

## Smoothing for non-smooth optimization, lecture 1

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### Goals:

- First-order (gradient-based) algorithms for huge convex optimization problems
- Two applications:
  - Game theory:  $\max_{y \in Q_2} \min_{x \in Q_1} \langle Ax, y \rangle$
  - Sparse principal components analysis:

$$\max_{X \in Q_1} \min_{Y \in Q_2} \{ \langle C, X \rangle - \langle X, Y \rangle \}$$

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### Course outline:

- Introduction & motivation
- Basics of convex analysis
- Gradient and subgradient schemes
- Smoothing techniques
- Applications

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### About algorithms for convex optimization

- Interior-point methods: based on Newton's method.  
 $O(\log(1/\epsilon))$  expensive iterations to find  $\epsilon$ -approximate solution
- Subgradient methods (non-smooth case)  
 $O(1/\epsilon^2)$  cheap iterations to find  $\epsilon$ -approximate solution
- Efficient gradient-based methods (smooth case)  
 $O(1/\sqrt{\epsilon})$  cheap iterations to find  $\epsilon$ -approximate solution
- Smoothing techniques (non-smooth case)  
 $O(1/\epsilon)$  cheap iterations to find  $\epsilon$ -approximate solution

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### Brief, incomplete, and biased history of convex optimization

- 19th century: optimization in physics, Gauss's least-squares method
- 1900-1970: math developments, convex optimization
- 1940s: simplex method for linear programming
- 1970s: ellipsoid method
- 1980s: interior-point methods
- 1990s: algorithms for semidefinite programming, and symmetric cone programming, new applications in stats, machine learning, combinatorial optimization, control, circuit design, etc.
- 2000s: methods with low computational cost for huge problems

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## Convex optimization

$$\begin{array}{ll}\min & f(x) \\ \text{s.t.} & x \in D\end{array}$$

where

- $f$ : convex function
- $D$ : convex set

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## Examples of convex optimization

- Least squares

$$\min_x \|Ax - b\|^2$$

- Linear programming

$$\begin{array}{ll}\min & c^\top x \\ \text{s.t.} & Ax = b \\ & x \geq 0\end{array}$$

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- Semidefinite programming

$$\begin{array}{ll}\min & C \bullet X \\ \text{s.t.} & AX = b \\ & X \succeq 0\end{array}$$

Here the variable  $X$  is an  $n \times n$  symmetric matrix and

$$X \succeq 0 \Leftrightarrow u^\top Xu \geq 0 \quad \forall u \in \mathbb{R}^n \Leftrightarrow \lambda(X) \geq 0.$$

Convention:  $\lambda(X) = (\lambda_1(X), \dots, \lambda_n(X))$  eigenvalues of  $X$  with

$$\lambda_1(X) \geq \dots \geq \lambda_n(X).$$

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## Today:

- Convex sets
- Convex functions
- Smooth convex functions
- Non-smooth convex functions
- Fenchel duality

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## Convex sets

Throughout this course

$E$ : finite-dimensional Euclidean space with inner product  $\langle \cdot, \cdot \rangle$ .

$\| \cdot \|$ : norm induced by  $\langle \cdot, \cdot \rangle$ .

Two special cases:

- $\mathbb{R}^n$ , with inner product  $\langle x, s \rangle = x^T s$
- $S^n$ : space of  $n \times n$  real symmetric matrices with inner product  $\langle X, S \rangle = \text{trace}(XS) =: X \bullet S$ .

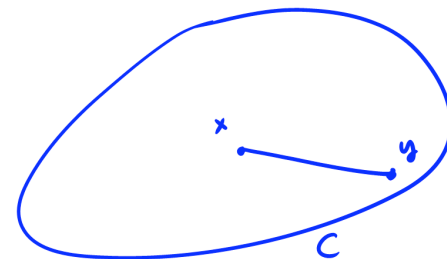
Recall:  $\text{trace}(M) = \sum_i \lambda_i(M) = \sum_i M_{ii}$ .

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## Example 2

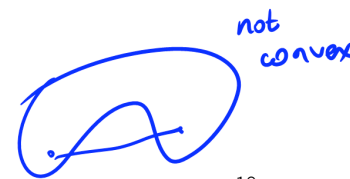
- Polyhedron:  $\{x \in \mathbb{R}^n : Ax \leq b\}$  for some given  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ .
- Non-negative orthant:  $\mathbb{R}_+^n := \{x \in \mathbb{R}^n : x \geq 0\}$
- Polytope:  $\{My : y \geq 0, \sum_{j=1}^n y_j = 1\}$  for some given  $M \in \mathbb{R}^{n \times m}$ .
- Simplex:  $\Delta_n := \{x \in \mathbb{R}^n : x \geq 0, \sum_{i=1}^n x_i = 1\}$
- Ball:  $\{x \in E : \langle x, x \rangle \leq r^2\}$  for some given  $r > 0$ .
- Positive semidefinite cone:  $S_+^n := \{X \in S^n : \lambda(X) \geq 0\}$
- Spectraplex:  $\{X \in S_+^n : \sum_{i=1}^n \lambda_i(X) = 1\}$ .

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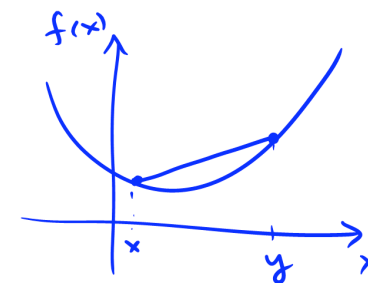


## Definition 1

- $C \subseteq E$  is *convex* if  $\alpha x + (1 - \alpha)y \in C$  whenever  $x, y \in C$  and  $\alpha \in [0, 1]$ .
- $K \subseteq E$  is a *cone* if  $\lambda x \in K$  whenever  $x \in K$  and  $\lambda \geq 0$ .



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## Convex functions

**Definition 3** Let  $C \subseteq E$  be a convex set and  $f : C \rightarrow \mathbb{R}$ . The function  $f$  is *convex* on  $C$  if for all  $x, y \in C$  and  $\alpha \in [0, 1]$

$$f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y).$$

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### Alternative definition of convex functions

- Consider extended-valued functions  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ .
- Can always do this by declaring  $f(x) = +\infty$  for  $x \notin \text{dom}(f)$
- $f$  is convex if and only its *epigraph*

$$\text{epi}(f) := \{(x, r) \in E \times \mathbb{R} : r \geq f(x)\},$$

is a convex set.

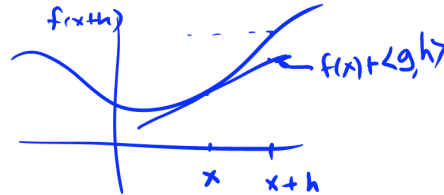
- Define  $\text{dom}(f) := \{x \in E : f(x) < \infty\}$  (projection of  $\text{epi}(f)$ )

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### Example 4 Some convex functions

- $f : E \rightarrow \mathbb{R}$  defined by  $x \mapsto \langle x, x \rangle = \|x\|^2$
- $f : E \rightarrow \mathbb{R}$  defined by  $x \mapsto \langle a, x \rangle$  for some  $a \in E$
- $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  defined by  $x \mapsto \sup_{i \in I} g_i(x)$  where each  $g_i : E \rightarrow \mathbb{R} \cup \{+\infty\}$ ,  $i \in I$  is convex
- $f : \mathbb{R} \rightarrow \mathbb{R}$  defined by  $x \mapsto e^x$
- $f : \mathbb{R}_+ \rightarrow \mathbb{R}$  defined by  $x \mapsto x \log x$
- $f : \mathbb{S}^n \rightarrow \mathbb{R}$  defined by  $X \mapsto \lambda_1(X) + \dots + \lambda_k(X)$ , where  $k \leq n$

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### Smooth convex functions

**Definition 5**  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  is *differentiable* at  $x \in \text{dom}(f)$  if there exists  $g \in E$  such that

$$\lim_{h \rightarrow 0} \frac{f(x+h) - f(x) - \langle g, h \rangle}{\|h\|} = 0,$$

In other words, if

$$f(x+h) = f(x) + \langle g, h \rangle + o(\|h\|).$$

When such  $g$  exists, it is unique. In this case,  $g$  is the *gradient* of  $f$  at  $x$ , and is denoted  $\nabla f(x)$ .

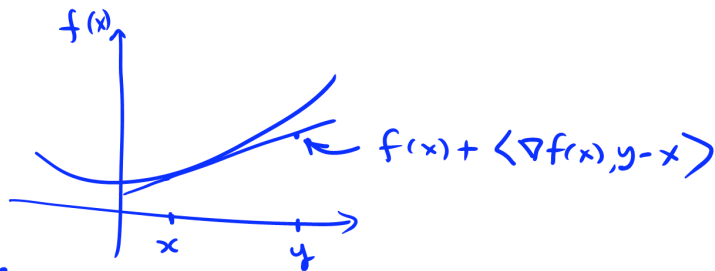
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**Proposition 6** Assume  $f$  is differentiable on an open set  $\Omega \subseteq E$  and  $C \subseteq \Omega$  is convex. Then the following are equivalent

- $f$  is convex on  $C$
- $f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$  for all  $x, y \in C$
- $\langle \nabla f(y) - \nabla f(x), y - x \rangle \geq 0$  for all  $x, y \in C$

**Corollary 7** Assume  $f : E \rightarrow \mathbb{R}$  is differentiable everywhere and convex. Then  $\bar{x} = \text{argmin}_x f(x)$  if and only if  $\nabla f(\bar{x}) = 0$ .

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To prove:

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$$

Pf. Consider  $\varphi(t) := f(x + t(y - x))$   
 $\varphi'(0) = \langle \nabla f(x), y - x \rangle$

$$\begin{aligned} \varphi(t) &= f(x + t(y - x)) = f(ty + (1-t)x) \leq tf(y) + (1-t)f(x) \\ &\leq f(x) + t(f(y) - f(x)) = \varphi(0) + t(f(y) - f(x)) \\ [\varphi(t) - \varphi(0)]/t &\leq f(y) - f(x) \end{aligned}$$

## Strong convexity

**Definition 9** Let  $C \subseteq E$  be a convex set and  $f : C \rightarrow \mathbb{R}$ . The function  $f$  is strongly convex on  $C$  if there exists  $\rho > 0$  such that for all  $x, y \in C$  and  $\alpha \in [0, 1]$

$$f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y) - \frac{1}{2}\rho\alpha(1 - \alpha)\|x - y\|^2.$$

In this case say that  $f$  is strongly convex with modulus  $\rho$ .

$$\nabla f : \Omega \subseteq E \rightarrow E$$

$$\nabla f \text{ Lipschitz if } \|\nabla f(x) - \nabla f(y)\| \leq L\|x - y\|$$

## Lipschitz continuity of the gradient

**Theorem 8** Assume  $f$  is differentiable on an open set  $\Omega \subseteq E$  and  $C \subseteq \Omega$  is convex. Then the following are equivalent:

- $f$  is convex and  $\nabla f$  is Lipschitz with constant  $L$  on  $C$
- $0 \leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle \leq \frac{L}{2}\|y - x\|^2$  for all  $x, y \in C$
- $\frac{1}{2L}\|\nabla f(y) - \nabla f(x)\|^2 \leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle$  for all  $x, y \in C$
- $\frac{1}{L}\|\nabla f(y) - \nabla f(x)\|^2 \leq \langle \nabla f(y) - \nabla f(x), y - x \rangle$  for all  $x, y \in C$

$$\rightarrow \frac{1}{2L} \|\nabla f(x) - \nabla f(y)\|^2 \leq f(x) - f(y) - \langle \nabla f(y), x - y \rangle$$

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**Proposition 10** Let  $C \subseteq E$  be a convex set and  $f : C \rightarrow \mathbb{R}$ . Then  $f$  is strongly convex with modulus  $\rho > 0$  if and only if  $f - \frac{1}{2}\rho\|\cdot\|^2$  is convex on  $C$ .

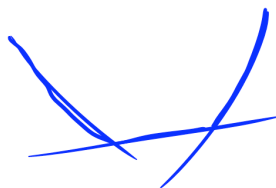
**Theorem 11** Assume  $f$  is differentiable on an open set  $\Omega \subseteq E$  and  $C \subseteq \Omega$  is convex. Then the following are equivalent

- $f$  is strongly convex on  $C$  with modulus  $\rho$
- $f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{1}{2}\rho\|x - y\|^2$  for all  $x, y \in C$
- $\langle \nabla f(y) - \nabla f(x), y - x \rangle \geq \rho\|x - y\|^2$  for all  $x, y \in C$

$f$  strongly convex w/ modulus  $\rho \Leftrightarrow$   
 $\rightarrow f - \frac{1}{2}\rho \|\cdot\|^2$  convex

$$f(\alpha x + (1-\alpha)y) - \frac{1}{2}\rho \|\alpha x + (1-\alpha)y\|^2 \leq \alpha \left[ f(x) - \frac{1}{2}\rho \|x\|^2 \right] + (1-\alpha) \left[ f(y) - \frac{1}{2}\rho \|y\|^2 \right]$$

$$f(\alpha x + (1-\alpha)y) \leq \alpha f(x) + (1-\alpha)f(y) - \frac{1}{2}\rho \left[ \underbrace{\alpha \|x\|^2 + (1-\alpha)\|y\|^2 - \|\alpha x + (1-\alpha)y\|^2}_{\alpha(1-\alpha)\|y-x\|^2} \right]$$



### Non-smooth convex functions

Non-smoothness is pervasive in optimization.

### Example 13

- $f(x) = \sup \{g_i(x) : i \in I\}$  where  $g_i, i \in I$  are convex.
- $f(x) = \inf \{\phi(u) : c_i(u) \leq x_i, i = 1, \dots, n\}$  where  $\phi, c_i, \dots, c_n$  are convex.

### Example 12 Some strongly convex functions

- $f : E \rightarrow \mathbb{R}$  defined by  $x \mapsto \langle x, x \rangle$
- $f : \Delta_n \rightarrow \mathbb{R}$  defined by  $x \mapsto \sum_{i=1}^n x_i \log x_i$
- $f : \{X \in \mathbb{S}^n : X \succeq 0, I \bullet X = 1\} \rightarrow \mathbb{R}$  defined by  $X \mapsto \sum_{i=1}^n \lambda_i(X) \log \lambda_i(X)$



**Proposition 14** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  convex and  $x \in \text{int}(\text{dom}(D))$ . Then there exists  $g \in E$  such that

$$f(y) - f(x) \geq \langle g, y - x \rangle \text{ for all } y \in E. \quad (1)$$

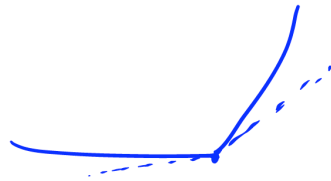
Generalization of the gradient:

$$\dots f(y) - f(x) \geq \langle \nabla f(x), y - x \rangle$$

**Definition 15** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  convex and  $x \in \text{dom}(D)$ . The *subdifferential* of  $f$  at  $x$  is

$$\partial f(x) := \{g \in E : (1) \text{ holds}\}$$

**Corollary 16** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  is convex. Then  $\bar{x} \in \text{argmin}_x f(x)$  if and only if  $0 \in \partial f(\bar{x})$ .



**Notice:**

If  $f$  is convex and differentiable at  $x$  then  $\partial f(x) = \{\nabla f(x)\}$ .

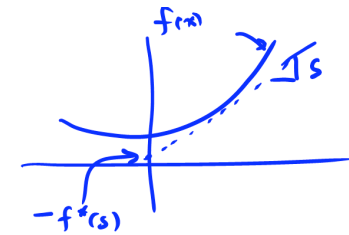
**Proposition 17** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  convex. Then  $f$  is differentiable at  $x$  if and only if  $\partial f(x)$  is a singleton.

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**Example 19**

$f(x)$	$f^*(s)$
$e^x$	$s \log s - s$
$\frac{ x ^p}{p}, p > 1$	$\frac{ s ^q}{q}, \frac{1}{q} = 1 - \frac{1}{p}$
$\sqrt{1+x^2}$	$-\sqrt{1-s^2}$
$-\log x$	$-1 - \log(-s)$

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**Fenchel conjugate**

**Definition 18** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ , with  $\text{dom}(f) \neq \emptyset$ . The conjugate of  $f$  is  $f^* : E \rightarrow \mathbb{R} \cup \{+\infty\}$  defined by

$$s \mapsto f^*(s) := \sup \{ \langle s, x \rangle - f(x) : x \in \text{dom}(f) \}.$$

**Notice:**

- $f^*(s) + f(x) \geq \langle s, x \rangle$  for all  $x, s \in E$
- $f^*(s) + f(x) = \langle s, x \rangle$  if and only if  $s \in \partial f(x)$

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**Definition 20**  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  is closed if  $\text{epi}(f)$  is closed.

**Proposition 21** Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ .

- $f^*$  is closed and convex
- $f^{**} \leq f$
- $f^{**} = f$  if and only if  $f$  is closed and convex.

$$f^*(s) = \sup_x \{ \langle s, x \rangle - f(x) \}$$

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## Fenchel duality

Assume  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ ,  $g : Y \rightarrow \mathbb{R} \cup \{+\infty\}$  and  $A \in L(E, Y)$ .


Consider the pair

$$p := \inf_{x \in E} \{f(x) + g(Ax)\} \quad (P)$$

and

$$d := \sup_{y \in Y} \{-f^*(A^*y) - g^*(-y)\} \quad (D)$$

### Theorem 22

- (Weak duality):  $p \geq d$  
- (Strong duality): If  $0 \in \text{relint}(\text{dom}(g) - A\text{dom}(f))$  then  $p = d$  and  $d$  is attained.

$$f(x) + g(Ax) + f^*(A^*y) + g^*(-y) \geq \langle x, A^*y \rangle + \langle Ax, -y \rangle = 0$$

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## Special case

Consider

$$\min_{x \in Q_1} \max_{y \in Q_2} \{\hat{f}(x) - \hat{\phi}(y) + \langle Ax, y \rangle\}. \quad (2)$$

where

- $E_1, E_2$  are finite dimensional Euclidean spaces,  $A \in L(E_1, E_2)$
- $Q_i \subseteq E_i$  are simple compact convex sets
- $\hat{f}$  and  $\hat{\phi}$  are convex and differentiable everywhere

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Can write (2) as

$$\min_{x \in E} \{f(x) + \phi^*(Ax)\}$$

for

- $f(x) = \hat{f}(x)$ ,  $\text{dom}(f) := Q_1$
- $\phi(y) = \hat{\phi}(y)$ ,  $\text{dom}(\phi) := Q_2$

Fenchel dual:

$$\max_{y \in E_2} \{-f^*(A^*y) - \phi^{**}(-y)\}$$

which can be written as

$$\max_{y \in Q_2} \min_{x \in Q_1} \{\hat{f}(x) - \hat{\phi}(y) + \langle Ax, y \rangle\}.$$

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Fenchel duality yields

$$\min_{x \in Q_1} \max_{y \in Q_2} \{\hat{f}(x) - \hat{\phi}(y) + \langle Ax, y \rangle\} = \max_{y \in Q_2} \min_{x \in Q_1} \{\hat{f}(x) - \hat{\phi}(y) + \langle Ax, y \rangle\}.$$

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## References for today's material

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