## Computer Organization: Boolean Logic

## Representing and Manipulating Data

## Last Unit

- How to represent data as a sequence of bits
- How to interpret bit representations
- Use of levels of abstraction in representing more complex information (music, pictures) using simpler building blocks (numbers)


## This Unit

- How sequences of bits are implemented using electrical signals, and manipulated by circuits
- Use of levels of abstraction in designing more complex computer components from simpler components

Foundations
Boolean logic is the logic of digital circuits

## Implementing Bits

- Computers compute by manipulating electricity according to specific rules.
$\square$ We associate electrical signals inside the machine with bits. Any electrical device with two distinct states (e.g. on/off switch, two distinct voltage or current levels) could implement our bits.
- The rules are implemented by electrical circuits.


## Conceptualizing bits and circuits

$\square$ ON or 1: true

O OFF or 0: false

- circuit behavior: expressed in Boolean logic or Boolean algebra


## Boolean Logic (Algebra)

- Computer circuitry works based on Boolean Logic (Boolean Algebra) : operations on True (1) and False (0) values.

| $A$ | $B$ | $A \wedge B$ <br> (A AND B) <br> (conjunction) | A $\vee B$ <br> (A OR B) <br> (disjunction) |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 |



- $A$ and $B$ in the table are Boolean variables, AND and $O R$ are operations (also called functions).


## Foundations of Digital Computing

- Boolean Algebra was invented by George Boole in 1854 (before digital computers)
- Variables and functions take on only one of two possible values: True (1) or False (0).
- The correspondence between Boolean Logic and circuits was not discovered until 1930s
- Shannon's thesis: A Symbolic Analysis of Relay and Switching Circuits argued that electrical applications of Boolean Algebra could construct any logical, numerical relationship.
- We forget about the logical (truth and falsehood) aspect of Boolean logic and just manipulate symbols.


## Boolean Logic \& Truth Tables

E Example: You can think of $A \wedge B$ below as 15110 is fun and 15110 is useful where A stands for the statement 15110 is fun, $B$ stands for the statement 15110 is useful.

| $A$ | $B$ | $A \wedge B$ <br> (A AND B) <br> (conjunction) | $A \vee B$ <br> (A OR B) <br> (disjunction) |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 |



## Logic gates

the basic elements of digital circuits

## Logic Gates

$\square$ A gate is a physical device that implements a Boolean operator by performing basic operations on electrical signals.
$\square$ Nowadays, gates are built from transistors.

${ }_{B}^{A}-D-{ }_{\text {"OR" }}^{A \vee B}$


## A Mechanical Implementation

## Push-pull logic AND gate

$\square$ For an input pushed-in lever represents 1
$\square$ For an output pushed-in lever represents 0


Source:
randomwraith.com by Martin Howard

## Combinational circuits

combinations of logic gates

## Combinational Circuits

The logic states of inputs at any given time determine the state of the outputs.


What is $Q$ ? $\quad(A \wedge B) \vee((B \vee C) \wedge(C \wedge B))$

## Truth Table of a Circuit



How do I know that there should be 8 rows in the truth table?
http://www.allaboutcircuits.com/vol_4/chpt_7/6.html

| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 |  |
| 0 | 0 | 1 |  |
| 0 | 1 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |
|  |  |  |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 |  |
| 0 | 1 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | Q |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 |  |

Describes the relationship between inputs and outputs of a device

## Truth Table of a Circuit



| $A$ | $B$ | $C$ | Q |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

Describes the relationship between inputs and outputs of a device

## Describing Behavior of Circuits

- Boolean expressions
- Circuit diagrams
- Truth tables

Equivalent notations

# Manipulating circuits 

Boolean algebra and logical equivalence

## Why manipulate circuits?

$\square$ The design process
■simplify a complex design for easier manufacturing, faster or cooler operation, ...
$\square$ Boolean algebra helps us find another design guaranteed to have same behavior

## Logical Equivalence



| A | B | C | Q |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

Can we come up with a simpler circuit implementing the same truth table? Simpler circuits are typically cheaper to produce, consume less energy etc.

## Logical Equivalence



| $A$ | $B$ | $C$ | $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |


$Q=B \wedge(A \vee C)$
This smaller circuit is logically equivalent to the one above: they have the same truth table. By using laws of Boolean Algebra we convert a circuit to another equivalent circuit.

## Laws for the Logical Operators $\wedge$ and $\vee$

 (Similar to $\times$ and + )$\square$ Commutative: $\quad A \wedge B=B \wedge A \quad A \vee B=B \vee A$

- Associative:

$$
\begin{aligned}
& A \wedge B \wedge C=(A \wedge B) \wedge C=A \wedge(B \wedge C) \\
& A \vee B \vee C=(A \vee B) \vee C=A \vee(B \vee C)
\end{aligned}
$$

- Distributive:

$$
\begin{aligned}
& A \wedge(B \vee B)=(A \wedge B) \vee(A \wedge C) \\
& A \vee(B \wedge C)=(A \vee B) \wedge(A \vee C)
\end{aligned}
$$

- Identity:

$$
A \wedge 1=A
$$

$$
A \vee 0=A
$$

- Dominance:

$$
A \wedge 0=0
$$

$$
A \vee 1=1
$$

- Idempotence:

$$
A \wedge A=A
$$

$$
A \vee A=A
$$

- Complementation: $\mathrm{A} \wedge \neg \mathrm{A}=0$ $A \vee \neg A=1$
- Double Negation: $\neg \neg \mathrm{A}=\mathrm{A}$


# Laws for the Logical Operators $\wedge$ and $\vee$ (Similar to $\times$ and + ) 

- Commutative:

$$
A \wedge B=B \wedge A \quad A \vee B=B \vee A
$$

- Associative:

$$
\begin{aligned}
& A \wedge B \wedge C=(A \wedge B) \wedge C=A \wedge(B \wedge C) \\
& A \vee B \vee C=(A \vee B) \vee C=A \vee(B \vee C)
\end{aligned}
$$

- Distributive:

$$
\begin{aligned}
& A \wedge(B \vee C)=(A \wedge B) \vee(A \wedge C) \\
& A \vee(B \wedge C)=(A \vee B) \wedge(A \vee C)
\end{aligned}
$$

- Identity:
$A \vee 0=A$

The A's and B's here are schematic variables! You can instantiate them with any expression that has a Boolean value:

$$
(x \vee y) \wedge z=z \wedge(x \vee y) \text { (by commutativity) }
$$



## Applying Properties for $\wedge$ and $\vee$

## Showing $(x \wedge y) \vee((y \vee z) \wedge(z \wedge y))=y \wedge(x$ or $z)$

Commutativity $A \wedge B=B \wedge A$

$$
(x \wedge y) \vee((z \wedge y) \wedge(y \vee z))
$$

Distributivity $A \wedge(B \vee C)=(A \wedge B) \vee(A \wedge C)$

$$
(x \wedge y) \vee(z \wedge y \wedge y) \vee(z \wedge y \wedge z)
$$

Associativity, Commutativity, Idempotence

$$
(x \wedge y) \vee((z \wedge y) \vee(y \wedge z))
$$

Commutativity, idempotence $A \wedge A=A$

$$
((y) \wedge) \vee(\square \wedge z)
$$

Distributivity (backwards) (A) $\wedge B) \vee(A) \wedge C)=(A) \wedge(B \vee C)$

$$
\text { (v) } \wedge(x \vee z)
$$

Conclusion:

$$
(x \wedge y) \vee((y \vee z) \wedge(z \wedge y))=y \wedge(x \vee z)
$$

## Extending the system

more gates and DeMorgan's laws

## More gates (NAND, NOR, XOR)

| A | B | A nand <br> B | A nor B | A xor B |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 |

$\square$ nand ("not and"): $A$ nand $B=\operatorname{not}(A$ and $B)$


- nor ("not or"): A nor B = not (A or B)

- xor ("exclusive or"):
$A$ xor $B=(A$ and $\operatorname{not} B)$ or $(B$ and not $A)$



## A curious fact

$\square$ Functional Completeness of NAND and NOR
$\square$ Any logical circuit can be implemented using NAND gates only
$\square$ Same applies to NOR

## DeMorgan's Law

Nand: $\quad \neg(\mathrm{A} \wedge \mathrm{B})=\neg \mathrm{A} \vee \neg \mathrm{B}$

Nor: $\quad \neg(\mathrm{A} \vee \mathrm{B})=\neg \mathrm{A} \wedge \neg \mathrm{B}$

## DeMorgan's Law

Nand: $\neg(A \wedge B)=\neg A \vee \neg B$
if not ( $\mathrm{x}>15$ and $\mathrm{x}<110$ ): ...
is logically equivalent to
if (not $x>15$ ) or (not $x<110$ ): ...
Nor: $\quad \neg(A \vee B)=\neg A \wedge \neg B$
if not ( $\mathrm{x}<15$ or $\mathrm{x}>110$ ): ...
is logically equivalent to
if (not $x<15$ ) and (not $x>110$ ): ...

## A circuit for parity checking

Boolean expressions and circuits

## A Boolean expression that checks parity

$\square$ 3-bit odd parity checker F: an expression that should be true when the count of 1 bits is odd: when 1 or 3 of the bits are 1 s .

$$
P=(\neg A \wedge \neg B \wedge C) \vee(\neg A \wedge B \wedge \neg C) \vee(A \wedge \neg B \wedge \neg C) \vee(A \wedge B \wedge C)
$$

| $A$ | $B$ | $C$ | $P$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

There are specific methods for obtaining canonical Boolean expressions from a truth table, such as writing it as a disjunction of conjunctions or as a conjunction of disjunctions.

Note we have four subexpressions above each of them corresponding to exactly one row of the truth table where $P$ is 1 .

## The circuit

## 3-bit odd parity checker

$P=(\neg A \wedge \neg B \wedge C) \vee(\neg A \wedge B \wedge \neg C) \vee(A \wedge \neg B \wedge \neg C) \vee(A \wedge B \wedge C)$

| A | B | C | P |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |



## Summary

You should be able to:

- Identify basic gates
- Describe the behavior of a gate or circuit using Boolean expressions, truth tables, and logic diagrams
- Transform one Boolean expression into another given the laws of Boolean algebra


## Next Time

- How circuits are combined to form a computer
-Von Neumann architecture revisited
- Fetch - Decode - Execute Cycle


