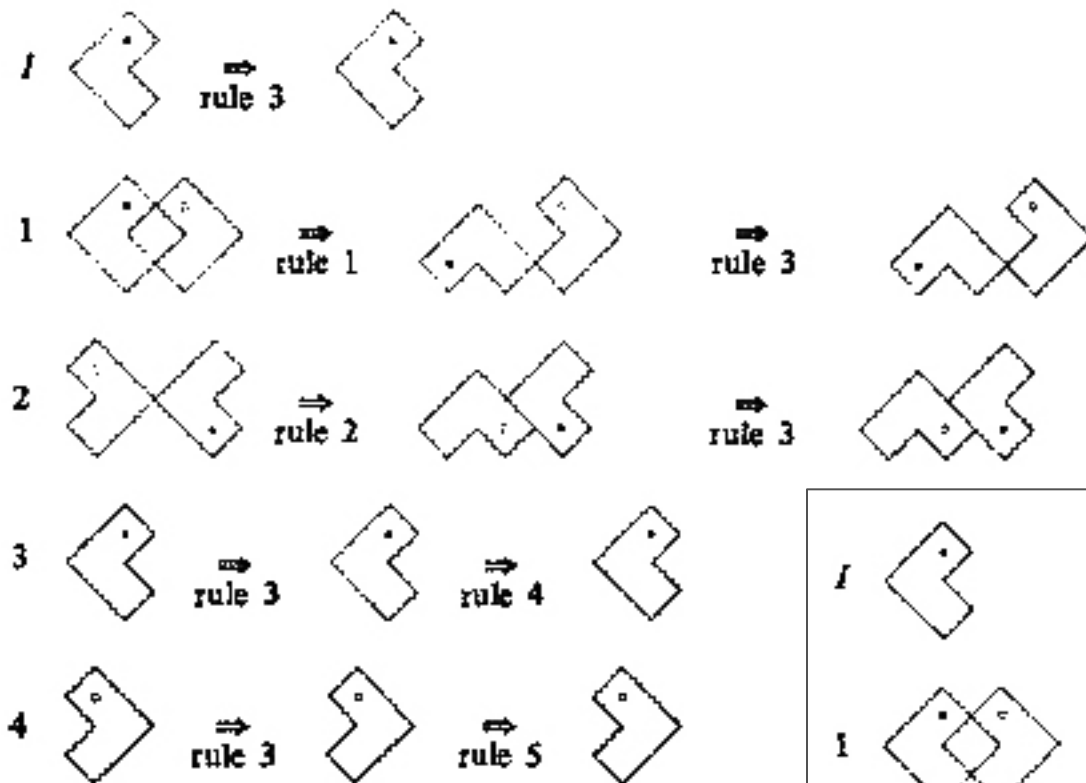
Comprises two independent stages

Defining shape change rules specifying transformation T_A between an initial and final set of relations

Defining state change rules specifying transformation T_B

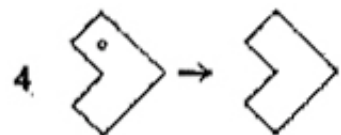
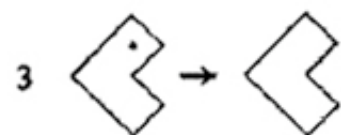
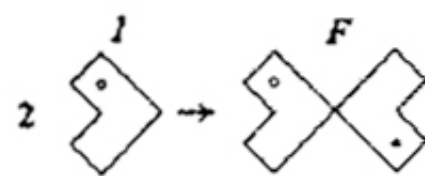
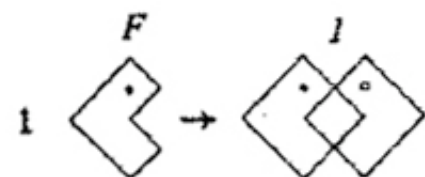
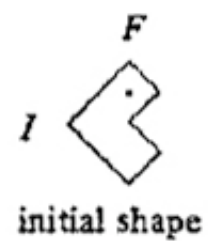
T_A and T_B are combined to produce a complete transformation T of G .

transformation of grammars



deriving a final set
of relations

G :

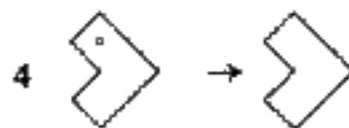
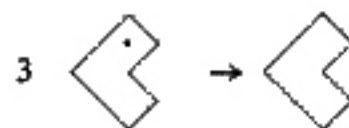
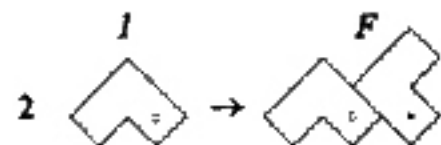
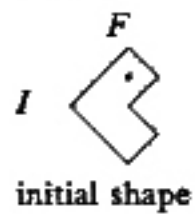


rules

final state: F

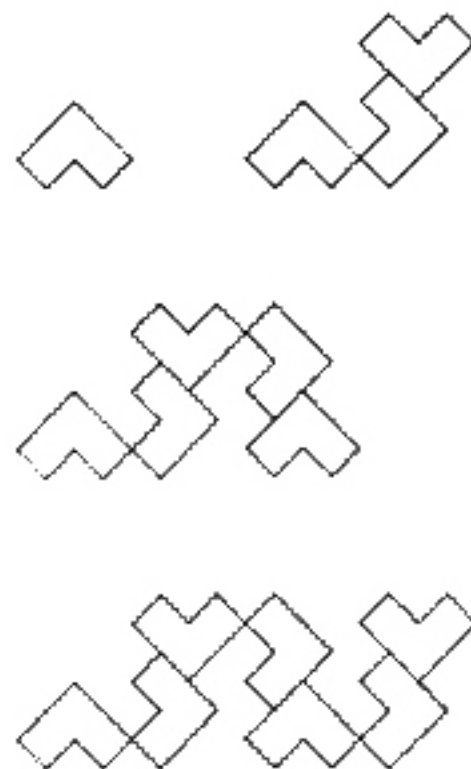
(a)

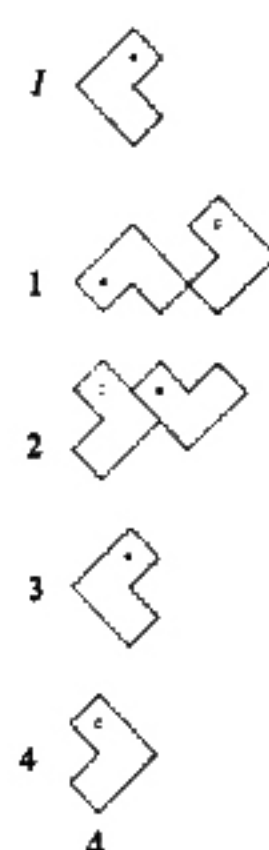
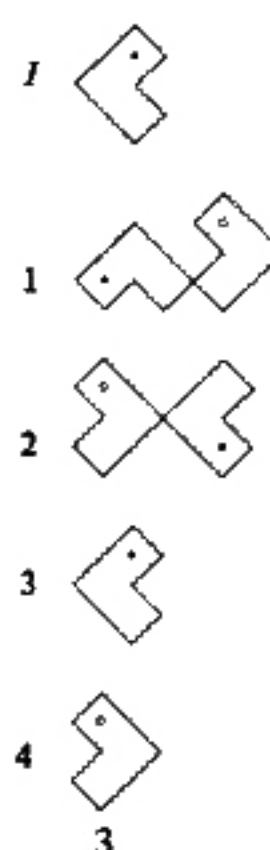
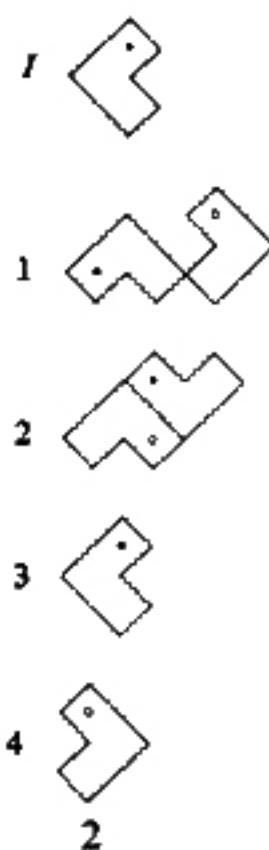
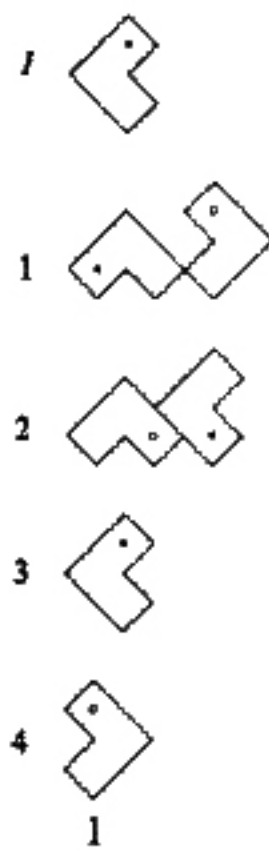
$T_A(G)$:

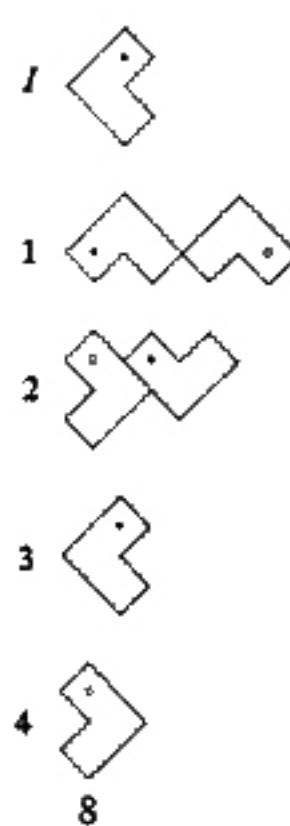
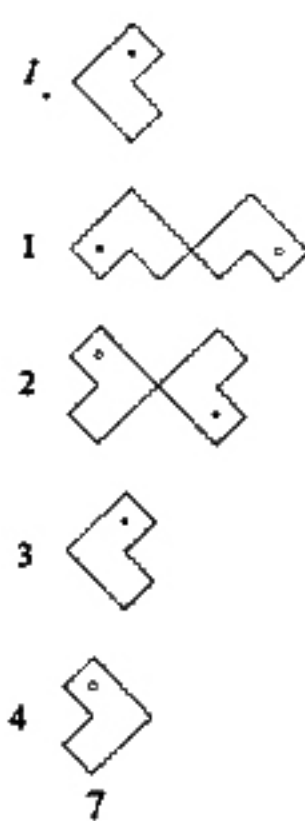
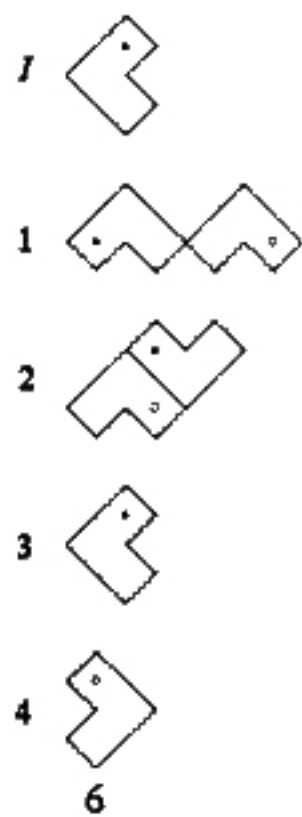
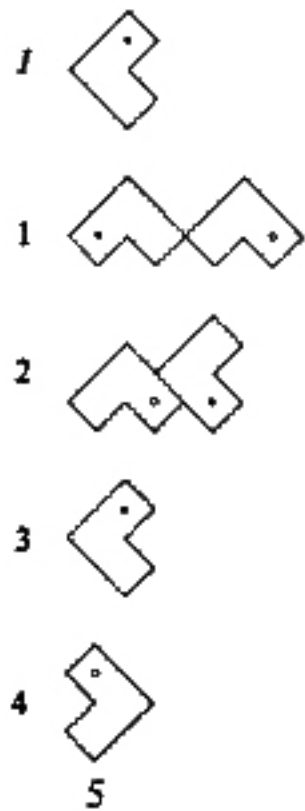


rules

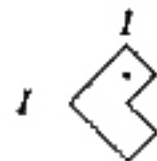
final state: F



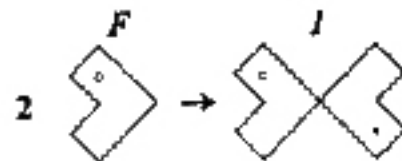
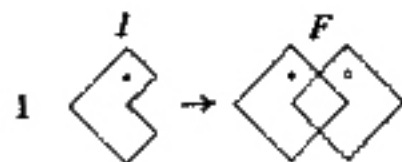




$T_B(G)$:



initial shape



rules

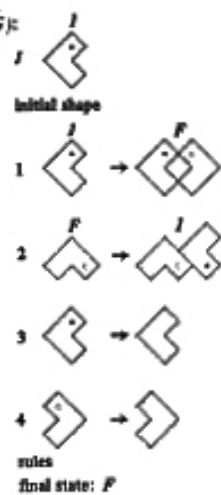
final state: F

(a)

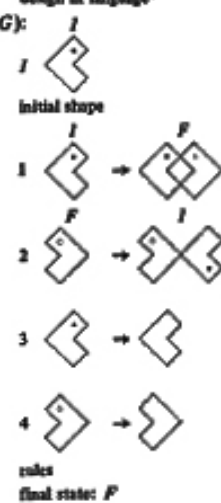


(b)

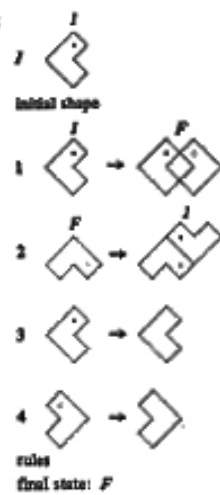
$T_9(G):$



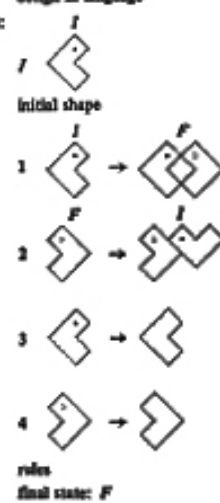
$T_{11}(G):$



$T_{10}(G):$

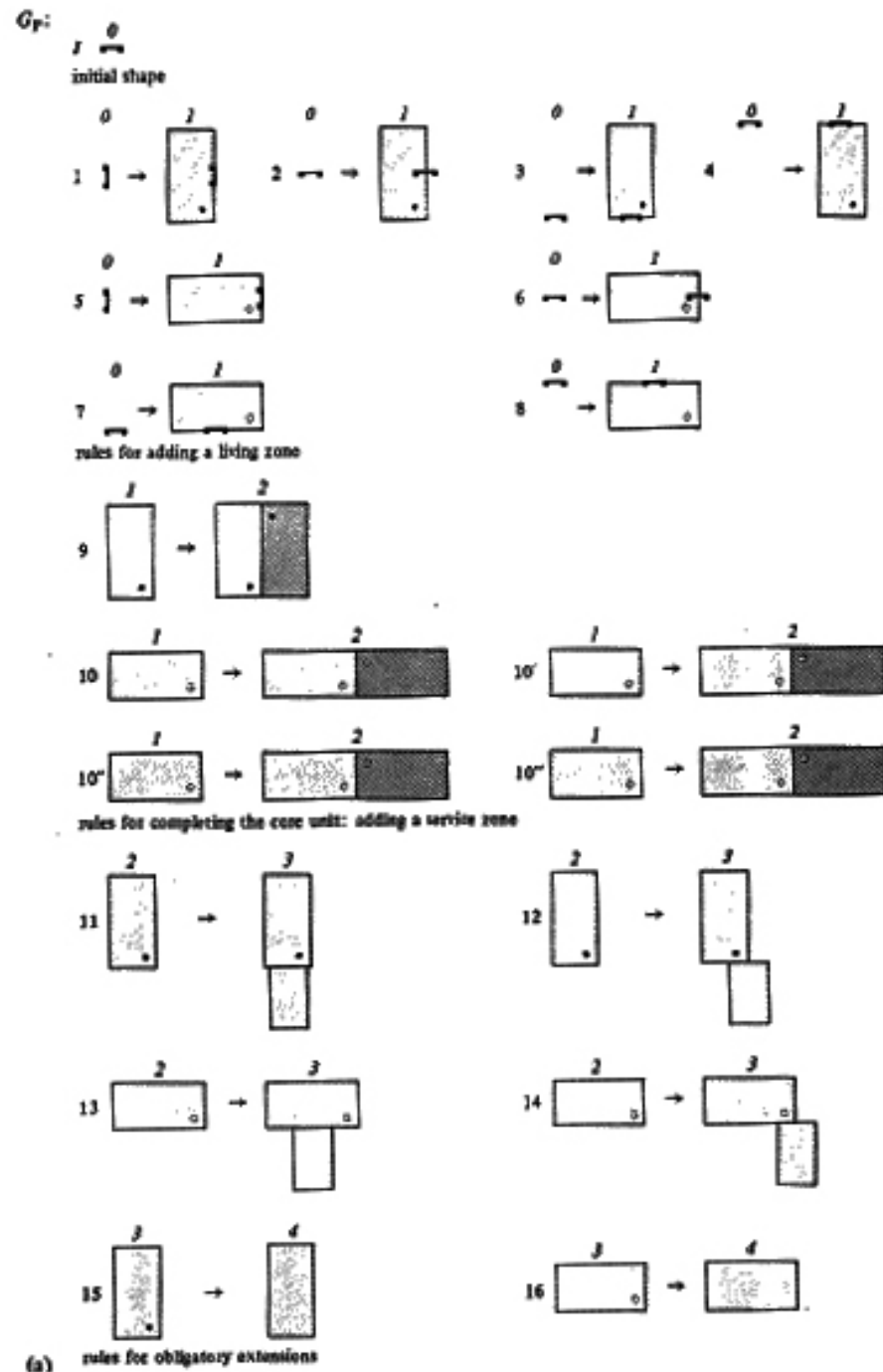


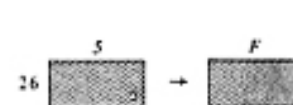
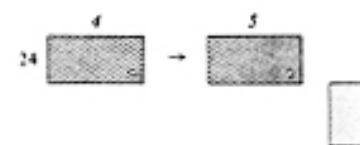
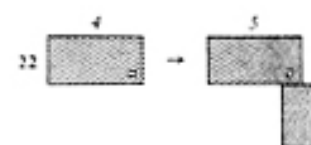
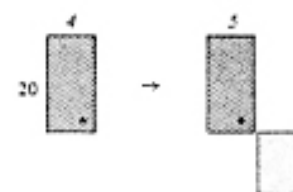
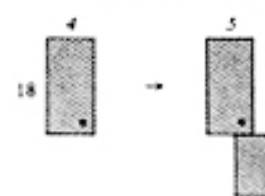
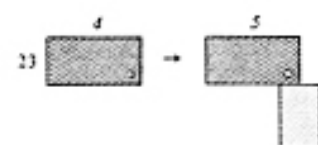
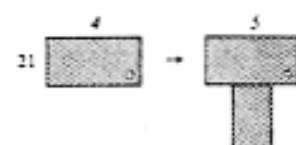
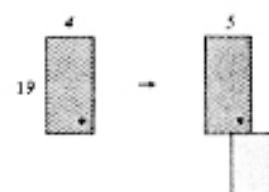
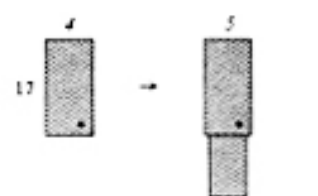
$T_{12}(G):$



Prairie to Usonian

the prairie house grammar in formal form

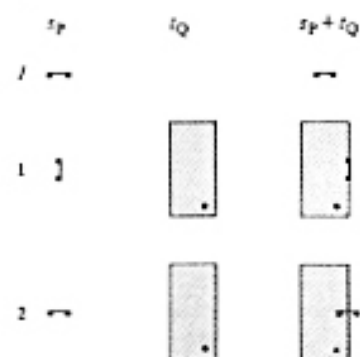




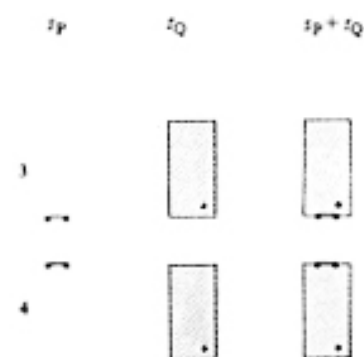
rules for obligatory extensions (continued)

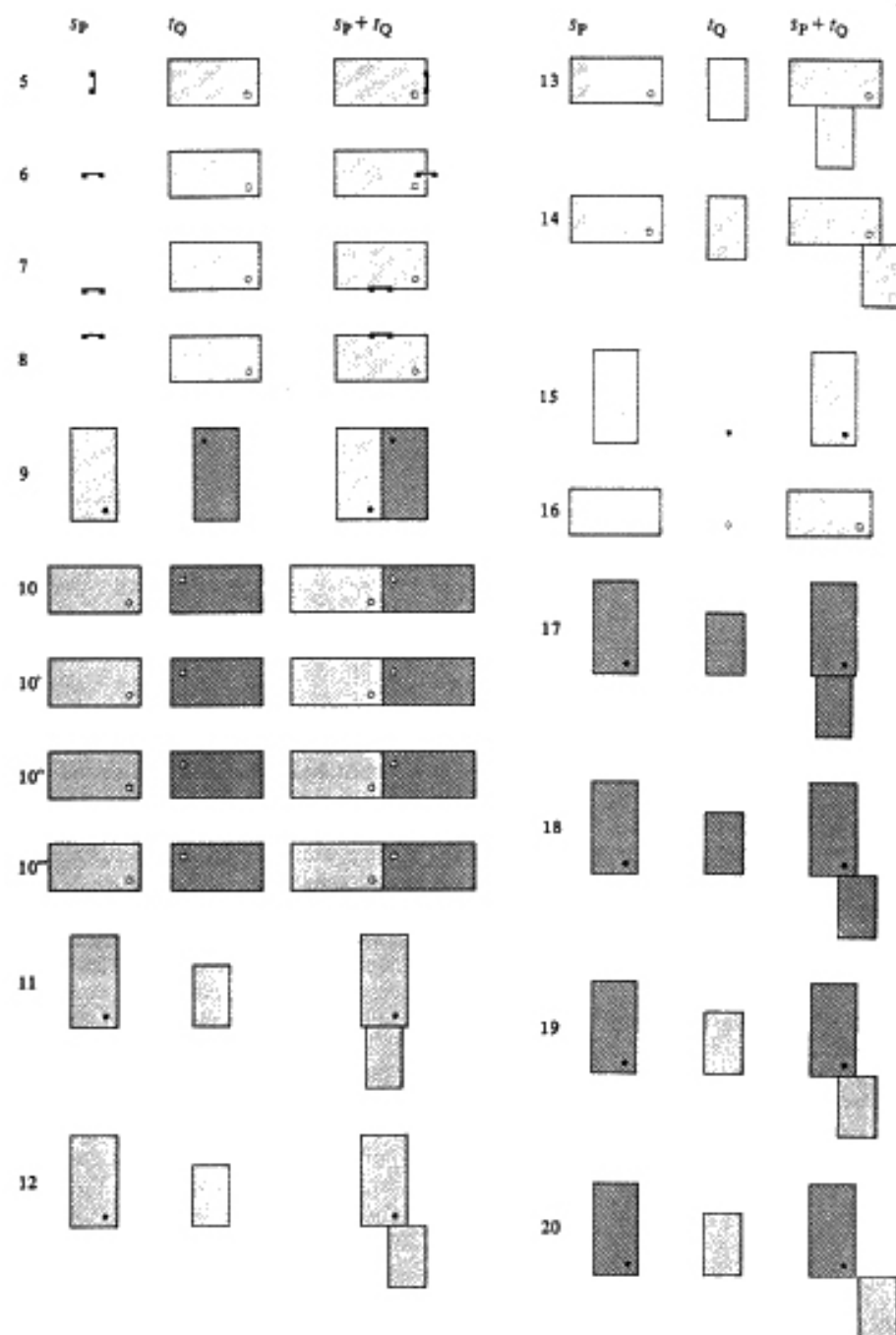
final state: F

(a) continued

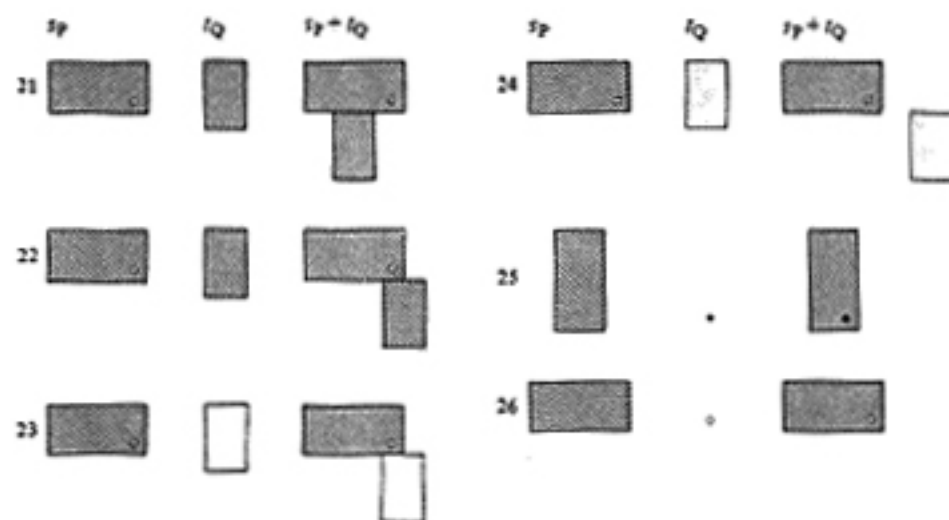


(b)





(b) continued

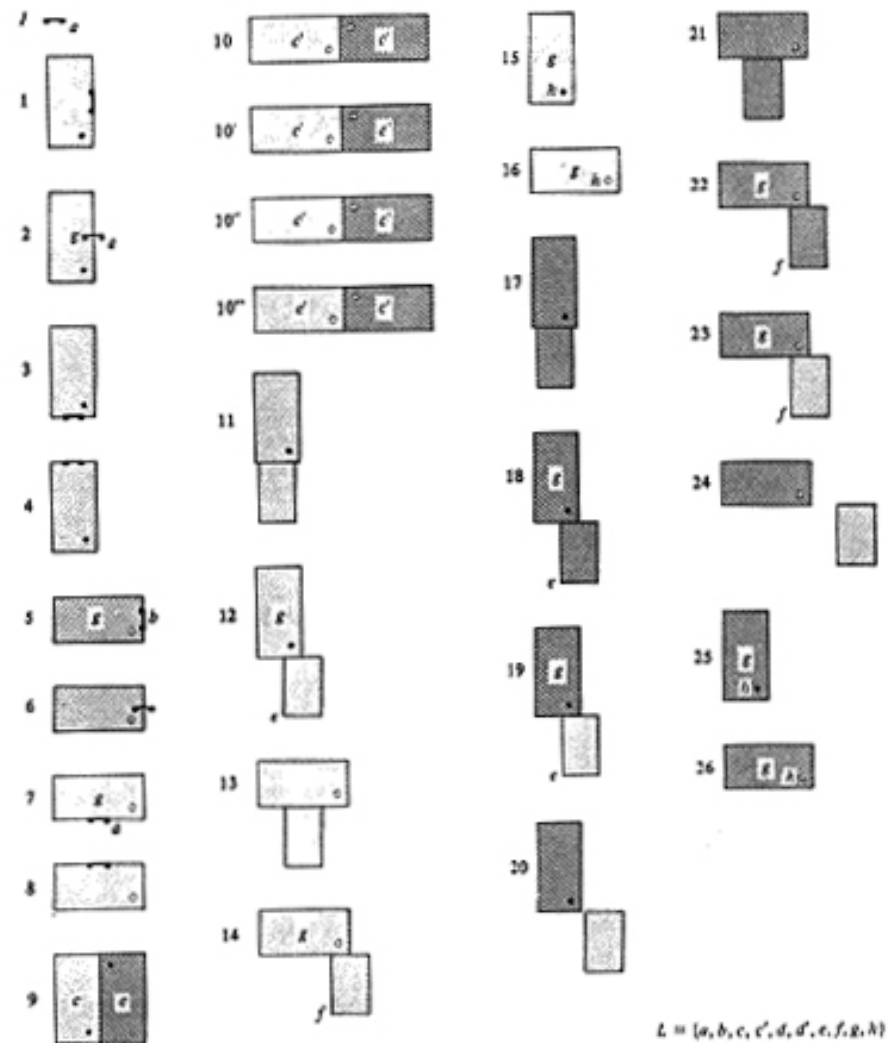
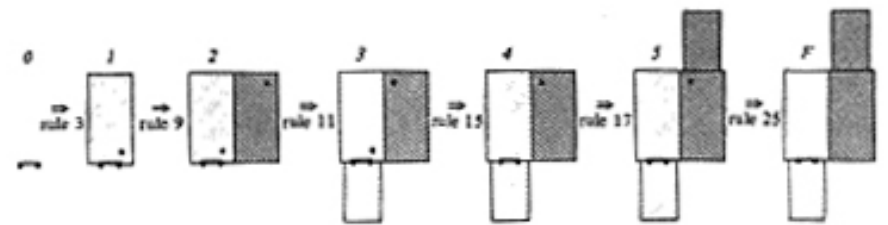


(b) continued

1	(a, 0)	6	(0, 1)	10'	(1, 2)	15	(2, 3)	21	(4, 5)
2	(0, 1)	7	(0, 1)	10''	(1, 2)	16	(2, 4)	22	(4, 5)
3	(0, 1)	8	(0, 1)	11	(2, 3)	17	(4, 5)	23	(4, 5)
4	(0, 1)	9	(1, 2)	12	(2, 3)	18	(4, 5)	24	(4, 5)
5	(0, 1)	10	(1, 2)	13	(2, 3)	19	(4, 5)	25	(5, 6)
		10'	(1, 2)	14	(2, 3)	20	(4, 5)	26	(5, 6)

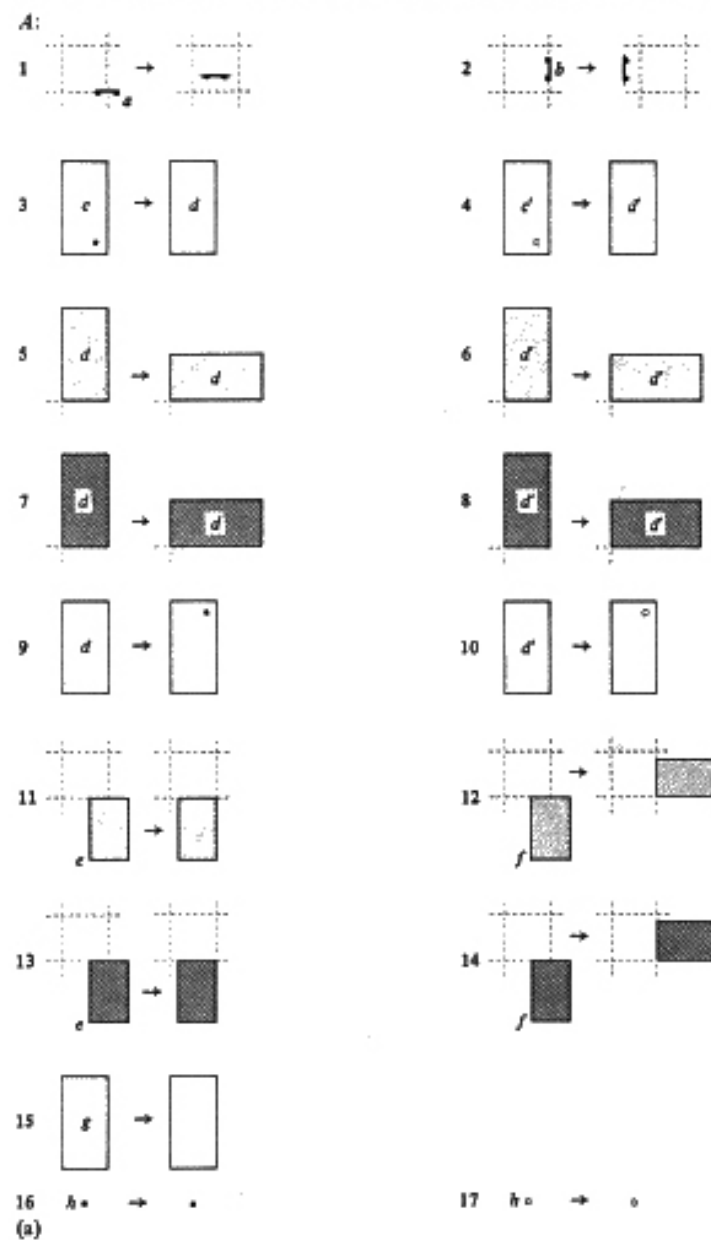
(c)

initial set of spatial relations



$L = \{a, b, c, c', d, d', e, f, g, h\}$

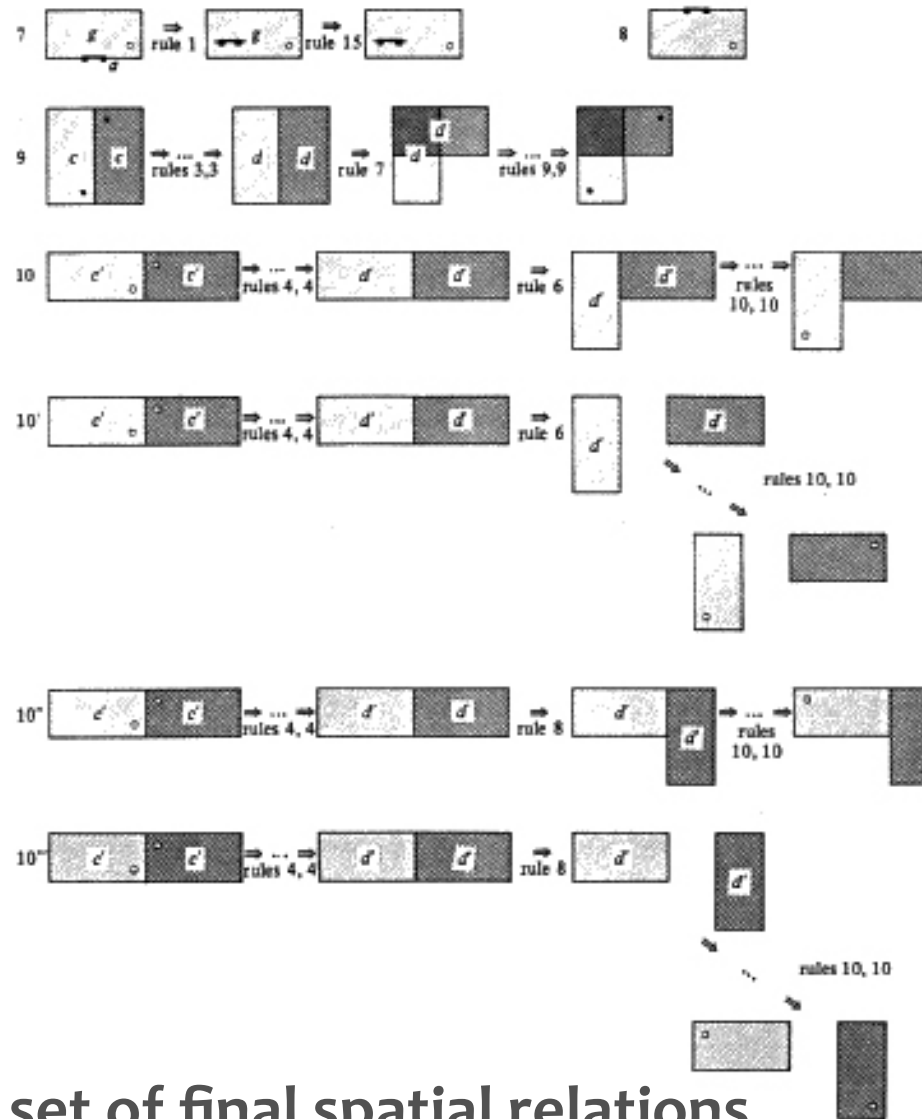
change rules



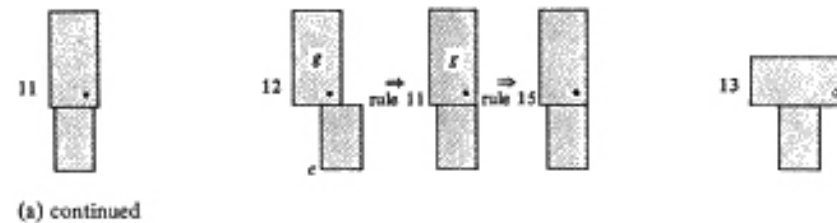
B:
 $F_1, F_2, F_3, F_7, F_9, F_{10}, F_{10'}, F_{10''}, F_{10'''}:$ identity state change rules

$F_{22}, F_{26}:$	$(2, 3) \rightarrow (2, 4)$
$F_{18}, F_{16}:$	$(3, 4) \rightarrow (8, 3)$
$F_{10}, F_{19}, F_{22}, F_{23}:$	$(4, 5) \rightarrow (3, F)$
$F_{20}, F_{24}:$	$(5, F) \rightarrow (8, F)$

(b)

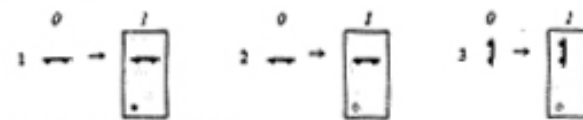


deriving a set of final spatial relations

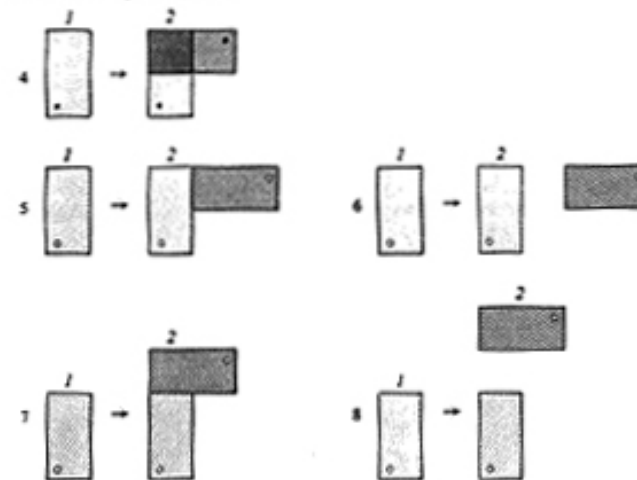


$$T(G_P) = G_U;$$

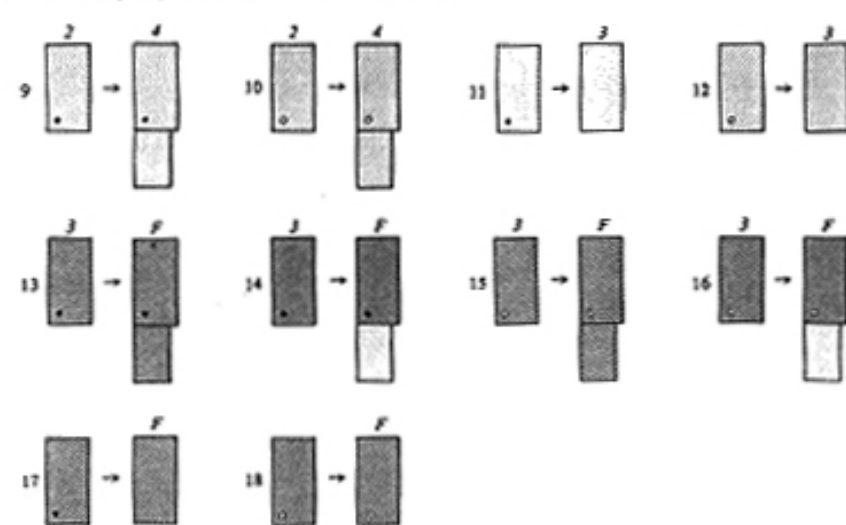
$I \xrightarrow{0}$
initial shape



rules for adding a living zone



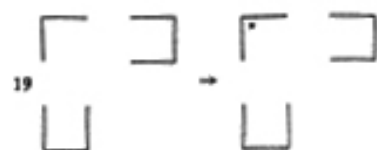
rules for completing the core unit: adding a bedroom zone



rules for optional extensions

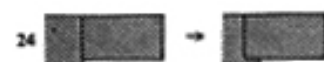
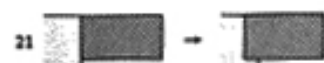
final state: F

transformation !

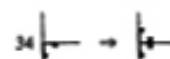
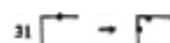
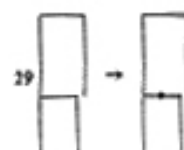


20

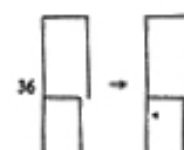
rules for extending the "hinge" of the core unit



rules for interpenetrating zones



rules for adding
secondary fireplaces



rules for extending
extensions



rules for indenting
exterior corners of
the plan

