

## 1. Isomorphisms

Show that there are eleven nonisomorphic simple graphs on four vertices.

**Solution:**

0            1            5            6

2 conn.    2 non-conn.    4 conn.    4 non-conn.

3 cycle    3 tree A    3 tree B

## 2. Euler's Formula

A soccer ball is a convex polyhedron whose faces are either hexagons or pentagons. Prove that a soccer ball has exactly 12 pentagonal faces.

**Solution:** Let the total number of faces on the soccer ball be  $F$ , the number of pentagonal faces be  $f_5$ , and the number of hexagonal faces be  $f_6$ . Consequently,  $F = f_5 + f_6$ .

First, note that every edge on the soccer ball is adjacent to two faces. Then we have the equation

$$2E = 5f_5 + 6f_6$$

which can be plugged into Euler's formula like so:

$$\begin{aligned} V - E + F - 2 &= V - \frac{5f_5 + 6f_6}{2} + f_5 + f_6 - 2 \\ &= V - \frac{3f_5 + 4f_6}{2} - 2 \end{aligned}$$

Next, note that the sum of the degrees of the angles at a vertex on a convex polyhedron must be less than  $360^\circ$ . Then there must be exactly three faces that meet at each vertex ( $108^\circ$  per pentagon,  $120^\circ$  per hexagon). This gives us the equation

$$3V = 5f_5 + 6f_6$$

which can be combined with our previous equation to find the following:

$$\begin{aligned} 3 \left( \frac{3f_5 + 4f_6}{2} + 2 \right) &= 5f_5 + 6f_6 \\ f_5 &= 12 \end{aligned}$$

### 3. Graph Induction

A tournament of  $n$  vertices is a complete directed graph with  $n$  vertices. Prove that there always exists a directed path that visits each vertex in the tournament exactly once.

**Solution:** Induction on the number of vertices in the tournament. If there are two vertices, there is exactly one edge and the directed path is immediately obvious. Consider the case of  $k$  vertices. Remove a vertex and suppose any tournament of  $k - 1$  vertices has a directed Hamiltonian path. If there is an edge pointing from the removed vertex to the tail of the  $k - 1$  path, we are done. Similarly, if there is an edge pointing from the head of the  $k - 1$  path to the removed vertex, we are done. Otherwise, step through the directed path until we find the first vertex  $i$  such that there is an edge pointing from the removed vertex to vertex  $i$ . Every vertex in the path before  $i$  must have an edge pointing to the removed vertex, so we can create a new directed path from the tail to the removed vertex to the  $i$ th vertex to the head.