

# Symbolic Summation

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## Gosper's Algorithm

"The ultimate goal of mathematics is to eliminate  
all need for intelligent thought"

D. Knuth [ 2, p. 56]

### ■ Introduction

Consider

$$S(n) = \sum_{k=1}^n k k!$$

It satisfies to the following recurrence equation for  $S$

$$S(n) - S(n-1) = n n!$$

How to solve it? Using the method of iteration will lead back to the original sum. Fortunately, we can use a simple relation for a factorial function:

$$(n+1)! = (n+1)n! = n n! + n!$$

It follows

$$n n! = (n+1)! - n!$$

Therefore,

$$\sum_{k=1}^n k k! = \sum_{k=1}^n ((k+1)! - k!) =$$

$$(2! - 1!) + (3! - 2!) + (4! - 3!) + \dots + ((n+1)! - n!) = (n+1)! - 1$$

and we can find a closed form for that sum without solving any recurrence equations.

## ■ Main Idea

Given

$$\sum_{k=1}^n s_k$$

we want to find a new sequence  $z_k$  such that

$$s_k = z_{k+1} - z_k \quad (1)$$

If we succeed in finding  $z_k$  then the definite sum can be computed by telescoping

$$\sum_{k=1}^n s_k = \sum_{k=1}^n (z_{k+1} - z_k) = z_{n+1} - z_1$$

Thus, the problem of summation is reduced to finding sequence  $z_k$ . Let us consider the ratio  $\frac{z_k}{s_k}$  and assume that is a rational function:

$$\frac{z_k}{s_k} = \frac{z_k}{z_{k+1} - z_k} = \frac{1}{\frac{z_{k+1}}{z_k} - 1} \in Q(k) \quad (2)$$

This is so if  $z_k$  is a hypergeometric term. Since (2), we can write

$$z_k = y_k s_k$$

where  $y_k \in Q(k)$  is a unknown rational function. Substituting it back into (1)

$$s_k = z_{k+1} - z_k = y_{k+1} s_{k+1} - y_k s_k$$

and dividing it by  $s_k$

$$1 = y_{k+1} \frac{s_{k+1}}{s_k} - y_k$$

we obtain a recurrence equation for  $y_k$ , where  $\frac{s_{k+1}}{s_k} \in Q(k)$ . The latter must be proved. Indeed,

$$\frac{s_{k+1}}{s_k} = \frac{z_{k+2} - z_{k+1}}{z_{k+1} - z_k} = \frac{z_{k+1}}{z_k} \frac{\frac{z_{k+2}}{z_{k+1}} - 1}{\frac{z_{k+1}}{z_k} - 1} \in Q(k)$$

since  $z_k$  is a hypergeometric term. Therefore, we reduce summation problem to finding *rational solutions* to

$$a_k y_{k+1} - y_k = 1$$

where  $a_k = \frac{s_{k+1}}{s_k}$  is some known rational function. How would you find a rational solution? Wouldn't it be much easier to find a polynomial solution?

In the next step, we will show that we can simplify the problem further to finding only polynomial solutions. Assume (this will be proven later-!), we can rewrite any rational function in the following special form:

$$a_k = \frac{s_{k+1}}{s_k} = \frac{p_{k+1} q_{k+1}}{p_k r_{k+1}} \quad (3)$$

where  $p$ ,  $q$  and  $r$  are polynomials and

$$\text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

Next, we define  $f_k$  such that

$$f_k = \frac{p_{k+1}}{r_{k+1}} \frac{z_{k+1}}{s_{k+1}}$$

Clearly  $f_k$  is rational

$$f_k = \frac{p_{k+1}}{r_{k+1}} \frac{z_{k+1}}{s_{k+1}} = \frac{p_{k+1}}{r_{k+1}} \frac{z_{k+1}}{z_{k+2} - z_{k+1}} = \frac{p_{k+1}}{r_{k+1}} \frac{1}{\frac{z_{k+2}}{z_{k+1}} - 1} \in Q(k)$$

But we can say more about  $f_k$  - it is actually a polynomial. Again, this will be proven later. Let us continue with the main idea. Using  $f_k$  we can rewrite sequence  $z_k$  from (1) as

$$z_k = \frac{r_k}{p_k} s_k f_{k-1} \quad (4)$$

and therefore,

$$s_k = z_{k+1} - z_k = \frac{r_{k+1}}{p_{k+1}} s_{k+1} f_k - \frac{r_k}{p_k} s_k f_{k-1}$$

Dividing it by  $s_k$

$$1 = \frac{r_{k+1}}{p_{k+1}} \frac{s_{k+1}}{s_k} f_k - \frac{r_k}{p_k} f_{k-1}$$

multiplying it by  $p_k$

$$p_k = \frac{p_k}{p_{k+1}} \frac{s_{k+1}}{s_k} r_{k+1} f_k - r_k f_{k-1}$$

and using (3)

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1} q_{k+1}}{p_k r_{k+1}} \implies \frac{p_k}{p_{k+1}} \frac{s_{k+1}}{s_k} = \frac{q_{k+1}}{r_{k+1}}$$

arrive at the following recurrence equation with polynomial coefficients for  $f_k$

$$p_k = q_{k+1} f_k - r_k f_{k-1} \quad (5)$$

Remember, that  $f_k$  is a polynomial! So, we are interested only in polynomial solutions to (5). This can be easily done by using a generic polynomial (what degree-?) with unknown coefficients, substituting it into (5), equating all coefficients by  $k$  to zero to get a system of linear equations from which we could find coefficients of the original generic polynomial. Once we have  $f_k$ , we get  $z_k$

$$z_k = \frac{r_k}{p_k} s_k f_{k-1}$$

and finally using (1), we find the result of summation.

$$\sum_{k=1}^n s_k = \sum_{k=1}^n (z_{k+1} - z_k) = z_{n+1} - z_1$$

#### ■ We need to prove

1. representation (3) always exists
2.  $f_k$  is a polynomial
3. a priori bound on the degree of  $f_k$

Observe, the algorithm is invariant to the top limit of summation. Therefore, we say that Gosper's algorithm is designed for **indefinite summation**

$$\sum_k s_k = z_k$$

We define indefinite summation in the following way: given a sequence  $s_k$ , find another sequence  $z_k$  such that

$$s_k = z_{k+1} - z_k = \Delta z_k$$

Note, indefinite summation is the inverse of  $\Delta$ . Similar, as integration is the inverse of differentiation.

Knowing sequence  $z_k$  it is easy to compute a definite sum:

$$\sum_{k=m}^n s_k = \sum_{k=m}^n \Delta z_k = z_{n+1} - z_m$$

■ **Algorithm for indefinite summation**

To sum

$$\sum_{k=1}^n s_k \text{ or } \sum_k s_k$$

1. Compute

$$\frac{s_{k+1}}{s_k}$$

and verify that it is hypergeometric.

2. Find polynomials  $p$ ,  $q$ , and  $r$  such that

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}} \quad \text{where } \text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

3. Find a polynomial solution to

$$p_k = q_{k+1} f_k - r_k f_{k-1}$$

4. If such solution exists, return

$$z_k = \frac{r_k}{p_k} s_k f_{k-1}$$

otherwise terminate the algorithm.

■ **Example 1.**

$$\sum_{k=0}^n \frac{(-1)^k}{1+k} \binom{n}{k}$$

1. Compute

$$\frac{s_{k+1}}{s_k} = \frac{k-n}{2+k}$$

2. Choose  $p$ ,  $q$ , and  $r$ , such that

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}, \quad \text{where } \text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

Obviously,

$$\begin{aligned} p_k &= 1; \\ q_{k+1} &= k-n, \end{aligned}$$

$$r_{k+1} = 2 + k$$

3. Find a polynomial solution to  $p_k = q_{k+1} f_k - r_k f_{k-1}$  which is

$$1 = (k - n) f_k - (k + 1) f_{k-1}$$

The polynomial solution is a constant  $f_k = c$

$$1 = (k - n) c - (k + 1) c$$

$$1 = -n c - c$$

$$c = -\frac{1}{n + 1}$$

4. Return  $z_k = \frac{r_k}{p_k} s_k f_{k-1}$  which is

$$z_k = \frac{r_k}{p_k} s_k f_{k-1} = \frac{1+k}{1} * \frac{(-1)^k}{1+k} \binom{n}{k} * \frac{(-1)}{n+1} = \binom{n}{k} \frac{(-1)^{k+1}}{n+1}$$

Therefore,

$$\sum_{k=0}^n s_k = z_{n+1} - z_0 = \binom{n}{n+1} \frac{(-1)^{n+2}}{n+1} - \binom{n}{0} \frac{(-1)}{n+1} = \frac{1}{n+1}$$

### ■ Example 2.

If we generalize the denominator a little

$$\sum_{k=0}^n \frac{(-1)^k \binom{n}{k}}{k+z}$$

the sum is not computable by Gosper's algorithm. Prove it! Though, there is a closed form solution for a definite sum

$$\sum_{k=0}^n \frac{(-1)^k}{k+z} \binom{n}{k} = \frac{1}{z \binom{n+z}{z}}$$

### ■ Mathematica session

In this section we provide all steps of Gosper's algorithm the way they can be done in *Mathematica*.

As an example, we evaluate

$$\sum_k \frac{(-1)^k}{\binom{n}{k}}$$

First we define the summand as a *Mathematica* function

```
s[k_] := (-1)^k / Binomial[n, k]
```

Then we compute the ratio

```
s[k+1] / s[k] // FunctionExpand // Simplify
```

```
1+k
-----
k-n
```

Since it's a hypergeometric term, we find a triple  $(p_k, q_k, r_k)$  such that

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}$$

```
p[k_] := 1
q[k_] := k
r[k_] := k - 1 - n
```

Verify the polynomial GCD (see the algorithm for that down below)

```
Table[PolynomialGCD[q[k], r[k+j]], {j, 0, 15}]
```

```
{1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1}
```

Here is a difference equation for  $f_k$

$$p_k = q_{k+1} f_k - r_k f_{k-1}$$

```
p[k] == q[k+1] f[k] - r[k] f[k-1]
```

```
1 == -(-1+k-n) f[-1+k] + (1+k) f[k]
```

The solution is obvious

```
Collect [ - (-1 + k - n) f [-1 + k] + (1 + k) f [k] - 1 /. f [k_] := a, k,
Factor]
```

```
-1 + 2 a + a n
```

Solving it, we find  $a$

```
Solve [% == 0, a]
```

```
{{a -> 1 / (2 + n)}}
```

Thus,

$$z_k = \frac{r_k}{p_k} f_{k-1} s_k$$

and

$$\sum_k \frac{(-1)^k}{\binom{n}{k}} = \frac{(-1)^k (k - n - 1)}{(n + 2) \binom{n}{k}}$$

In particular,

$$\sum_{k=0}^n \frac{(-1)^k}{\binom{n}{k}} = \frac{(1 + (-1)^n) (n + 1)}{n + 2}$$

■  $\text{GCD}(q_k, r_{k+j}) = 1, j \in N_0$ .

If they have different degrees, we factor them over rationals and apply the procedure to each pair of factors. Therefore, we assume that  $q$  and  $r$  are of the same degree. Let

$$q_k = a k^n + b k^{n-1} + \dots$$

$$r_k = c k^n + d k^{n-1} + \dots$$

Consider

$$\frac{c}{a} q_k = c k^n + \frac{bc}{a} k^{n-1} + \dots \quad (6)$$

From the  $\text{GCD}(q_k, r_{k+j}) \neq 1$



$$\frac{c}{a} q_k = r_{k+j} \quad (7)$$

Evaluating the right hand side, we obtain

$$r_{k+j} = c(k+j)^n + d(k+j)^{n-1} + \dots = ck^n + ncjk^{n-1} + dk^{n-1} + \dots$$

Next, we put this and (6), back to equation (7)

$$ck^n + \frac{bc}{a} k^{n-1} + \dots = ck^n + ncjk^{n-1} + dk^{n-1} + \dots$$

Comparing coefficients by  $k^{n-1}$ , yields the value for  $j$

$$\frac{bc}{a} = ncj + d$$

$$j = \frac{bc - ad}{acn}$$

Note,  $j$  must be non-negative integer in order for the GCD not to be one. Finally, we have to verify (why-?)

$$cq_k - ar_{k+j} = 0$$

**Example.**

```

q[k_] := 5 k^2 + 2 k + 5
r[k_] := 5 k^2 - 48 k + 120
n := Exponent[q[k], k]
a := Coefficient[q[k], k, n]
b := Coefficient[q[k], k, n - 1]
c := Coefficient[r[k], k, n]
d := Coefficient[r[k], k, n - 1]
j = (bc - ad) / (acn)

```

```
5
```

```
c q[k] - a r[k + j] // Expand
```

```
0
```

## Gosper's Algorithm (proofs)

"In the first place, the beginner must be convinced that proofs deserve to be studied, that they have a purpose, that they are interesting" -

George Polya

### ■ We need to prove.

1. for any rational function

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}$$

where  $p$ ,  $q$  and  $r$  are polynomials. and

$$\text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

2.  $f_k$  is a polynomial

$$f_k = \frac{p_{k+1}}{r_{k+1}} \frac{z_{k+1}}{s_{k+1}}$$

3. a priori bound on the degree of  $f_k$

$$p_k = q_{k+1} f_k - r_k f_{k-1}$$

### ■ Proof of 1.

Let us choose

$$p_k = 1, \quad q_k = s_k, \quad \text{and} \quad r_k = s_{k-1}$$

$$\frac{s_{k+1}}{s_k} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}$$

If

$$\text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

there is nothing to prove. Assume that

$$\text{GCD}(q_k, r_{k+j}) = \pi_k \neq 1, \quad \text{for some } j \in J.$$

We choose new functions  $P_k$ ,  $Q_k$ , and  $R_k$  such that

$$\frac{P_{k+1}}{P_k} \frac{Q_{k+1}}{R_{k+1}} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}$$

They are

$$Q_k = \frac{q_k}{\pi_k}, \quad R_k = \frac{r_k}{\pi_{k-j}} \quad \text{and} \quad P_k = p_k \pi_k \pi_{k-1} \dots \pi_{k-j+1},$$

Simple check

$$\frac{P_{k+1}}{P_k} \frac{Q_{k+1}}{R_{k+1}} = \frac{p_{k+1} \pi_{k+1} \pi_k \dots \pi_{k-j+2}}{p_k \pi_k \pi_{k-1} \dots \pi_{k-j+2} \pi_{k-j+1}} \frac{q_{k+1} \pi_{k-j+1}}{\pi_{k+1} r_{k+1}} = \frac{p_{k+1}}{p_k} \frac{q_{k+1}}{r_{k+1}}$$

It follows then

$$\text{GCD}(Q_k, R_{k+j}) = \text{GCD}\left(\frac{q_k}{\pi_k}, \frac{r_{k+j}}{\pi_k}\right) = 1$$

Doing this for all  $j \in J$  we find triples  $(p_k, q_k, r_k)$  we conclude the proof. ■

### ■ Proof of 2.

The proof is by contradiction to  $\text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0$ . Assume that  $f_k$  is rational

$$f_k = \frac{a_k}{b_k}, \quad \text{gcd}(a_k, b_k) = 1$$

where  $b_k$  is NOT a constant. Substitute this  $f_k$  into the equation

$$p_k = q_{k+1} f_k - r_k f_{k-1}$$

we obtain

$$p_k = q_{k+1} \frac{a_k}{b_k} - r_k \frac{a_{k-1}}{b_{k-1}}$$

or (after multiplying by  $b_k b_{k-1}$ )

$$b_k b_{k-1} p_k = b_{k-1} q_{k+1} a_k - b_k r_k a_{k-1} \tag{8}$$

Suppose  $N$  is the largest integer such that (why the largest  $N$  exists -?)

$$\text{gcd}(b_k, b_{k+N}) = g_k \neq 1 \tag{9}$$

It follows from (9) that  $g_k$  divides  $b_k$ . Therefore,  $g_k$  divides  $b_{k-1} q_{k+1} a_k$ .

$g_k$  does not divide  $a_k$ , because  $\text{gcd}(a_k, b_k) = 1$

$g_k$  does not divide  $b_{k-1}$  by (9)

thus  $g_k$  divides  $q_{k+1}$  which is the same as  $g_{k-1}$  divides  $q_k$

It follows from (9) that  $g_{k-N-1}$  divides  $b_{k-1}$  (by performing simple shifting  $k \rightarrow k - N - 1$ ). Therefore, in view of (6)  $g_{k-N-1}$  divides  $b_k r_k a_{k-1}$

$g_{k-N-1}$  does not divide  $b_k$  by (9) ( $N$  is the maximum)

$g_{k-N-1}$  does not divide  $a_{k-1}$ , because otherwise  $g_k$  would divide  $a_{k+N}$  and then  $\gcd(a_{k+N}, b_{k+N}) \neq 1$

thus  $g_{k-N-1}$  divides  $r_k$  which is the same as  $g_{k-1}$  divides  $r_{k+N}$

This contradicts to

$$\text{GCD}(q_k, r_{k+j}) = 1, \quad j \in N_0.$$

for  $j = N$ . QED

### ■ Proof of 3.

Here we find the upper bound on the degree of  $f_k$

$$p_k = q_{k+1} f_k - r_k f_{k-1}$$

The equation can be rewritten as

$$p_k = (q_{k+1} - r_k) \frac{f_k + f_{k-1}}{2} + (q_{k+1} + r_k) \frac{f_k - f_{k-1}}{2} \quad (10)$$

It is easy to see that

$$\deg \left[ \frac{f_k + f_{k-1}}{2} \right] > \deg \left[ \frac{f_k - f_{k-1}}{2} \right]$$

because the highest degree term will be canceled in  $f_k - f_{k-1}$ .

Which one of two terms in the rhs of (10) has a higher degree?

#### Case 1

If

$$\deg(q_{k+1} - r_k) \geq \deg(q_{k+1} + r_k)$$

then

$$\deg(f_k) = \deg(p_k) - \deg(q_{k+1} - r_k)$$

*Proof.*

Since

$$\deg \left[ (q_{k+1} - r_k) \frac{f_k + f_{k-1}}{2} \right] > \deg \left[ (q_{k+1} + r_k) \frac{f_k - f_{k-1}}{2} \right]$$

the second term in the rhd of (10) can be neglected

$$p_k = (q_{k+1} - r_k) \frac{f_k + f_{k-1}}{2}$$

Therefore,

$$\deg(p_k) = \deg(q_{k+1} - r_k) + \deg\left(\frac{f_k + f_{k-1}}{2}\right) = \deg(q_{k+1} - r_k) + \deg(f_k)$$

This concludes the first case.

### Case 2

If  $n = \deg(q_{k+1} + r_k) > \deg(q_{k+1} - r_k)$

then

a) if  $\frac{2b}{a}$  is not negative integer

$$\deg(f_k) = \deg(p_k) - n + 1$$

b) if  $\frac{2b}{a}$  is negative integer

$$\deg(f_k) \leq \max\left[-2 * \frac{b}{a}, \deg(p_k) - n + 1\right]$$

*Proof* Let

$$f_k = c * k^m + \dots$$

$$q_{k+1} + r_k = a * k^n + \dots, \text{ where } n = \deg(q_{k+1} + r_k)$$

$$q_{k+1} - r_k = b * k^{n-1} + \dots \quad \text{why } (n-1)?$$

Then by (10) we have

$$p_k = (b * k^{n-1} + \dots) \frac{2 * c * k^m + \dots}{2} + (a * k^n + \dots) \frac{c * m * k^{m-1} + \dots}{2}$$

Collecting terms

$$p_k = \left(b + \frac{a * m}{2}\right) * c * k^{n+m-1} + \dots$$

Depending on whether or not  $\frac{2b}{a}$  is negative, we have two choices. QED

## Not Gosper-summable

$$\sum_{k=0}^m \binom{n}{k}$$

Compute the ratio

$$\frac{s_{k+1}}{s_k} = \frac{n-k}{k+1}$$

Obviously

$$\begin{aligned} p_k &= 1; \\ q_{k+1} &= n-k, \\ r_{k+1} &= k+1 \end{aligned}$$

Compute the degree of  $f_k$

$$\deg(q_{k+1} - r_k) = \deg(n - k - k) = 1$$

$$\deg(q_{k+1} + r_k) = \deg(n - k + k) = 0$$

This is the first case, therefore

$$\deg(f_k) = \deg(p_k) - \deg(q_{k+1} - r_k)$$

$$\deg(f_k) = -1$$

but  $f_k$  must be a polynomial, thus, the above sum is not Gosper-summable.

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