



Machine-Level Programming I: Basics

14-513/18-613: Introduction to Computer Systems
5th Lecture, May 27, 2020

Today: Machine Programming I: Basics

- History of Intel processors and architectures
- **Assembly Basics: Registers, operands, move**
- Arithmetic & logical operations
- C, assembly, machine code

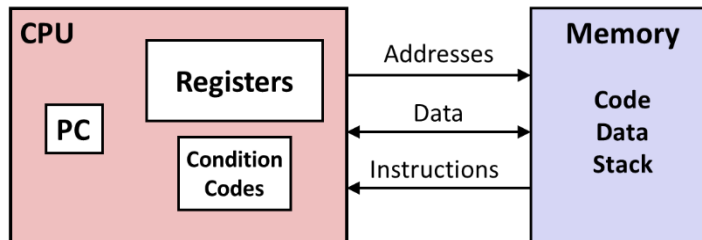
Levels of Abstraction

C programmer

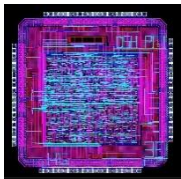
```
#include <stdio.h>
int main(){
    int i, n = 10, t1 = 0, t2 = 1, nxt;
    for (i = 1; i <= n; ++i){
        printf("%d, ", t1);
        nxt = t1 + t2;
        t1 = t2;
        t2 = nxt; }
    return 0; }
```

**Nice clean layers,
but beware...**

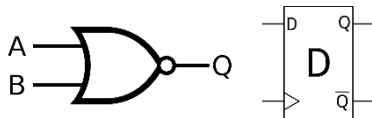
Assembly programmer



Computer Designer



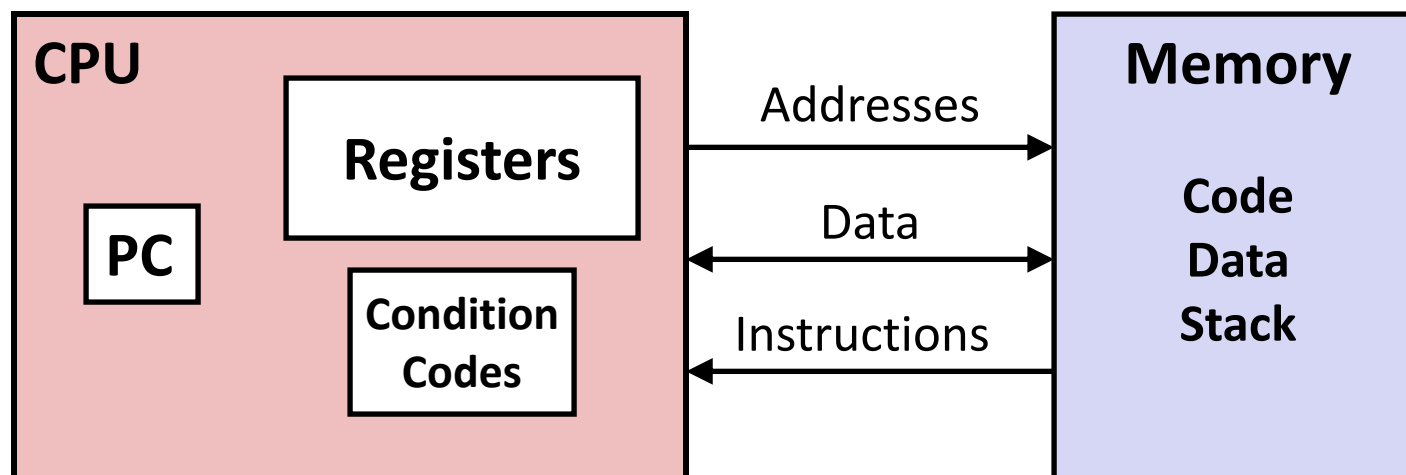
Gates, clocks, circuit layout, ...



Definitions

- **Architecture:** (also ISA: instruction set architecture) The parts of a processor design that one needs to understand for writing correct machine/assembly code
 - Examples: instruction set specification, registers
 - **Machine Code:** The byte-level programs that a processor executes
 - **Assembly Code:** A text representation of machine code
- **Microarchitecture: Implementation of the architecture**
 - Examples: cache sizes and core frequency
- **Example ISAs:**
 - Intel: x86, IA32, Itanium, x86-64
 - ARM: Used in almost all mobile phones
 - RISC V: New open-source ISA

Assembly/Machine Code View



Programmer-Visible State

- **PC: Program counter**
 - Address of next instruction
 - Called “RIP” (x86-64)
- **Register file**
 - Heavily used program data
- **Condition codes**
 - Store status information about most recent arithmetic or logical operation
 - Used for conditional branching
- **Memory**
 - Byte addressable array
 - Code and user data
 - Stack to support procedures

Assembly Characteristics: Data Types

- **“Integer” data of 1, 2, 4, or 8 bytes**
 - Data values
 - Addresses (untyped pointers)
- **Floating point data of 4, 8, or 10 bytes**
- **(SIMD vector data types of 8, 16, 32 or 64 bytes)**
- **Code: Byte sequences encoding series of instructions**
- **No aggregate types such as arrays or structures**
 - Just contiguously allocated bytes in memory

x86-64 Integer Registers

<code>%rax</code>	<code>%eax</code>
<code>%rbx</code>	<code>%ebx</code>
<code>%rcx</code>	<code>%ecx</code>
<code>%rdx</code>	<code>%edx</code>
<code>%rsi</code>	<code>%esi</code>
<code>%rdi</code>	<code>%edi</code>
<code>%rsp</code>	<code>%esp</code>
<code>%rbp</code>	<code>%ebp</code>

<code>%r8</code>	<code>%r8d</code>
<code>%r9</code>	<code>%r9d</code>
<code>%r10</code>	<code>%r10d</code>
<code>%r11</code>	<code>%r11d</code>
<code>%r12</code>	<code>%r12d</code>
<code>%r13</code>	<code>%r13d</code>
<code>%r14</code>	<code>%r14d</code>
<code>%r15</code>	<code>%r15d</code>

- Can reference low-order 4 bytes (also low-order 1 & 2 bytes)
- Not part of memory (or cache)

Assembly Characteristics: Operations

- **Transfer data between memory and register**
 - Load data from memory into register
 - Store register data into memory

- **Perform arithmetic function on register or memory data**

- **Transfer control**
 - Unconditional jumps to/from procedures
 - Conditional branches
 - Indirect branches

Memory Addressing Modes

■ Most General Form

$D(Rb, Ri, S)$ $Mem[Reg[Rb]+S*Reg[Ri]+ D]$

- D: Constant “displacement” 1, 2, or 4 bytes
- Rb: Base register: Any of 16 integer registers
- Ri: Index register: Any, except for `%rsp`
- S: Scale: 1, 2, 4, or 8 (*why these numbers?*)

■ Special Cases

(Rb, Ri) $Mem[Reg[Rb]+Reg[Ri]]$

$D(Rb, Ri)$ $Mem[Reg[Rb]+Reg[Ri]+D]$

(Rb, Ri, S) $Mem[Reg[Rb]+S*Reg[Ri]]$

Address Computation Examples

<code>%rdx</code>	<code>0xf000</code>
<code>%rcx</code>	<code>0x0100</code>

 $D(Rb, Ri, S)$
 $Mem[Reg[Rb]+S*Reg[Ri]+ D]$

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Expression	Address Computation	Address
<code>0x8 (%rdx)</code>		
<code>(%rdx, %rcx)</code>		
<code>(%rdx, %rcx, 4)</code>		
<code>0x80 (, %rdx, 2)</code>		

Address Computation Examples

<code>%rdx</code>	<code>0xf000</code>
<code>%rcx</code>	<code>0x0100</code>

Expression	Address Computation	Address
<code>0x8 (%rdx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%rdx, %rcx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%rdx, %rcx, 4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(, %rdx, 2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

Moving Data

■ Moving Data

`movq Source, Dest`

■ Operand Types

- **Immediate:** Constant integer data
 - Example: `$0x400`, `$-533`
 - Like C constant, but prefixed with ``$'`
 - Encoded with 1, 2, or 4 bytes
- **Register:** One of 16 integer registers
 - Example: `%rax`, `%r13`
 - But `%rsp` reserved for special use
 - Others have special uses for particular instructions
- **Memory** 8 consecutive bytes of memory at address given by register
 - Simplest example: `(%rax)`
 - Various other “addressing modes”

`%rax`

`%rcx`

`%rdx`

`%rbx`

`%rsi`

`%rdi`

`%rsp`

`%rbp`

`%rN`

**Warning: Intel docs use
`mov Dest, Source`**

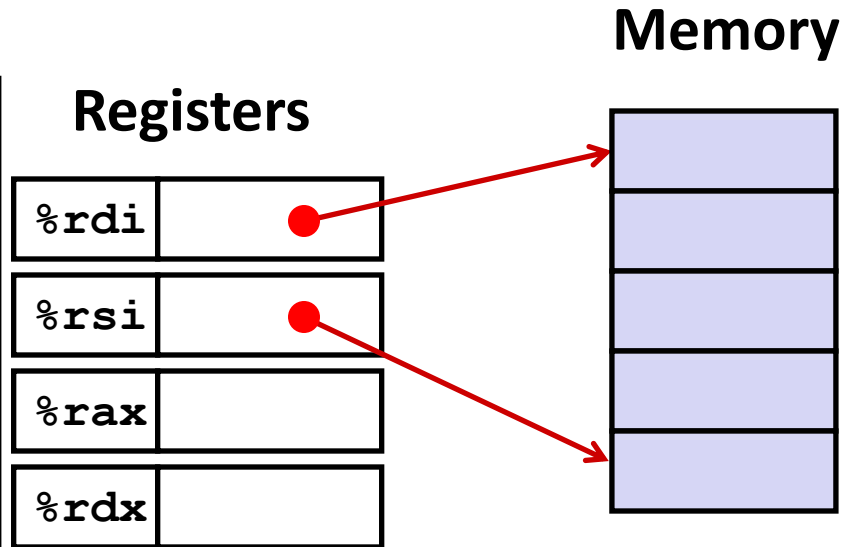
movq Operand Combinations

	Source	Dest	Src, Dest	C Analog
movq	Imm	Reg	movq \$0x4, %rax	temp = 0x4;
		Mem	movq \$-147, (%rax)	*p = -147;
	Reg	Reg	movq %rax, %rdx	temp2 = temp1;
		Mem	movq %rax, (%rdx)	*p = temp;
	Mem	Reg	movq (%rax), %rdx	temp = *p;

Cannot do memory-memory transfer with a single instruction

Understanding Swap()

```
void swap
(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```



Register	Value
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret
```

Understanding Swap()

Registers

<code>%rdi</code>	<code>0x120</code>
<code>%rsi</code>	<code>0x100</code>
<code>%rax</code>	
<code>%rdx</code>	

Memory

	Address
123	<code>0x120</code>
	<code>0x118</code>
	<code>0x110</code>
	<code>0x108</code>
456	<code>0x100</code>

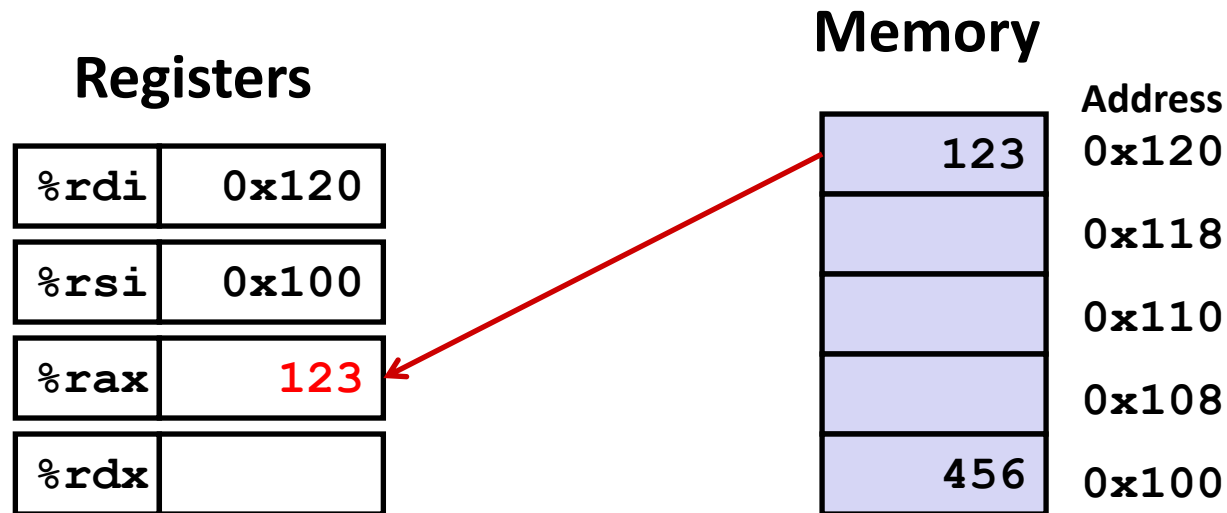
swap:

```

movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret

```


Understanding Swap()



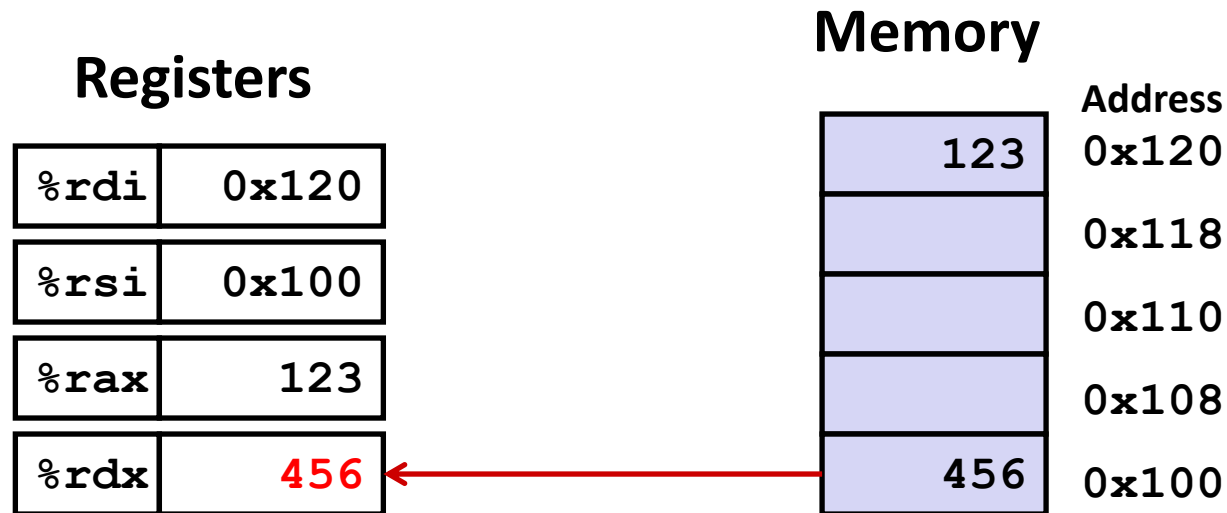
swap:

```

movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret

```

Understanding Swap()



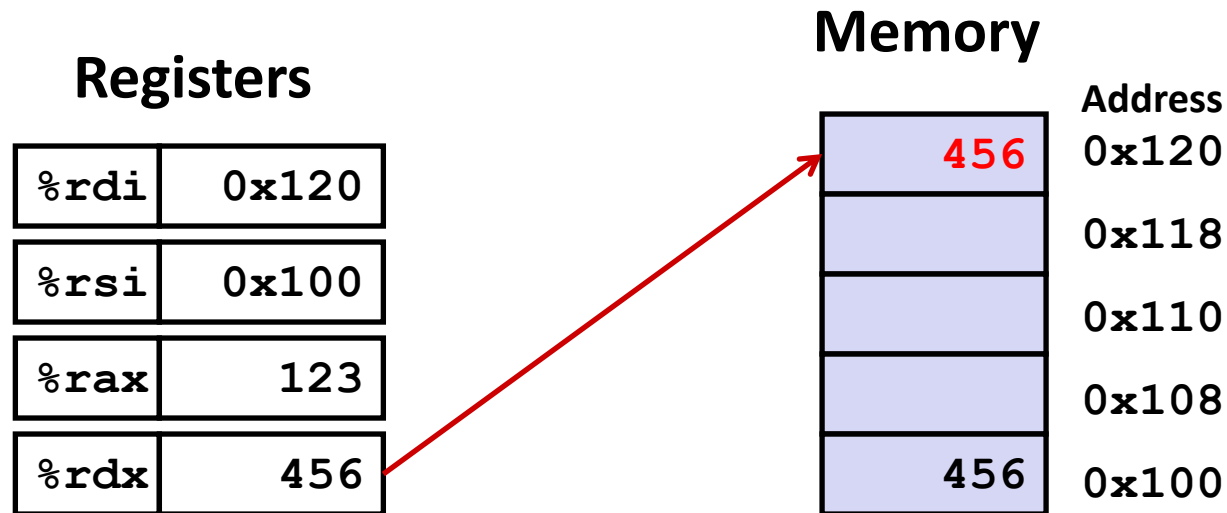
swap:

```

movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
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ret

```

Understanding Swap()



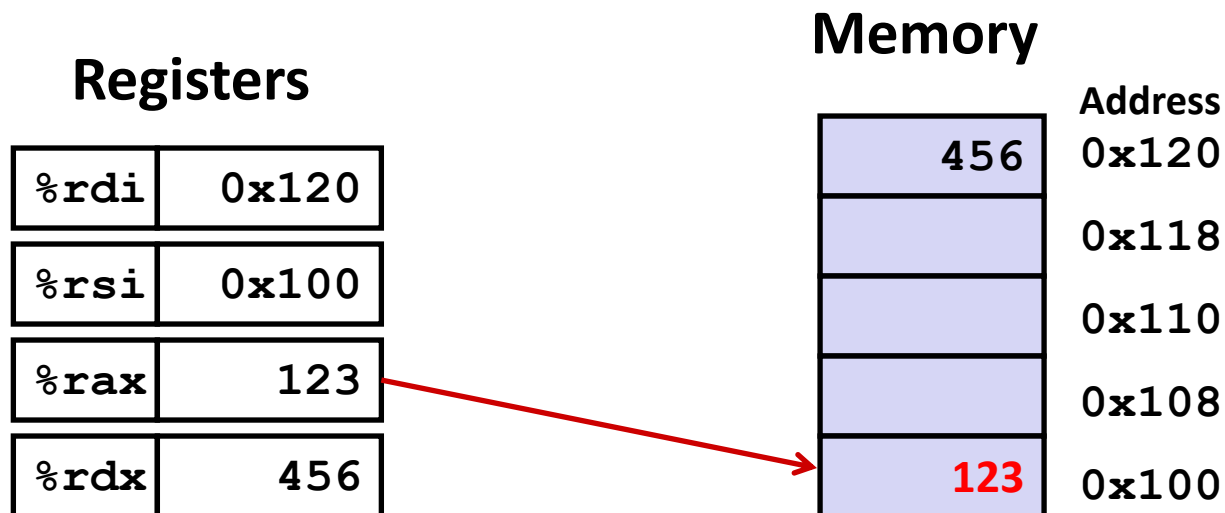
swap:

```

movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret

```

Understanding Swap()



swap:

```

movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret

```

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- History of Intel processors and architectures
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- **Arithmetic & logical operations**
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Address Computation Instruction

■ `leaq Src, Dst`

- *Src* is address mode expression
- Set *Dst* to address denoted by expression

■ Uses

- Computing addresses without a memory reference
 - E.g., translation of `p = &x[i];`
- Computing arithmetic expressions of the form $x + k*y$
 - $k = 1, 2, 4, \text{ or } 8$

■ Example

```
long m12(long x)
{
    return x*12;
}
```

Converted to ASM by compiler:

```
leaq (%rdi,%rdi,2), %rax # t = x+2*x
salq $2, %rax           # return t<<2
```

Some Arithmetic Operations

■ Two Operand Instructions:

Format

Computation

addq *Src, Dest* Dest = Dest + Src

subq *Src, Dest* Dest = Dest – Src

imulq *Src, Dest* Dest = Dest * Src

shlq *Src, Dest* Dest = Dest << Src

Synonym: salq

sarq *Src, Dest* Dest = Dest >> Src

Arithmetic

shrq *Src, Dest* Dest = Dest >> Src

Logical

xorq *Src, Dest* Dest = Dest ^ Src

andq *Src, Dest* Dest = Dest & Src

orq *Src, Dest* Dest = Dest | Src

■ Watch out for argument order! *Src, Dest* (Warning: Intel docs use “op *Dest, Src*”)

■ No distinction between signed and unsigned int (why?)

Some Arithmetic Operations

■ One Operand Instructions

`incq` *Dest* $Dest = Dest + 1$

`decq` *Dest* $Dest = Dest - 1$

`negq` *Dest* $Dest = -Dest$

`notq` *Dest* $Dest = \sim Dest$

■ See book for more instructions

- Depending how you count, there are 2,034 total x86 instructions
- (If you count all addr modes, op widths, flags, it's actually 3,683)

Arithmetic Expression Example

```
long arith
(long x, long y, long z)
{
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
    leaq    (%rdi,%rsi), %rax
    addq    %rdx, %rax
    leaq    (%rsi,%rsi,2), %rdx
    salq    $4, %rdx
    leaq    4(%rdi,%rdx), %rcx
    imulq   %rcx, %rax
    ret
```

Interesting Instructions

- **leaq**: address computation
- **salq**: shift
- **imulq**: multiplication
 - Curious: only used once...

Understanding Arithmetic Expression

Example

```

long arith
(long x, long y, long z)
{
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}

```

```

arith:
    leaq    (%rdi,%rsi), %rax    # t1
    addq    %rdx, %rax          # t2
    leaq    (%rsi,%rsi,2), %rdx
    salq    $4, %rdx           # t4
    leaq    4(%rdi,%rdx), %rcx  # t5
    imulq   %rcx, %rax         # rval
    ret

```

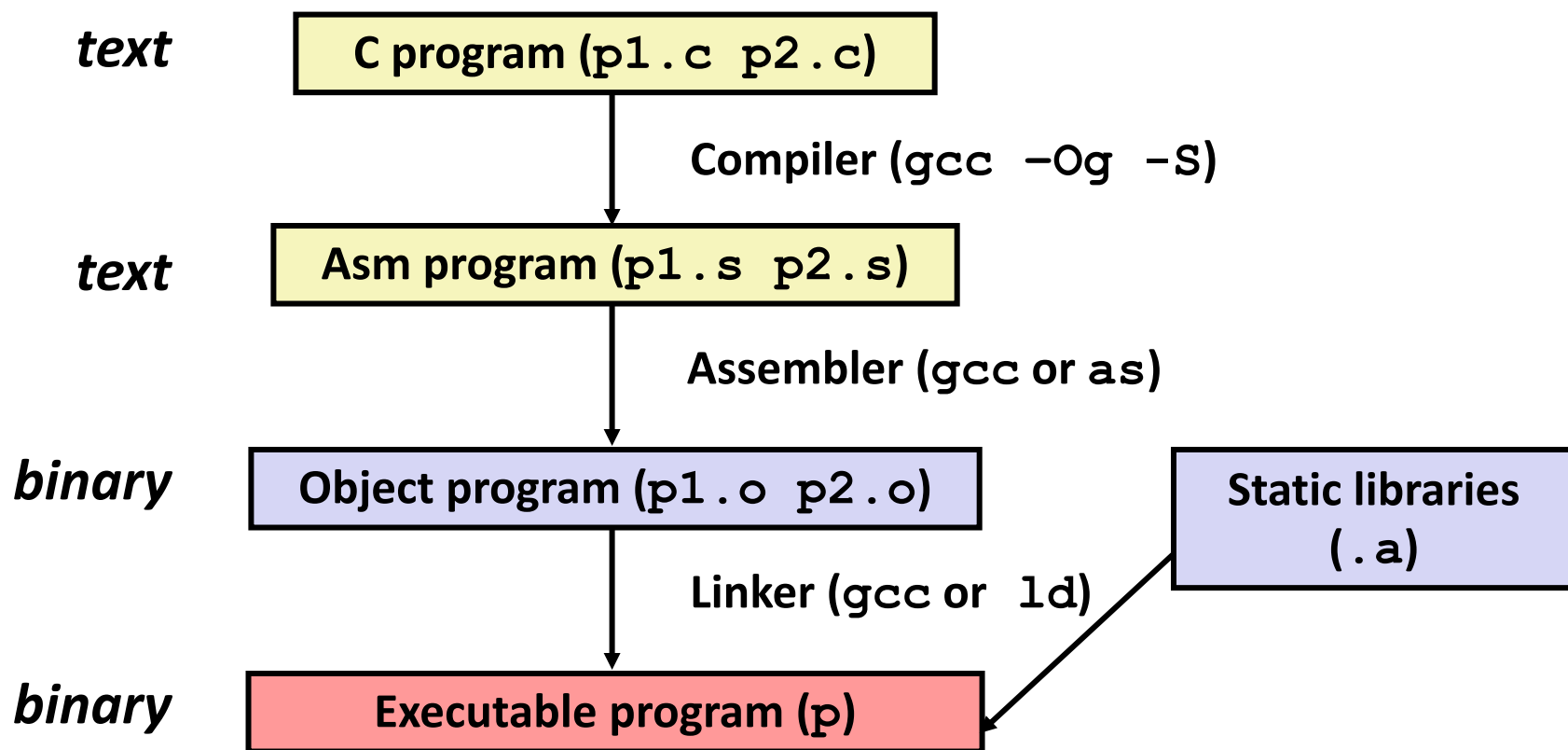
Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z , t4
%rax	t1, t2, rval
%rcx	t5

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Turning C into Object Code

- Code in files `p1.c` `p2.c`
- Compile with command: `gcc -Og p1.c p2.c -o p`
 - Use basic optimizations (`-Og`) [New to recent versions of GCC]
 - Put resulting binary in file `p`



Compiling Into Assembly

C Code (sum.c)

```
long plus(long x, long y);

void sumstore(long x, long y,
              long *dest)
{
    long t = plus(x, y);
    *dest = t;
}
```

Generated x86-64 Assembly

```
sumstore:
    pushq    %rbx
    movq    %rdx, %rbx
    call    plus
    movq    %rax, (%rbx)
    popq    %rbx
    ret
```

Obtain (on shark machine) with command

```
gcc -Og -S sum.c
```

Produces file `sum.s`

Warning: Will get very different results on non-Shark machines (Andrew Linux, Mac OS-X, ...) due to different versions of gcc and different compiler settings.

What it really looks like

```

        .globl  sumstore
        .type   sumstore, @function
sumstore:
.LFB35:
        .cfi_startproc
pushq   %rbx
        .cfi_def_cfa_offset 16
        .cfi_offset 3, -16
movq    %rdx, %rbx
call    plus
movq    %rax, (%rbx)
popq    %rbx
        .cfi_def_cfa_offset 8
ret
        .cfi_endproc
.LFE35:
        .size   sumstore, .-sumstore

```

Things that look weird
and are preceded by a ‘
are generally directives.

```

sumstore:
    pushq   %rbx
    movq    %rdx, %rbx
    call    plus
    movq    %rax, (%rbx)
    popq    %rbx
    ret

```

Assembly Characteristics: Data Types

- **“Integer” data of 1, 2, 4, or 8 bytes**
 - Data values
 - Addresses (untyped pointers)
- **Floating point data of 4, 8, or 10 bytes**
- **(SIMD vector data types of 8, 16, 32 or 64 bytes)**
- **Code: Byte sequences encoding series of instructions**
- **No aggregate types such as arrays or structures**
 - Just contiguously allocated bytes in memory

Assembly Characteristics: Operations

- **Transfer data between memory and register**
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- **Perform arithmetic function on register or memory data**

- **Transfer control**
 - Unconditional jumps to/from procedures
 - Conditional branches

Object Code

Code for `sumstore`

0x0400595:

0x53

0x48

0x89

0xd3

0xe8

0xf2

0xff

0xff

0xff

0x48

0x89

0x03

0x5b

0xc3

- **Total of 14 bytes**
- **Each instruction 1, 3, or 5 bytes**
- **Starts at address 0x0400595**

■ Assembler

- Translates `.s` into `.o`
- Binary encoding of each instruction
- Nearly-complete image of executable code
- Missing linkages between code in different files

■ Linker

- Resolves references between files
- Combines with static run-time libraries
 - E.g., code for `malloc`, `printf`
- Some libraries are *dynamically linked*
 - Linking occurs when program begins execution

Machine Instruction Example

```
*dest = t;
```

```
movq %rax, (%rbx)
```

```
0x40059e: 48 89 03
```

■ C Code

- Store value `t` where designated by `dest`

■ Assembly

- Move 8-byte value to memory
 - Quad words in x86-64 parlance
- Operands:
 - `t`: Register `%rax`
 - `dest`: Register `%rbx`
 - `*dest`: Memory `M[%rbx]`

■ Object Code

- 3-byte instruction
- Stored at address `0x40059e`

Disassembling Object Code

Disassembled

```

0000000000400595 <sumstore>:
 400595: 53                push   %rbx
 400596: 48 89 d3          mov    %rdx,%rbx
 400599: e8 f2 ff ff ff   callq 400590 <plus>
 40059e: 48 89 03          mov    %rax, (%rbx)
 4005a1: 5b                pop    %rbx
 4005a2: c3                retq

```

■ Disassembler

```
objdump -d sum
```

- Useful tool for examining object code
- Analyzes bit pattern of series of instructions
- Produces approximate rendition of assembly code
- Can be run on either a `.out` (complete executable) or `.o` file

Alternate Disassembly

Disassembled

```
Dump of assembler code for function sumstore:
0x0000000000400595 <+0>: push    %rbx
0x0000000000400596 <+1>: mov     %rdx,%rbx
0x0000000000400599 <+4>: callq  0x400590 <plus>
0x000000000040059e <+9>: mov     %rax, (%rbx)
0x00000000004005a1 <+12>: pop     %rbx
0x00000000004005a2 <+13>: retq
```

- **Within gdb Debugger**
 - Disassemble procedure
- ```
gdb sum
disassemble sumstore
```

# Alternate Disassembly

Object  
Code

```
0x0400595:
 0x53
 0x48
 0x89
 0xd3
 0xe8
 0xf2
 0xff
 0xff
 0xff
 0x48
 0x89
 0x03
 0x5b
 0xc3
```

Disassembled

```
Dump of assembler code for function sumstore:
0x000000000400595 <+0>: push %rbx
0x000000000400596 <+1>: mov %rdx,%rbx
0x000000000400599 <+4>: callq 0x400590 <plus>
0x00000000040059e <+9>: mov %rax, (%rbx)
0x0000000004005a1 <+12>: pop %rbx
0x0000000004005a2 <+13>: retq
```

## ■ Within gdb Debugger

- Disassemble procedure

```
gdb sum
```

```
disassemble sumstore
```

- Examine the 14 bytes starting at `sumstore`

```
x/14xb sumstore
```

# Machine Programming I: Summary

- **History of Intel processors and architectures**
  - Evolutionary design leads to many quirks and artifacts
- **C, assembly, machine code**
  - New forms of visible state: program counter, registers, ...
  - Compiler must transform statements, expressions, procedures into low-level instruction sequences
- **Assembly Basics: Registers, operands, move**
  - The x86-64 move instructions cover wide range of data movement forms
- **Arithmetic**
  - C compiler will figure out different instruction combinations to carry out computation