

Algebra II: Fields, Polynomials

Recap: Definition of a group

G is a "group under operation \bullet " if:

- [Closure] G is closed under \bullet
 i.e., $a \bullet b \in G \quad \forall a, b \in G$
- [Associativity] Operation \bullet is *associative*:
 i.e., $a \bullet (b \bullet c) = (a \bullet b) \bullet c \quad \forall a, b, c \in G$
- [Identity] There exists an element $e \in G$
 (called the "identity element") such that
 $a \bullet e = a, e \bullet a = a \quad \forall a \in G$
- [Inverse] For each $a \in G$ there is an element $a^{-1} \in G$
 (called the "inverse of a ") such that
 $a \bullet a^{-1} = e, a^{-1} \bullet a = e$

Examples of groups

Any group of transformations is a group.

E.g., the 'mattress group' (AKA Klein 4-group)

\bullet	Id	R	F	H
Id	Id	R	F	H
R	R	Id	H	F
F	F	H	Id	R
H	H	F	R	Id

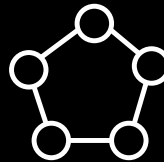
identity element is Id

$$R^{-1} = R$$

$$F^{-1} = F$$

$$H^{-1} = H$$

Symmetries of undirected cycle: dihedral group



$$G = \{ \text{Id}, r_1, r_2, r_3, r_4, f_1, f_2, f_3, f_4, f_5 \}$$

\bullet	Id	r_1	r_2	r_3	r_4	f_1	f_2	f_3	f_4	f_5
Id	Id	r_1	r_2	r_3	r_4	f_1	f_2	f_3	f_4	f_5
r_1	r_1	Id	r_4	r_2	r_3	f_4	f_3	f_2	f_1	f_5
r_2	r_2	r_4	Id	r_1	r_3	f_3	f_4	f_1	f_5	f_2
r_3	r_3	r_4	Id	r_1	r_2	f_2	f_4	f_5	f_3	f_1
r_4	r_4	Id	r_1	r_2	r_3	f_1	f_3	f_4	f_2	f_5
f_1	f_1	f_2	f_3	f_4	f_5	Id	r_3	r_1	r_4	r_2
f_2	f_2	Id	r_4	r_2	r_3	r_1	Id	r_3	r_1	r_4
f_3	f_3	r_4	Id	r_1	r_2	r_3	r_1	Id	r_3	r_2
f_4	f_4	r_3	r_4	Id	r_1	r_2	r_4	r_2	Id	r_3
f_5	f_5	r_2	r_3	r_4	Id	r_3	r_4	r_2	r_3	Id

Symmetries of directed cycle: Cyclic group

C_4

\bullet	Id	R_1	R_2	R_3
Id	Id	R_1	R_2	R_3
R_1	R_1	Id	R_3	R_2
R_2	R_2	R_3	Id	R_1
R_3	R_3	Id	R_1	R_2

In a group table, every row and every column is a permutation of the group elements
 Follows because (i) $a \bullet b = a \bullet c \Rightarrow b = c$
 (ii) $b \bullet a = c \bullet a \Rightarrow b = c$

Abelian groups

In a group we do NOT NECESSARILY have

$$a \bullet b = b \bullet a$$

Definition:

" $a, b \in G$ commute" means $ab = ba$.

Definition:

A group is said to be *abelian* if all pairs $a, b \in G$ commute.

Order of a group element

Let G be a *finite* group. Let $a \in G$.

Definition: The **order** of x , denoted $\text{ord}(a)$, is the smallest $m \geq 1$ such that $a^m = 1$.

Note that $a, a^2, a^3, \dots, a^{m-1}, a^m=1$ all distinct.

Order Theorem: For every $a \in G$, $\text{ord}(a)$ divides $|G|$.

Corollary: $a^{|G|} = 1$ for all $a \in G$.

Corollary (Euler's Theorem): For $a \in \mathbb{Z}_n^*$, $a^{\phi(n)} = 1$
That is, if $\gcd(a,n)=1$, then $a^{\phi(n)} \equiv 1 \pmod{n}$

Corollary (Fermat's little theorem):
For prime p , if $\gcd(a,p)=1$, then
 $a^{p-1} \equiv 1 \pmod{p}$

Cyclic groups

A finite group G of order n is cyclic if $G = \{e, b, b^2, \dots, b^{n-1}\}$ for some group element b

In such a case, we say b "**generates**" G , or b is a "**generator**" of G .

Examples:

- $(\mathbb{Z}_n, +)$ (1 is a generator)
- C_4 (Rot_{90} is a generator)

Non-examples: Matrix group;
any non-abelian group.

How many generators does $(\mathbb{Z}_n, +)$ have?

Answer: $\phi(n)$

b generates $\mathbb{Z}_n \Leftrightarrow \exists a$ s.t. $ba \equiv 1 \pmod{n}$
($ba = b+b+\dots+b$ (a times))

Same holds for *any* cyclic group with n elements

Subgroups

Q: Is (Even integers, $+$) a group?

A: Yes. It is a "subgroup" of $(\mathbb{Z}, +)$

Definition: Suppose (G, \bullet) is a group.

If $H \subseteq G$, and if (H, \bullet) is also a group, then H is called a **subgroup** of G .

To check H is a subgroup of G , check:

1. H is closed under \bullet
2. $e \in H$
3. If $h \in H$ then $h^{-1} \in H$
 - (3rd condition follows from 1,2 if H is finite)

Examples

Every G has two trivial subgroups: $\{e\}, G$
Rest are called "proper" subgroups

Suppose $k, 1 < k < n$, divides n .

Q1. Is $(\{0, k, 2k, 3k, \dots, (n/k-1)k\}, +_n)$ subgroup of $(\mathbb{Z}_n, +_n)$?
Yes!

Q2. Is $(\mathbb{Z}_{kr}, +_k)$ a subgroup of $(\mathbb{Z}_n, +_n)$?
No! it doesn't even have the same operation

Q3. Is $(\mathbb{Z}_{kr}, +_n)$ a subgroup of $(\mathbb{Z}_n, +_n)$?
No! \mathbb{Z}_k is not closed under $+_n$

Lagrange's Theorem

Theorem: If G is a finite group, and H is a subgroup then $|H|$ divides $|G|$.

Proof similar to order theorem.

Corollary (order theorem): If $x \in G$, then $\text{ord}(x)$ divides $|G|$.

Proof of Corollary:

Consider the set $T_x = \{x, x^2, x^3, \dots\}$

(i) $\text{ord}(x) = |T_x|$

(ii) (T_x, \bullet) is a subgroup of (G, \bullet) (check!)

A useful corollary

If G is a finite group and H is a proper subgroup of G , then $|H| \leq |G|/2$

Example:

$G = (\{0,1\}^n, +)$ ($+$: coordinate-wise addition mod 2)

$H = \{(x_1, x_2, \dots, x_n) : x_1 + x_2 + \dots + x_n = 0\}$

H is a proper subgroup (check!)

So $|H| \leq 2^{n-1}$

(in fact, in this case, $|H| = 2^{n-1}$)

It was nice meeting you, groups!

Number Theory Interlude II

Bezout's identity

Let a, b be arbitrary positive integers. There exist integers r and s such that

$$r a + s b = \gcd(a, b)$$

Follows from
Extended
Euclid Algorithm

A non-algorithmic proof:

- Consider the set L of all *positive* integers that can be expressed as $r a + s b$ for some integers r, s .
- L is non-empty (eg. $a \in S$)
- So L has a minimum element d
(well-ordering principle \Leftrightarrow principle of induction)

Claim: $d = \gcd(a, b)$

Claim: $\gcd(a, b) = d$ (the minimum positive integer expressible as $ra+sb$)

1. $\gcd(a, b)$ divides both a and b , and hence also divides d . So $d \geq \gcd(a, b)$
2. d divides both a and b , and hence $d \leq \gcd(a, b)$

Let's show $d \mid a$.

Write $a = q d + t$, with $0 \leq t < d$.

$t = a - q d$ is also expressible as a combination $r' a + s' b$.

Contradicts minimality of d .

Extended Euclid & Unique Factorization

Lemma: If $\gcd(a, b) = 1$ and $a \mid bc$, then $a \mid c$.

Proof: Let r, s be such that $r a + s b = 1$

$$r a c + s b c = c$$

$a \mid bc$ and $a \mid r a c$, so $a \mid c$. \square

Corollary: If p is a prime and $p \mid q_1 q_2 \dots q_k$, then p must divide some q_i .

If the q_i 's are also prime, then $p = q_i$ for some i .

Unique prime factorization follows from this!

Extended Euclid and Chinese Remaindering

Chinese Remainder Theorem: Suppose n_1, n_2, \dots, n_k are pairwise coprime. Then, for all integers a_1, a_2, \dots, a_k , there exists an integer x solving the below system of simultaneous congruences

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\vdots \\ x &\equiv a_k \pmod{n_k}. \end{aligned}$$

Further, all solutions x are congruent modulo $N = \prod_{i=1}^k n_i$.

Uniqueness of solutions modulo N

If x, y are two solutions, then n_i divides $x-y$, for $i=1, 2, \dots, k$

Since the n_i are coprime, this means the product $N = n_1 n_2 \dots n_k$ divides $(x-y)$, thus $x \equiv y \pmod{N}$

Extended Euclid and Chinese Remaindering

Chinese Remainder Theorem: Suppose n_1, n_2, \dots, n_k are pairwise coprime. Then, for all integers a_1, a_2, \dots, a_k , there exists an integer x solving the below system of simultaneous congruences

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Further, all solutions x are congruent modulo $N = \prod_{i=1}^k n_i$.

Proof for $k=2$:

Take $x = a_1 (n_2^{-1} \pmod{n_1}) n_2 + a_2 (n_1^{-1} \pmod{n_2}) n_1$

Can compute x efficiently (by computing modular inverses)

Extended Euclid and Chinese Remaindering

Chinese Remainder Theorem: Suppose n_1, n_2, \dots, n_k are pairwise coprime. Then, for all integers a_1, a_2, \dots, a_k , there exists an integer x solving the below system of simultaneous congruences

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\vdots \\ x &\equiv a_k \pmod{n_k}. \end{aligned}$$

Further, all solutions x are congruent modulo $N = \prod_{i=1}^k n_i$.

For arbitrary k : Let $m_i = N/n_i$ Note $\gcd(m_i, n_i) = 1$
 $n_i \mid m_j$ for $j \neq i$

Take $x = a_1 (m_1^{-1} \pmod{n_1}) m_1 + a_2 (m_2^{-1} \pmod{n_2}) m_2 + \dots + a_k (m_k^{-1} \pmod{n_k}) m_k$

Field Theory

Find out about the wonderful world of \mathbb{F}_{2^k}
 where two equals zero, plus is minus,
 and squaring is a linear operator!

– Richard Schroepel



A group is a set with a single binary operation.

Number-theoretic sets often have more than one operation defined on them.

For example, in \mathbb{Z} , we can do both addition and multiplication.

Same in \mathbb{Z}_n (we can add and multiply modulo n)

For reals \mathbb{R} or rationals \mathbb{Q} , we can also divide (inverse operation for multiplication).

Fields

Informally, it's a place where you can add, subtract, multiply, and divide.

Examples: Real numbers \mathbb{R}
 Rational numbers \mathbb{Q}
 Complex numbers \mathbb{C}
 Integers mod *prime* \mathbb{Z}_p (Why?)

NON-examples: Integers \mathbb{Z} division??
 Non-negative reals \mathbb{R}^+ subtraction??

Field – formal definition

A **field** is a set F with two binary operations, called $+$ and \cdot .

$(F, +)$ an abelian group, with identity element called 0

$(F \setminus \{0\}, \cdot)$ an abelian group, identity element called 1

Distributive Law holds:
 $a \cdot (b+c) = a \cdot b + a \cdot c$

Example:
 $F_3 = \mathbb{Z}_3^*$

+	0	1	2
0	0	1	2
1	1	2	0
2	2	0	1

•	0	1	2
0	0	0	0
1	0	1	2
2	0	2	1

Example

Quadratic “number field”

$$\mathbb{Q}(\sqrt{2}) = \{ a + b\sqrt{2} : a, b \in \mathbb{Q} \}$$

Addition: $(a + b\sqrt{2}) + (c + d\sqrt{2}) = (a+c) + (b+d)\sqrt{2}$

Multiplication:
 $(a + b\sqrt{2}) \cdot (c + d\sqrt{2}) = (ac+2bd) + (ad+bc)\sqrt{2}$

Exercise: Prove above defines a field.

Finite fields

Some familiar *infinite* fields: $\mathbb{Q}, \mathbb{R}, \mathbb{C}$ (now $\mathbb{Q}(\sqrt{2})$)

Finite fields we know: \mathbb{Z}_p aka F_p for p a prime

Is there a field with 2 elements? **Yes**

Is there a field with 3 elements? **Yes**

Is there a field with 4 elements? **Yes**

F_4

+	0	1	a	b
0	0	1	a	b
1	1	0	b	a
a	a	b	0	1
b	b	a	1	0

•	0	1	a	b
0	0	0	0	0
1	0	1	a	b
a	0	a	b	1
b	0	b	1	a

Finite fields

Is there a field with 2 elements? **Yes**

Is there a field with 3 elements? **Yes**

Is there a field with 4 elements? **Yes**

Is there a field with 5 elements? **Yes**

Is there a field with 6 elements? **No**

Is there a field with 7 elements? **Yes**

Is there a field with 8 elements? **Yes**

Is there a field with 9 elements? **Yes**

Is there a field with 10 elements? **No**



Evariste **Galois** (1811–1832) introduced the concept of a finite field (also known as a **Galois Field** in his honor)

Finite fields

Theorem (which we won't prove):

There is a field with q elements if and only if q is a **power of a prime**.

Up to *isomorphism*, it is **unique**.

That is, all fields with q elements have the same addition and multiplication tables, after renaming elements.

This field is denoted \mathbb{F}_q (also $\text{GF}(q)$)

Finite fields

Question:

If q is a prime power but not just a prime, what **are** the addition and multiplication tables of \mathbb{F}_q ?

Answer:

It's a bit hard to describe.

We'll tell you later, but for 251's purposes, you only need to know about prime q .

Polynomials

Polynomials

Informally, a polynomial is an expression that looks like this:

$$6x^3 - 2.3x^2 + 5x + 4.1$$

x is a symbol, called the *variable*

the 'numbers' standing next to powers of x are called the *coefficients*

Polynomials

Informally, a polynomial is an expression that looks like this:

$$6x^3 - 2.3x^2 + 5x + 4.1$$

$\in \mathbb{R}[x]$

Actually, coefficients can come from any *field*.

Can allow multiple variables, but we won't.

Set of polynomials with variable x and coefficients from field F is denoted $\mathbb{F}[x]$.

Polynomials – formal definition

Let F be a field and let x be a variable symbol.

$F[x]$ is the set of **polynomials over F** , defined to be expressions of the form

$$c_d x^d + c_{d-1} x^{d-1} + \dots + c_2 x^2 + c_1 x + c_0$$

where each c_i is in F , and $c_d \neq 0$.

We call d the **degree** of the polynomial.

Also, the expression 0 is a polynomial.

(By convention, we call its degree $-\infty$.)

Adding and multiplying polynomials

You can add and multiply polynomials.

Example. Here are two polynomials in $F_{11}[x]$

$$P(x) = x^2 + 5x - 1$$

$$Q(x) = 3x^3 + 10x$$

$$\begin{aligned} P(x) + Q(x) &= 3x^3 + x^2 + 15x - 1 \\ &= 3x^3 + x^2 + 4x - 1 \\ &= 3x^3 + x^2 + 4x + 10 \end{aligned}$$

Adding and multiplying polynomials

You can add and multiply polynomials (they are a “ring” but we’ll skip a formal treatment of rings)

Example. Here are two polynomials in $F_{11}[x]$

$$P(x) = x^2 + 5x - 1$$

$$Q(x) = 3x^3 + 10x$$

$$\begin{aligned} P(x) \cdot Q(x) &= (x^2 + 5x - 1)(3x^3 + 10x) \\ &= 3x^5 + 15x^4 + 7x^3 + 50x^2 - 10x \\ &= 3x^5 + 4x^4 + 7x^3 + 6x^2 + x \end{aligned}$$

Adding and multiplying polynomials

Polynomial addition is associative and commutative.

$$0 + P(x) = P(x) + 0 = P(x).$$

$$P(x) + (-P(x)) = 0.$$

So $(F[x], +)$ is an abelian group!

Polynomial multiplication is associative and commutative.

$$1 \cdot P(x) = P(x) \cdot 1 = P(x).$$

Multiplication distributes over addition:

$$P(x) \cdot (Q(x) + R(x)) = P(x) \cdot Q(x) + P(x) \cdot R(x)$$

If $P(x) / Q(x)$ were always a polynomial, then $F[x]$ would be a field! *Alas...*

Dividing polynomials?

$P(x) / Q(x)$ is not necessarily a polynomial.

So $F[x]$ is not quite a field.

(It’s an “integral domain”)

Same with \mathbb{Z} , the integers:

it has everything except division.

Actually, there are many analogies between $F[x]$ and \mathbb{Z} .

- starting point for rich interplay between algebra, arithmetic, and geometry in modern mathematics

Dividing polynomials?

\mathbb{Z} has the concept of “division with remainder”:

Given $a, b \in \mathbb{Z}$, $b \neq 0$, can write

$$a = q \cdot b + r,$$

where r is “smaller than” b .

$F[x]$ has the same concept:

Given $A(x), B(x) \in F[x]$, $B(x) \neq 0$, can write

$$A(x) = Q(x) \cdot B(x) + R(x),$$

where $\deg(R(x)) < \deg(B(x))$.

“Division with remainder” for polynomials

Example: Divide $6x^4 + 8x + 1$ by $2x^2 + 4$ in $F_{11}[x]$

$$\begin{array}{r} 3x^2 + 5 \\ 2x^2 + 4 \overline{) 6x^4 + 8x + 1} \\ \underline{- 6x^4 + x^2} \\ -x^2 + 8x + 1 \\ \underline{- -x^2 + 9} \\ 8x + 3 \end{array}$$

Check:

$$\begin{aligned} &6x^4 + 8x + 1 \\ &= (3x^2 + 5)(2x^2 + 4) + (8x + 3) \\ &\quad (\text{in } F_{11}[x]) \end{aligned}$$

Integers \mathbb{Z}

"size" = abs. value

"division":

$$a = qb + r, \quad |r| < |b|$$

can use **Euclid's Algorithm** to find GCDs

p is "prime":

no nontrivial divisors

$\mathbb{Z} \bmod p$:

a field iff p is prime

Polynomials $F[x]$

"size" = degree

"division":

$$A(x) = Q(x)B(x) + R(x), \\ \deg(R) < \deg(B)$$

can use **Euclid's Algorithm** to find GCDs

$P(x)$ is "irreducible":

no nontrivial divisors

$F[x] \bmod P(x)$:

a field iff $P(x)$ is irreducible
(with $|F|^{\deg(P)}$ elements)

The field with 4 elements

Degree < 2 polynomials $\{0, 1, x, 1+x\} \subseteq F_2[x]$

Addition and multiplication modulo $1+x+x^2$

F_4

$a=x$
 $b=1+x$

+	0	1	a	b
0	0	1	a	b
1	1	0	b	a
a	a	b	0	1
b	b	a	1	0

•	0	1	a	b
0	0	0	0	0
1	0	1	a	b
a	0	a	b	1
b	0	b	1	a

Enough algebraic theory.

Let's play with polynomials!

Evaluating polynomials

Given a polynomial $P(x) \in F[x]$,
 $P(a)$ means its **evaluation** at element a .

E.g., if $P(x) = x^2 + 3x + 5$ in $F_{11}[x]$

$$P(6) = 6^2 + 3 \cdot 6 + 5 = 36 + 18 + 5 = 59 = 4$$

$$P(4) = 4^2 + 3 \cdot 4 + 5 = 16 + 12 + 5 = 33 = 0$$

Definition: α is a **root** of $P(x)$ if $P(\alpha) = 0$.

Polynomial roots

Theorem:

Let $P(x) \in F[x]$ have degree **1**.
Then $P(x)$ has exactly **1** root.

Proof:

Write $P(x) = cx + d$ (where $c \neq 0$).

$$\begin{aligned} \text{Then } P(r) = 0 &\Leftrightarrow cr + d = 0 \\ &\Leftrightarrow cr = -d \\ &\Leftrightarrow r = -d/c. \end{aligned}$$

Polynomial roots

Theorem:

Let $P(x) \in F[x]$ have degree **2**.
Then $P(x)$ has... how many roots??

E.g.: $x^2 + 1 \dots$

of roots over $F_2[x]$: **1** (namely, 1)

of roots over $F_3[x]$: **0**

of roots over $F_5[x]$: **2** (namely, 2 and 3)

of roots over $\mathbb{R}[x]$: **0**

of roots over $\mathbb{C}[x]$: **2** (namely, i and $-i$)

The single most important theorem about polynomials over fields:

A nonzero degree- d polynomial has at most d roots.

Theorem: Over a field, for all $d \geq 0$, a nonzero degree- d polynomial P has at most d roots.

Proof by induction on $d \in \mathbb{N}$:

Base case: If $P(x)$ is degree-0 then $P(x) = a$ for some $a \neq 0$. This has 0 roots.

Recall our convention: $\deg(0) = -\infty$

Induction:

Assume true for $d \geq 0$. Let $P(x)$ have degree $d+1$.

If $P(x)$ has 0 roots: we're done! Else let b be a root.

Divide with remainder: $P(x) = Q(x)(x-b) + R(x)$. (*)

$\deg(R) < \deg(x-b) = 1$, so $R(x)$ is a constant. Say $R(x) = r$.

Plug $x = b$ into (*): $0 = P(b) = Q(b)(b-b) + r = 0 + r = r$.

So $P(x) = Q(x)(x-b)$. Now, $\deg(Q) = d$. $\therefore Q$ has $\leq d$ roots.

$\therefore P(x)$ has $\leq d+1$ roots, completing the induction.

A useful corollary

Theorem: Over a field F , for all $d \geq 0$, degree- d polynomials have at most d roots.

Corollary: Suppose a polynomial $R(x) \in F[x]$ is such that

- (i) R has degree $\leq d$ and
- (ii) R has $> d$ roots

Then R must be the 0 polynomial

I've used the above corollary *several times* in my research.

Theorem:

Over a field, degree- d polynomials have at most d roots.

Reminder:

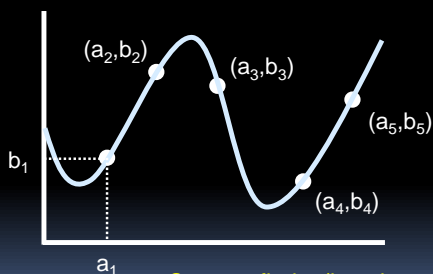
This is only true over a field.

E.g., consider $P(x) = 3x$ over \mathbb{Z}_6 .

It has degree 1, but 3 roots: 0, 2, and 4.

Interpolation

Say you're given a bunch of "data points"



Can you find a (low-degree) polynomial which "fits the data"?

Interpolation

Let pairs $(a_1, b_1), (a_2, b_2), \dots, (a_{d+1}, b_{d+1})$ from a field F be given (with all a_i 's distinct).

Theorem:

There is **exactly one** polynomial $P(x)$ of **degree at most d** such that $P(a_i) = b_i$ for all $i = 1 \dots d+1$.

E.g., through 2 points there is a unique linear polynomial.

Interpolation

There are two things to prove.

1. There is at *least* one polynomial of degree $\leq d$ passing through all $d+1$ data points.
2. There is at *most* one polynomial of degree $\leq d$ passing through all $d+1$ data points.

Let's prove #2 first.

Interpolation

Theorem: Let pairs $(a_1, b_1), (a_2, b_2), \dots, (a_{d+1}, b_{d+1})$ from a field F be given (with all a_i 's distinct).

Then there is **at most one** polynomial $P(x)$ of **degree at most d** with $P(a_i) = b_i$ for all i .

Proof: Suppose $P(x)$ and $Q(x)$ both do the job.

Let $R(x) = P(x) - Q(x)$.

Since $\deg(P), \deg(Q) \leq d$ we must have $\deg(R) \leq d$.

But $R(a_i) = b_i - b_i = 0$ for all $i = 1 \dots d+1$.

Thus $R(x)$ has more roots than its degree.

$\therefore R(x)$ must be the 0 polynomial, i.e., $P(x) = Q(x)$.

Interpolation

Now let's prove the other part, that there is **at least one** polynomial.

Theorem:

Let pairs $(a_1, b_1), (a_2, b_2), \dots, (a_{d+1}, b_{d+1})$ from a field F be given (with all a_i 's distinct).

Then there **exists** a polynomial $P(x)$ of **degree at most d** with $P(a_i) = b_i$ for all i .

Interpolation

The method for constructing the polynomial is called **Lagrange Interpolation**.

Discovered in 1779
by Edward Waring.



Rediscovered in 1795
by J.-L. Lagrange.



Lagrange Interpolation

a_1	b_1
a_2	b_2
a_3	b_3
\dots	\dots
a_d	b_d
a_{d+1}	b_{d+1}

Want $P(x)$
(with $\text{degree} \leq d$)
such that $P(a_i) = b_i \forall i$.

Lagrange Interpolation

a_1	1
a_2	0
a_3	0
\dots	\dots
a_d	0
a_{d+1}	0

Can we do this special case?

Promise: once we solve this special case, the general case is very easy.

Lagrange Interpolation

a_1	1
a_2	0
a_3	0
...	...
a_d	0
a_{d+1}	0



Lagrange Interpolation

a_1	1
a_2	0
a_3	0
...	...
a_d	0
a_{d+1}	0

Just divide $P(x)$ by this number.

Idea #1: $P(x) = (x-a_2)(x-a_3)\cdots(x-a_{d+1})$

Degree is d . ✓

$P(a_2) = P(a_3) = \cdots = P(a_{d+1}) = 0$. ✓

$P(a_1) = (a_1-a_2)(a_1-a_3)\cdots(a_1-a_{d+1})$. ??

Lagrange Interpolation

Numerator is a deg. d polynomial

a_1	1
a_2	0
a_3	0
...	...
a_d	0
a_{d+1}	0

Denominator is a nonzero field element

Idea #2:

$$S_1(x) = \frac{(x-a_2)(x-a_3)\cdots(x-a_{d+1})}{(a_1-a_2)(a_1-a_3)\cdots(a_1-a_{d+1})}$$

Call this the **selector polynomial** for a_1 .

Lagrange Interpolation

a_1	0
a_2	1
a_3	0
...	...
a_d	0
a_{d+1}	0

Great! But what about **this** data?

$$S_2(x) = \frac{(x-a_1)(x-a_3)\cdots(x-a_{d+1})}{(a_2-a_1)(a_2-a_3)\cdots(a_2-a_{d+1})}$$

Lagrange Interpolation

a_1	0
a_2	0
a_3	0
...	...
a_d	0
a_{d+1}	1

Great! But what about **this** data?

$$S_{d+1}(x) = \frac{(x-a_1)(x-a_2)\cdots(x-a_d)}{(a_{d+1}-a_1)(a_{d+1}-a_2)\cdots(a_{d+1}-a_d)}$$

Lagrange Interpolation

a_1	b_1
a_2	b_2
a_3	b_3
...	...
a_d	b_d
a_{d+1}	b_{d+1}

Great! But what about **this** data?

$$P(x) = b_1 \cdot S_1(x) + b_2 \cdot S_2(x) + \cdots + b_{d+1} \cdot S_{d+1}(x)$$

Lagrange Interpolation – example

Over \mathbb{Z}_{11} , find a polynomial P of degree ≤ 2 such that $P(5) = 1$, $P(6) = 2$, $P(7) = 9$.

$$S_5(x) = \frac{1}{(5-6)(5-7)}(x-6)(x-7)$$

$$S_6(x) = \frac{1}{(6-5)(6-7)}(x-5)(x-7)$$

$$S_7(x) = \frac{1}{(7-5)(7-6)}(x-5)(x-6)$$

$$P(x) = 1 S_5(x) + 2 S_6(x) + 9 S_7(x)$$

$$P(x) = 6(x^2-13x+42) - 2(x^2-12x+35) + 54(x^2-11x+30)$$

$$P(x) = 3x^2+x+9$$

The Chinese Remainder Theorem had a very similar proof



Not a coincidence:

algebraically, integers & polynomials share many common properties

Lagrange interpolation is the *exact analog* of Chinese Remainder Theorem for polynomials.

Chinese Remainder Theorem: Suppose n_1, n_2, \dots, n_k are pairwise coprime. Then, for all integers a_1, a_2, \dots, a_k , there exists an integer x solving the below system of simultaneous congruences

$$x \equiv a_1 \pmod{n_1}$$

$$x \equiv a_2 \pmod{n_2}$$

$$\vdots$$

$$x \equiv a_k \pmod{n_k}$$

Further, all solutions x are congruent modulo $N = \prod_{i=1}^k n_i$.

Let $m_i = N/n_i$

i 'th "selector" number: $T_i = (m_i^{-1} \pmod{n_i}) m_i$

$$x = a_1 T_1 + a_2 T_2 + \dots + a_k T_k$$

Recall: Interpolation

Let pairs $(a_1, b_1), (a_2, b_2), \dots, (a_{d+1}, b_{d+1})$ from a field F be given (with all a_i 's distinct).

Theorem:

There is a unique degree d polynomial $P(x)$ satisfying $P(a_i) = b_i$ for all $i = 1 \dots d+1$.

A linear algebra view

Let $p(x) = p_0 + p_1 x + p_2 x^2 + \dots + p_d x^d$
Need to find the coefficient vector (p_0, p_1, \dots, p_d)

$$p(a) = p_0 + p_1 a + \dots + p_d a^d \\ = 1 \cdot p_0 + a \cdot p_1 + a^2 \cdot p_2 + \dots + a^d \cdot p_d$$

Thus we need to solve:

$$\begin{pmatrix} 1 & a_1 & a_1^2 & \dots & a_1^d \\ 1 & a_2 & a_2^2 & \dots & a_2^d \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & a_{d+1} & a_{d+1}^2 & \dots & a_{d+1}^d \end{pmatrix} \cdot \begin{pmatrix} p_0 \\ p_1 \\ \vdots \\ p_d \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_{d+1} \end{pmatrix}$$

Lagrange interpolation

The $(d+1) \times (d+1)$ Vandermonde matrix

$$M = \begin{pmatrix} 1 & a_1 & a_1^2 & \dots & a_1^d \\ 1 & a_2 & a_2^2 & \dots & a_2^d \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & a_{d+1} & a_{d+1}^2 & \dots & a_{d+1}^d \end{pmatrix}$$

is invertible.

• The determinant of M is nonzero when a_i 's are distinct

Thus can recover coefficient vector as $\vec{p} = M^{-1} \vec{b}$

The columns of M^{-1} are given by the coefficients of the various "selector" polynomials we constructed in Lagrange interpolation.

Representing Polynomials

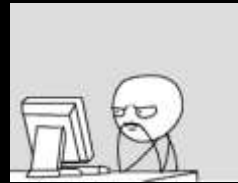
Let $P(x) \in F[x]$ be a degree- d polynomial.

Representing $P(x)$ using $d+1$ field elements:

1. List the $d+1$ coefficients.
2. Give P 's value at $d+1$ different elements.

Rep 1 to Rep 2: Evaluate at $d+1$ elements

Rep 2 to Rep 1: Lagrange Interpolation



Study Guide

Group Theory:

- Abelian
- Order theorem
- Isomorphism
- Subgroups

Number Theory:

- Euler's theorem
- Chinese Remainder theorem

Fields:

- Definitions
- Examples
- Finite fields of prime order

Polynomials:

- Degree- d polys have $\leq d$ roots.
- Polynomial division with remainder
- Lagrange Interpolation