Constructive Logic (15-317), Fall 2012 Assignment 4: Sequent Calculus and Natural Numbers

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Out: Thursday, September 27, 2012 Due: Thursday, October 4, 2012 (before class)

In this assignment, you will prove metatheorems about sequent calculus, and propositions about natural numbers.

1 Contraction (20 points)

The contraction lemma for sequent calculus states that two equal propositions on the left of a sequent can be replaced by a single instance:

Lemma 1 (Contraction). If Δ , A, $A \rightarrow D$, then Δ , $A \rightarrow D$.

We can prove contraction by structural induction on the derivation Δ , A, $A \rightarrow D$, in a similar fashion to the proofs of the substitution and weakening lemmas for natural deduction.

Task 1 (20 points). Prove contraction for the \supset , \land fragment of sequent calculus. That is, show the cases of the proof where the last rule to be applied is:

- init
- $\supset R$
- $\supset L$
- $\wedge R$
- $\wedge L_1$

You don't need to show the $\wedge L_2$ case, because it is very similar to the $\wedge L_1$ case.

2 Derivability and Admissibility (10 points)

Task 2 (10 points). For each of the following sequent calculus rules, state whether or not it is derivable, and whether or not it is admissible. Explain.

$$\frac{\Delta, A, A \longrightarrow D}{\Delta, A \longrightarrow D} (1) \qquad \frac{A \text{ atomic}}{\Delta, A \land B \longrightarrow A} (2)$$

$$\frac{\Delta, A, A \supset B \longrightarrow B}{\Delta, A, A \supset C} (3) \qquad \frac{\Delta \longrightarrow A \quad \Delta, A \longrightarrow C}{\Delta \longrightarrow C} (4)$$

3 Natural numbers (10 points)

In recitation, we saw how to define the natural numbers as a data type, and prove theorems quantified over them. We also defined the even and odd predicates on natural numbers:

	n:nat	odd(n) true	n:nat	even(<i>n</i>) <i>true</i>
even(0)	even(s(n)) true		odd(s(<i>n</i>)) <i>true</i>	

Task 3 (5 points). Derive the following judgment:

 $\forall n : nat.(even(n) \supset even(s(s(n)))) true$

Task 4 (5 points). Derive the following judgment:

 $\forall n : nat.(even(n) \lor even(s(n))) true$