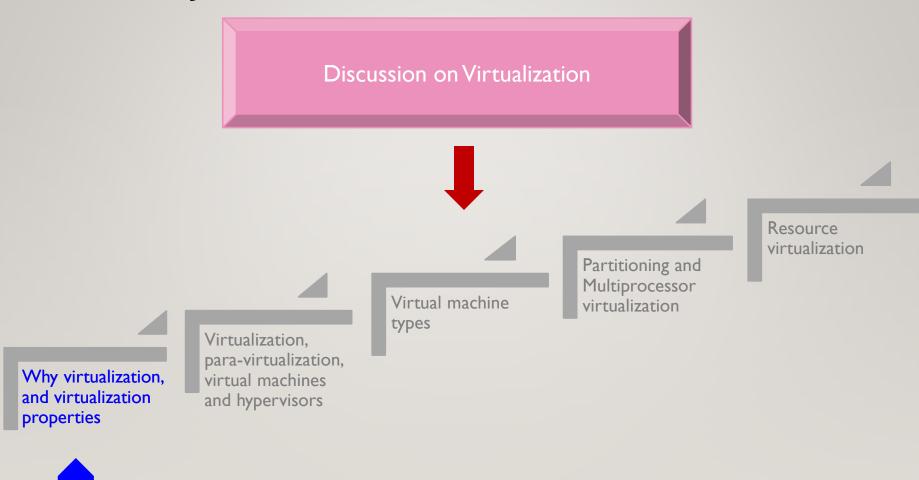
DISTRIBUTED SYSTEMS 14-736

GREGORY KESDEN

ADAPTED FROM 14-440 @ CMU-Q:

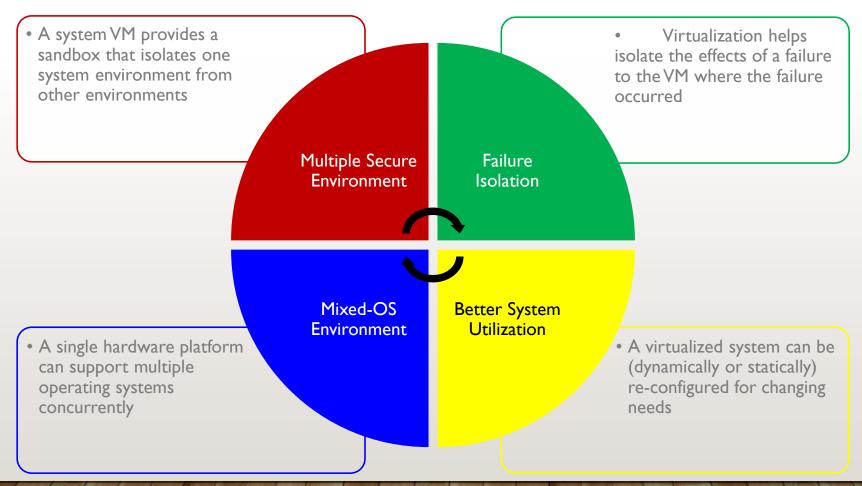
MAJD F. SAKR, MOHAMMAD HAMMOUD ANDVINAY KOLAR

OBJECTIVES



BENEFITS OF VIRTUALIZATION

 Here are <u>some</u> of the benefits that are typically provided by a virtualized system



OPERATING SYSTEMS LIMITATIONS

- OSs provide a way of virtualizing hardware resources among processes
- This may help isolate processes from one another
- However, this does not provide a <u>virtual machine</u> to a user who may wish to run a different OS
- Having hardware resources managed by a single OS limits the flexibility of the system in terms
 of available software, security, and
 failure isolation
- Virtualization typically provides a way of relaxing constraints and increasing flexibility

VIRTUALIZATION PROPERTIES

- Fault Isolation
- Software Isolation
- Performance Isolation (accomplished through scheduling and resource allocation)

Isolation

- All VM state can be captured into a file (i.e., you can operate on VM by operating on file—cp, rm)
- Complexity is proportional to virtual HW model and independent of guest software configuration

Encapsulation

2

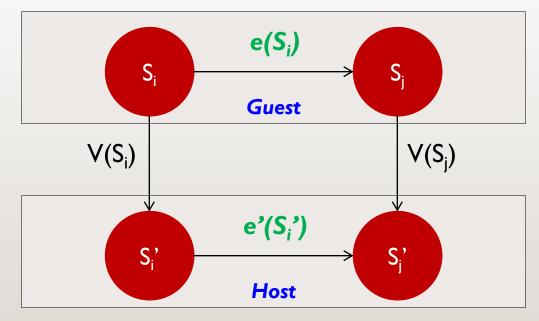
 All guest actions go through the virtualizing software which can inspect, modify, and deny operations

Interposition

5

WHAT IS VIRTUALIZATION?

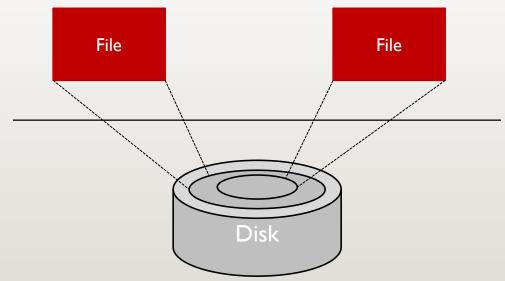
- Informally, a virtualized system (or subsystem) is a <u>mapping</u> of its interface, and all resources visible through that interface, to the interface and resources of a real system
- Formally, virtualization involves the construction of an isomorphism that <u>maps</u> a virtual guest system to a real <u>host</u> system (Popek and Goldberg 1974)
- ✓ Function V maps the guest state to the host state
- ✓ For a sequence of operations, e, that modifies a guest state, there is a corresponding e' in the host that performs an equivalent modification
- ✓ How can this be managed?



ABSTRACTION

 The key to managing complexity in computer systems is their division into levels of abstraction separated by well-defined interfaces

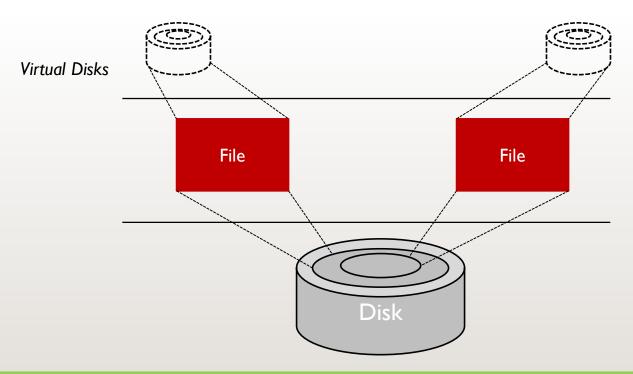
 Levels of abstraction allow implementation details at lower levels of a design to be ignored or simplified



- √ Files are an abstraction of a Disk
- ✓ A level of abstraction provides a simplified interface to underlying resources

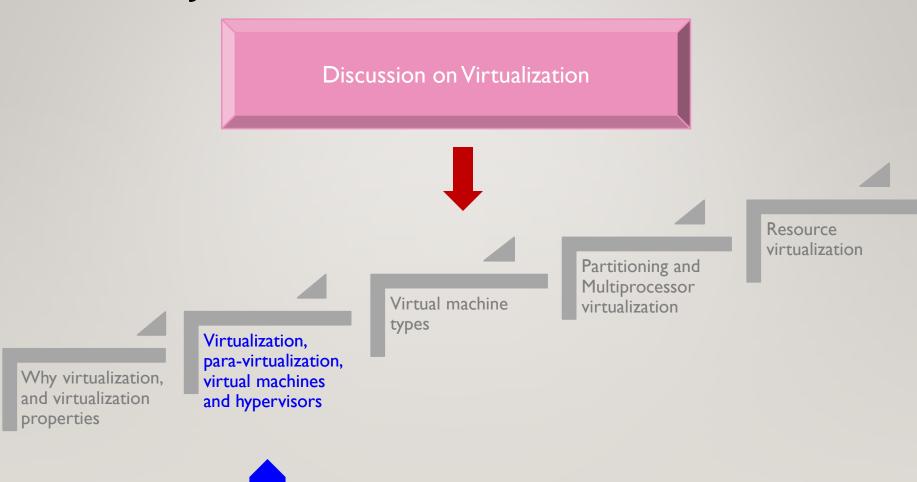
VIRTUALIZATION AND ABSTRACTION

• Virtualization uses abstraction but is different in that it doesn't necessarily hide details; the level of detail in a virtual system is often the same as that in the underlying real system



✓ Virtualization provides a different interface and/or resources at the same level of abstraction

OBJECTIVES

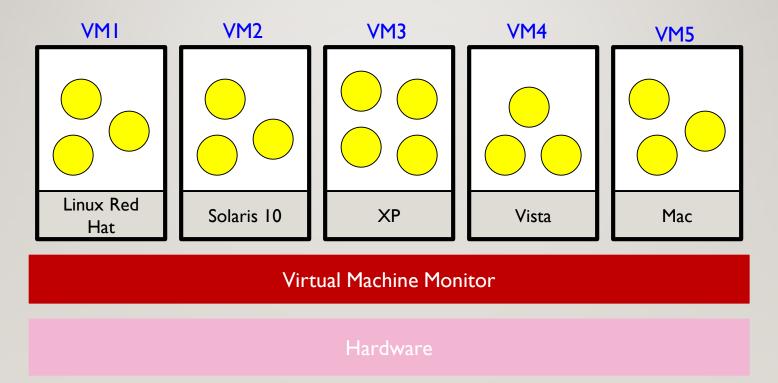


VIRTUAL MACHINES AND HYPERVISORS

- The concept of virtualization can be applied not only to subsystems such as disks, but to an entire machine denoted as a virtual machine (VM)
- A VM is implemented by adding a <u>layer of software</u> to a real machine so as to support the desired VM's architecture
- This layer of software is often referred to as virtual machine monitor (VMM)
- Early VMMs are implemented in firmware
- Today, VMMs are often implemented as a co-designed firmware-software layer, referred to as the hypervisor

A MIXED OS ENVIRONMENT

 Multiple VMs can be implemented on a single hardware platform to provide individuals or user groups with their own OS environments



FULL VIRTUALIZATION

- Traditional VMMs provide full-virtualization:
 - The functionally provided is identical to the underlying physical hardware
 - The functionality is exposed to the VMs
 - They allow unmodified guest OSs to execute on the VMs
 - This might result in some performance degradation
 - E.g., VMWare provides full virtualization

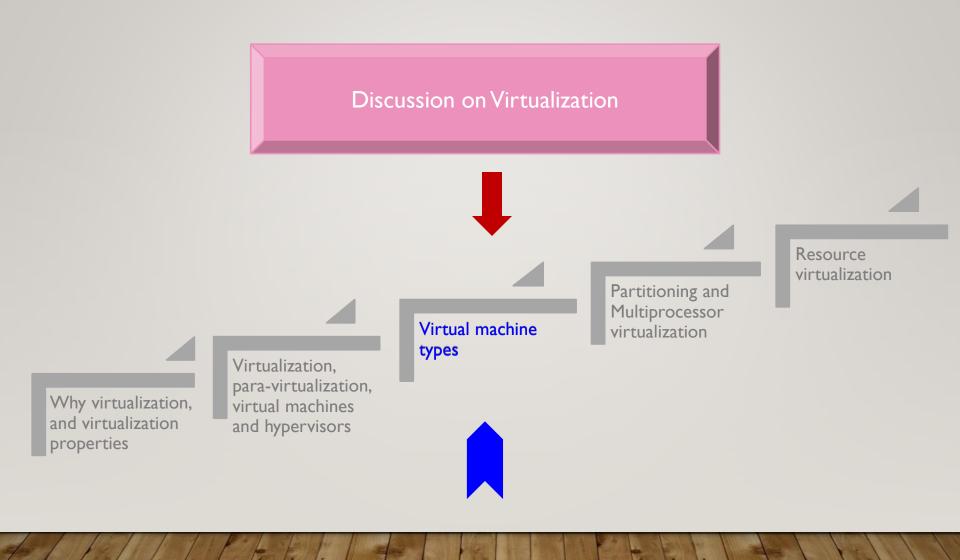
PARA-VIRTUALIZATION

- Other types of VMMs provide para-virtualization:
 - They provide a virtual hardware abstraction that is <u>similar, but not</u> <u>identical</u> to the real hardware
 - They modify the guest OS to cooperate with the VMM
 - They result in lower overhead leading to better performance
 - E.g., Xen provides both para-virtualization as well as full-virtualization

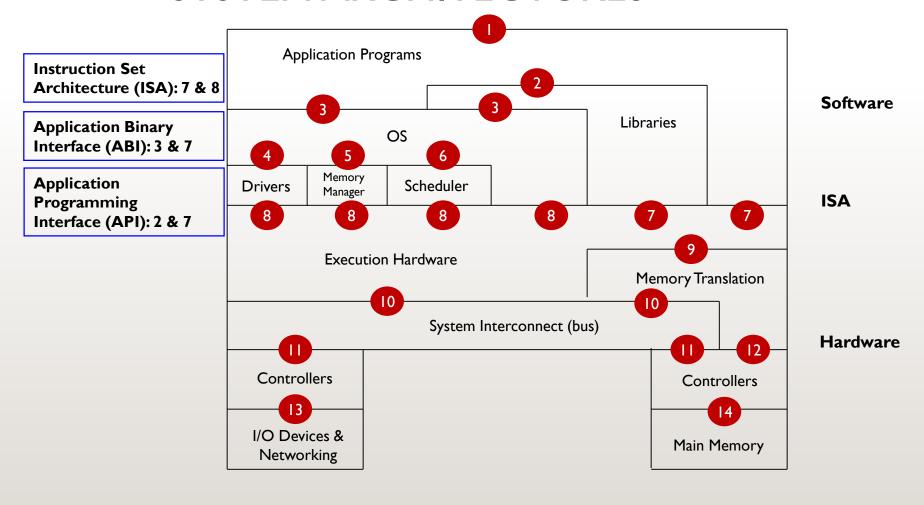
VIRTUALIZATION AND EMULATION

- VMs can employ emulation techniques to support cross-platform software compatibility
- Compatibility can be provided either at the system level (e.g., to run a Windows OS on Macintosh) or at the program or process level (e.g., to run Excel on a Sun Solaris/SPARC platform)
- Emulation is the process of implementing the interface and functionality of one system on a system having a different interface and functionality
- It can be argued that virtualization itself is simply a form of emulation

OBJECTIVES



BACKGROUND: COMPUTER SYSTEM ARCHITECTURES



TYPES OF VIRTUAL MACHINES

- As there is a process perspective and a system perspective of machines, there are also process-level and system-level VMs
- Virtual machines can be of two types:

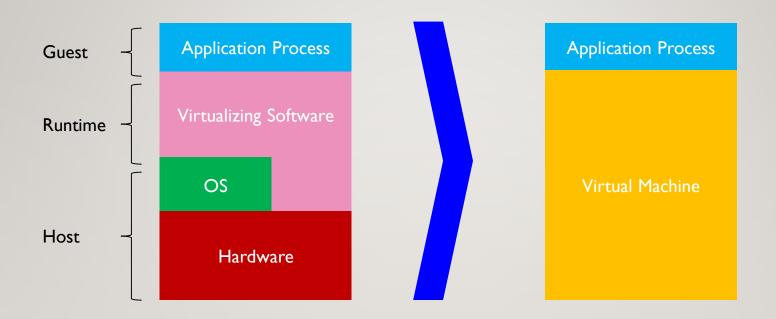
1. Process VM

Capable of supporting an individual process

2. System VM

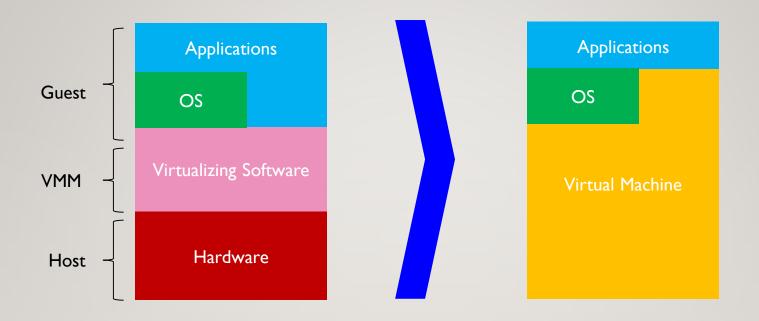
- Provides a complete system environment
- Supports an OS with potentially many types of processes

PROCESS VIRTUAL MACHINE



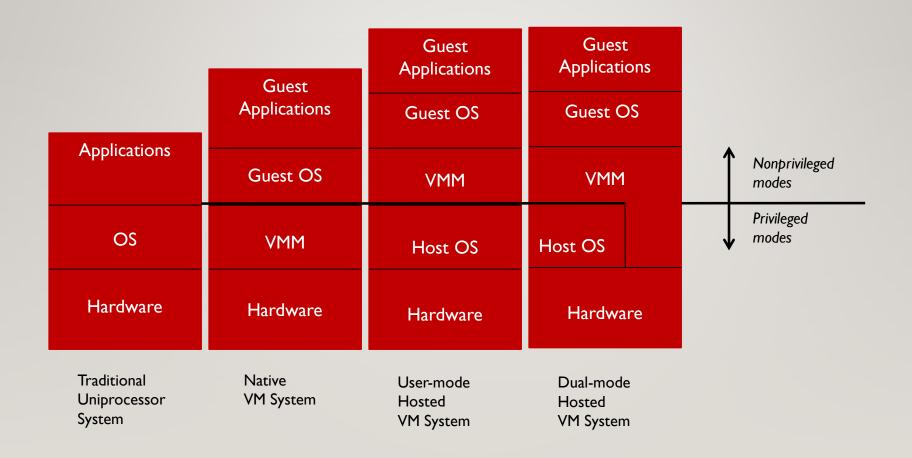
- ✓ Runtime is placed at the ABI interface
- ✓ Runtime emulates both user-level instructions and OS system calls

SYSTEM VIRTUAL MACHINE

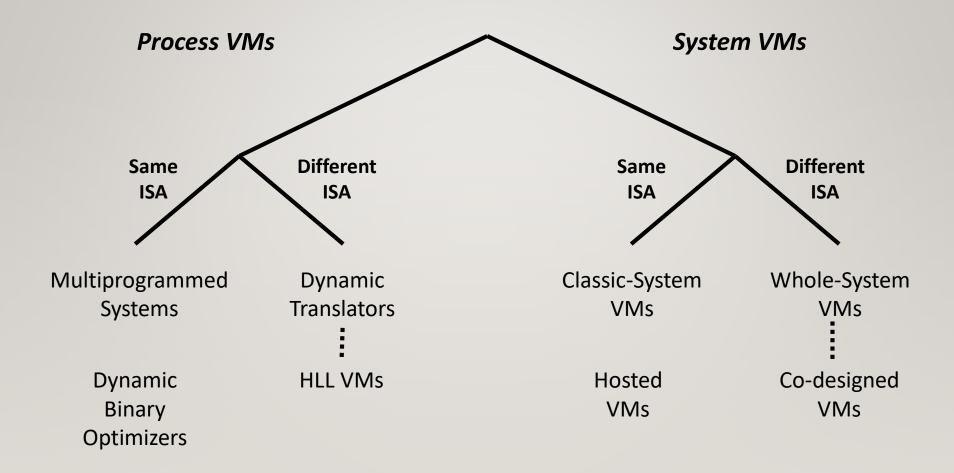


- ✓ VMM emulates the ISA used by one hardware platform to another, forming a system VM
- ✓ A system VM is capable of executing a system software environment developed for a different set of hardware

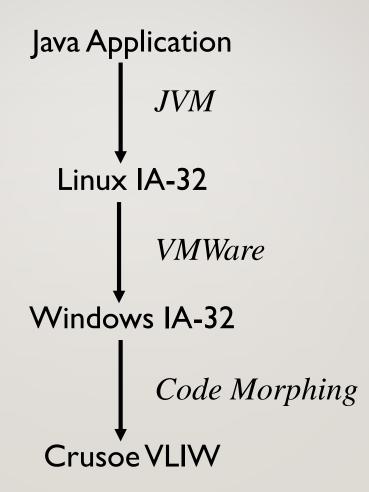
NATIVE AND HOSTED VM SYSTEMS



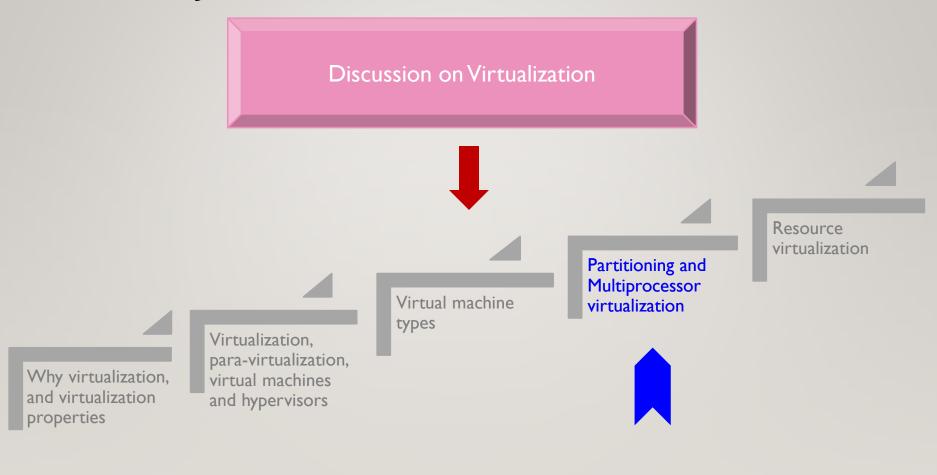
A TAXONOMY



THE VERSATILITY OF VMS



OBJECTIVES



MULTIPROCESSOR SYSTEMS

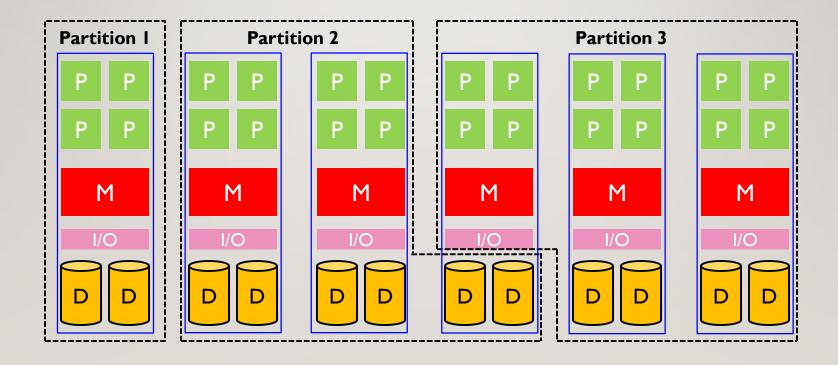
- Multiprocessor systems might have 1000s of processors connected to TBs of memory and PBs of disk capacity
- Often there is a mismatch between the ideal number of processors an application needs and the actual number of physical processors available
- It is more often the case that applications cannot exploit more than a fraction of the processors available. The is mainly because of:
 - Limitations in the parallelism available in the programs
 - Limitations in the scalability of applications due to the overhead of communication between processors

PARTITIONING

- The increasing availability of multiprocessor systems has led to the examination of techniques that can help utilize them more effectively
- Techniques have been developed in which the multiprocessor system can be partitioned into multiple partitions
 - A partition is given a subset of the resources available on the system
- Hence, using partitioning, multiple applications can simultaneously exploit the available resources of the system
- Partitioning can be achieved:
 - Either in-space (referred to as physical partitioning)
 - Or in-time (referred to as logical partitioning)

PHYSICAL PARTITIONING

 With physical partitioning, each partition is assigned resources that are physically distinct from the resources used by the other partitions



PHYSICAL PARTITIONING

- Physical partitioning allows a partition to own its resources physically
- It is not permissible for two partitions to share the resources of a single system board
- Partitions are configured by a central control unit that receives commands from the console of the system admin and provisions hardware resources accordingly
- The number of partitions that can be supported in physically partitioned systems is limited to the number of available physical processors

PHYSICAL PARTITIONING-ADVANTAGES

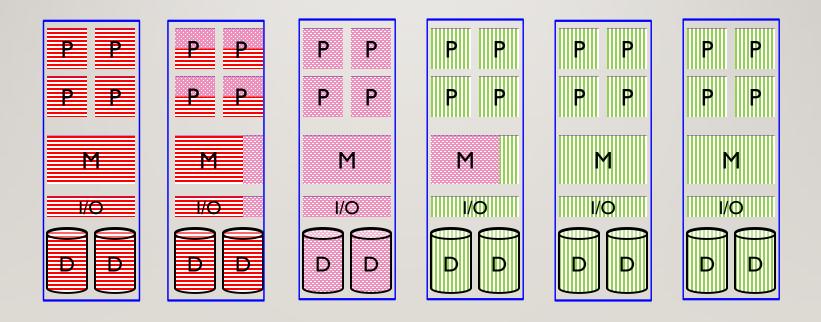
- Physical partitioning provides:
 - Failure Isolation: it ensures that in the event of a failure, only the part of the physical system that houses the failing partition will be affected
 - Better security isolation: Each partition is protected from the possibility of intentional or unintentional denial-of-service attacks by other partitions
 - Better ability to meet system-level objectives (these result from contracts between system owners and users of the system)
 - Easier management of resources: no need of sophisticated algorithms for scheduling and management of resources

PHYSICAL PARTITIONING-DISADVANTAGES

- While physical partitioning has a number of attractive features, it has some major disadvantages:
 - System utilization: Physical partitioning is probably not the ideal solution if system utilization is to be optimized
 - It is often the case that each of the physical partitions is underutilized
 - Load balancing: with physical partitioning, dynamic workload balancing becomes difficult to implement

LOGICAL PARTITIONING

With logical partitioning, partitions share some of the physical resources, usually in a time-multiplexed manner

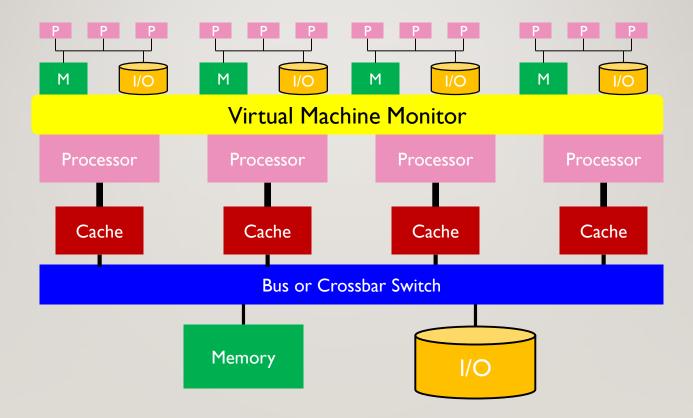


LOGICAL PARTITIONING

- With logical partitioning it is permissible for two partitions to share the resources of a single system board
- Logical partitioning makes it possible to partition an n-way system into a system with more than n partitions, if so desired
- Logical partitioning is more flexible than physical partitioning but needs additional mechanisms to provide safe and efficient way of sharing resources
- Logical partitioning is usually done through a VMM or a hypervisor and provides what is referred to as multiprocessor virtualization

MULTIPROCESSOR VIRTUALIZATION

 A virtualized multiprocessor gives the appearance of a system that may or may not reflect the exact configuration of the underlying physical system



IMPORTANCE IN DISTRIBUTED SYSTEMS

- Abstract environment from resources
 - Model need not reflect reality w.r.t. network, hardware, OS, software stack.
- Abstract application from environment
 - Each application can bring its own environment
- Enable homogeneity over heterogeneity
 - Different environments need not affect users
- Migration
 - Enable moving of applications
- Scaling
 - Add instances as needed
- Protection
 - Clear model
- Sharing
 - Partitioning easier and more fluid that physical resources.