

# Compressive (or Compressed) Sensing "Detección Comprimida"

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Materials available at  
<http://andrew.cmu.edu/user/jfp/UNencuentro>

# Compressive Sensing, Lecture 1

Consider the following two pictures

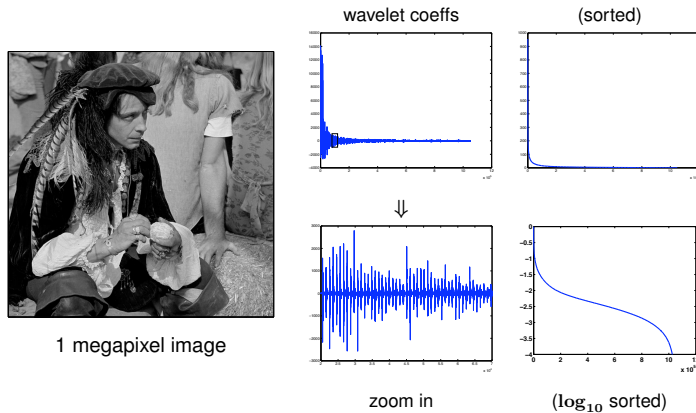


One of these is a raw jpg 2.7MB file. The other one is a compressed jpg 300KB version.

Can you tell them apart?

## Wavelets and images

Compressibility corresponds to sparsity in a suitable basis.



1 megapixel image

zoom in

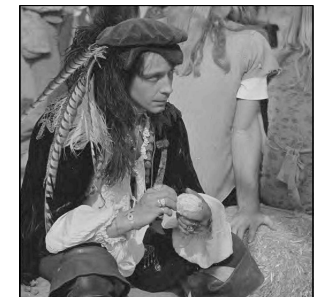
(log<sub>10</sub> sorted)

## Compression

- Take original 1 megapixel image
- Compute all 1 million wavelet coefficients
- Keep only the 25K largest and set the others to zero
- Invert the wavelet transform

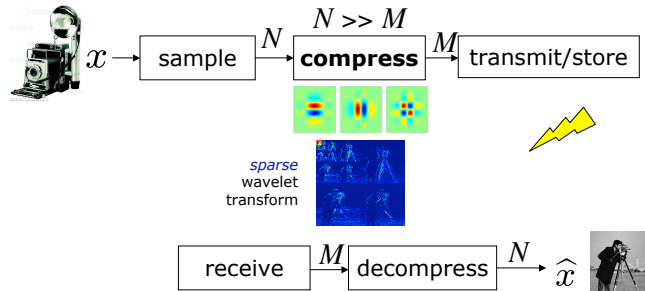


original image



25K approximation

## Conventional signal acquisition paradigm



### Question

If the signal is compressible, can it be acquired efficiently?

### Plan

- Introduction to compressive sensing, undetermined systems of equations,  $\ell_1$  minimization
- Main theoretical and computational techniques
- Matrix completion, undetermined linear matrix equations, nuclear norm minimization.

## Compressive sensing

- Term coined by Donoho
- Body of theory and algorithms for sparse signal acquisition and recovery
- Seminal papers:
  - Candès, Romberg and Tao (2006)
  - Candès and Tao (2006)
  - Donoho (2006)
- Hot area of research spanning information theory, signal processing, statistics, mathematics, etc.
- Applications where measurements are
  - slow or costly (MRI)
  - missing or wasteful
  - beyond other capabilities such as memory

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## Undetermined systems of equations

### Problem

Recover a signal  $\bar{x} \in \mathbb{R}^n$  from  $m \ll n$  linear measurements

$$b_k = \langle a_k, \bar{x} \rangle, \quad k = 1, \dots, m \rightsquigarrow b = A\bar{x}$$

- In general this is impossible.
- Suppose we know that  $\bar{x}$  is sparse. Does that help?

### Example

Suppose only one component of  $\bar{x}$  is different from zero. Can we get by with fewer than  $n$  measurements?

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## $\ell_0$ versus $\ell_1$ -optimization

Possible approach to recover sparse  $\bar{x}$

Take  $m \ll n$  measurements  $b = A\bar{x}$  and then solve

$$\min_{Ax = b} \|x\|_0$$

Here  $\|\cdot\|_0$  stands for the  $\ell_0$  quasi-norm:  $\|x\|_0 = |\{i : x_i \neq 0\}|$ .

### Relevant questions

- Does this work (provided we take enough measurements)?
- Suppose  $\bar{x}$  is  $k$ -sparse, i.e.,  $\|\bar{x}\|_0 = k$ . How many measurements suffice?
- How hard is it to solve the above  $\ell_0$ -minimization problem?

$$\min_{Ax = b} \|x\|_0$$

computationally hard

$$\min_{Ax = b} \|x\|_1$$

computationally tractable  
(linear program)

- $\ell_1$  norm is the convex envelope of the  $\ell_0$  quasi-norm.
- The  $\ell_1$  minimization problem is a *convex relaxation* of the  $\ell_0$  minimization problem.
- In many cases the above  $\ell_1$  minimization problem yields the same solution as the  $\ell_0$  problem.

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## A bit of history of $\ell_1$ optimization

$\ell_1$  minimization often finds the “right” answer

- Seismology
- Lasso regression
- Bandlimited deconvolution
- Total variation (TV) denoising
- Basis pursuit

## Acquiring a sparse signal

Suppose  $\bar{x} \in \mathbb{R}^n$  is  $s$ -sparse.

- Take  $m$  random and nonadaptive measurements

$$b_k = \langle a_k, \bar{x} \rangle, \quad k = 1, \dots, m \rightsquigarrow b = A\bar{x}$$

- Try to reconstruct  $\bar{x}$  via  $\ell_1$  minimization.

### First fundamental result

If  $m \gtrsim s \cdot \log n$  and the  $a_k$  are suitably chosen, then the recovery is exact.

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# Nonadaptive sensing of compressible signals

## Classical approach

- Measure the full signal  $\bar{x}$  (all coefficients)
- Store  $s$  largest coefficients
- Distorsion  $\|\bar{x} - \bar{x}_s\|_2$

## Second fundamental result

If  $m \gtrsim s \cdot \log n$  then

$$\|\hat{x} - \bar{x}\|_2 \lesssim \|\bar{x} - \bar{x}_s\|_2$$

## Compressive sensing

- Take  $m$  random measurements
- Reconstruct  $\hat{x}$  via  $\ell_1$  minimization

## Optimality of compressive sensing

It is not possible to do better

- with fewer measurements
- with other reconstruction algorithms

## Key features of compressive sensing

- Obtain compressible signals from few sensors
- Sensing is **nonadaptive**: no knowledge about the signal
- Simple acquisition followed by  $\ell_1$  decoder

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## Next: formal statements

- Probabilistic approach
  - Isotropy and incoherence
  - Incoherent sampling theorems
- Deterministic approach
  - Restricted isometry property
  - Signal recovery theorem
- Robustness to noise
- Optimality

## Probabilistic approach: random sensing

Acquire  $\bar{x} \in \mathbb{R}^n$  by measuring

$$b_k = \langle a_k, \bar{x} \rangle, \quad k = 1, \dots, m \rightsquigarrow b = A\bar{x}$$

where  $a_k$  are iid  $F$  for some distribution  $F$  in  $\mathbb{C}^n$ .

## Two key properties

- Isotropy:  $\mathbb{E}(aa^*) = I$
- Coherence measure  $\mu(F)$ : smallest number such that

$$\max_{i=1, \dots, n} |\langle a, e_i \rangle|^2 \leq \mu(F) \quad \text{with high probability}$$

We want low coherence  $\mu(F)$ .

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## Coherence

### Observe

- $\mathbb{E}(aa^*) = I$  implies  $\mu(F) \geq 1$ .
- We would like  $\mu(F)$  to be as close as possible to 1 (incoherent sensing).

### Examples of isotropic incoherent sensing

- Gaussian sensing
- Binary sensing
- Partial Fourier transform

Notation: for  $a \in \mathbb{C}^n$  and  $i = 1, \dots, n$

$$a[i] := \langle a, e_i \rangle.$$

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## An example of coherent sensing

### Sample random components of $x$

- Select  $j$  uniformly at random in  $\{1, \dots, n\}$
- Set  $a = \sqrt{n}e_j$ .

In this case

$$\mathbb{E}(aa^*) = I$$

and

$$\max_{i=1, \dots, n} |a[i]|^2 = n.$$

For this type of sensing, how many samples do we need to recover a 1-sparse vector with high probability?

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## Isotropic incoherent sensing

### Gaussian sensing

$a \sim \mathcal{N}(0, I)$ , that is,  $a[1], \dots, a[n]$  are iid  $\mathcal{N}(0, 1)$ .

In this case  $\mu(F) = \log n$ .

### Binary sensing

$a[1], \dots, a[n]$  are iid with distribution  $\mathbb{P}(a[i] = \pm 1) = 1/2$ .

In this case  $\mu(F) = 1$ .

### Partial discrete Fourier transform

- Select  $k$  uniformly at random in  $\{0, 1, \dots, n-1\}$
- Set  $a[t] := e^{i2\pi kt/n}$ ,  $t = 0, 1, \dots, n-1$ .

In this case  $\mu(F) = 1$ .

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## Incoherent sampling theorems

### Compressive sampling approach

measure

$$b_k = \langle a_k, \bar{x} \rangle, \quad k = 1, \dots, m \rightsquigarrow b = A\bar{x}$$

perform  $\ell_1$  recovery

$$\hat{x} := \underset{x}{\operatorname{argmin}} \{ \|x\|_1 : Ax = b \}$$

### Theorem (Candès & Plan)

Assume  $\bar{x}$  is  $s$ -sparse. Then recovery is exact with probability at least  $1 - 5/n - e^{-\beta}$  provided that

$$m \geq C_0 \cdot (1 + \beta) \cdot \mu(F) \cdot s \cdot \log n$$

for some constant  $C_0$ .

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Related predecessors:

### Theorem (Candès & Tao)

Assume  $\bar{x}$  is  $s$ -sparse and sample  $m$  Fourier coefficients selected at random. Then recovery is exact with probability at least  $1 - \mathcal{O}(n^{-\beta})$  provided that

$$m \geq C_\beta \cdot s \cdot \log n$$

for some constant  $C_\beta$  that depends only on the desired accuracy  $\beta$ .

This result is optimal: any reliable recovery method would require at least  $s \cdot \log n$  samples.

### Theorem (Candès & Tao)

Assume  $\bar{x} \in \mathbb{R}^n$  is  $s$ -sparse,  $n$  is prime, and we sample  $m$  Fourier coefficients. Then  $\bar{x}$  can be reconstructed from the  $m$  samples if  $m \geq 2s$ .

## Deterministic approach: restricted isometry

### Restricted isometry property (RIP)

Given  $A \in \mathbb{R}^{m \times n}$  and  $k \in \{1, \dots, m\}$ , the  $k$ -isometry constant  $\delta_k$  is the smallest  $\delta \geq 0$  such that

$$(1 - \delta)\|x\|_2^2 \leq \|Ax\|_2^2 \leq (1 + \delta)\|x\|_2^2$$

for all  $k$ -sparse  $x \in \mathbb{R}^n$ .

If  $\delta_k < 1$ , we say that  $A$  satisfies the RIP with constant  $\delta_k$ .

## Sampling of non-sparse signals

Given  $x \in \mathbb{R}^n$  and  $s < n$  define

$$x_s := \operatorname{argmin}_{\|z\|_0 \leq k} \|z - x\|_2$$

### Theorem (Candès & Plan)

Let  $x \in \mathbb{R}^n$ ,  $\beta > 0$  and  $\bar{s}$  be such that  $m \geq C_\beta \cdot \bar{s} \cdot \log n$ . Then with probability at least  $1 - 6/n - 6e^{-\beta}$  the  $\ell_1$  solution  $\hat{x}$  satisfies

$$\|\hat{x} - x\|_1 \leq \min_{1 \leq s \leq \bar{s}} C \cdot (1 + \alpha) \cdot \|x - x_s\|_1$$

for some constant  $C$  and

$$\alpha = \sqrt{\frac{(1 + \beta)s\mu(F) \log n \log m \log^2 s}{m}}$$

## Signal recovery with a RIP matrix

### Compressive sampling approach

measure

$$b_k = \langle a_k, \bar{x} \rangle, \quad k = 1, \dots, m \rightsquigarrow b = A\bar{x}$$

perform  $\ell_1$  recovery

$$\hat{x} := \operatorname{argmin}_x \{\|x\|_1 : Ax = b\}$$

### Theorem (Candès, Romberg, Tao)

Assume  $\delta_{2s} \leq \sqrt{2} - 1$ . Then the solution  $\hat{x}$  satisfies

$$\|\hat{x} - \bar{x}\|_1 \leq C \cdot \|\bar{x} - \bar{x}_s\|_1$$

and

$$\|\hat{x} - \bar{x}\|_2 \leq C \cdot \frac{\|\bar{x} - \bar{x}_s\|_1}{\sqrt{s}}$$

for some constant  $C$ .

## Matrices that satisfy RIP

With high probability an  $m \times n$  matrix  $A$  satisfies the RIP in the following cases:

- $m \gtrsim s \log(n/s)$  and  $A$  is Gaussian
- $m \gtrsim s \log(n/s)$  and  $A$  is binary
- $m \gtrsim s \log^4 n$  and  $A$  is partial DFT
- $m \gtrsim \mu(F)s \log^4 n$  and rows of  $A$  are iid  $F$ .

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## Sampling with noise

- Measurements in real life are generally noisy
- More appropriate model

$$b = A\bar{x} + z$$

noise term  $z \sim \mathcal{N}(0, \sigma^2 I)$

- Assume all columns of  $A$  have Euclidean norm equal to one.
- Modify  $\ell_1$  minimization to account for noise

### Lasso

$$\min \frac{1}{2} \|b - Ax\|_2^2 + \lambda \cdot \sigma \cdot \|x\|_1$$

### Dantzig selector

$$\min \|x\|_1 \quad \text{subject to} \quad \|A^*(b - Ax)\|_\infty \leq \lambda \cdot \sigma$$

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## Noise aware recovery (random sensing)

### Theorem (Candès & Plan)

Let  $\bar{x} \in \mathbb{R}^n$ ,  $\beta > 0$  and  $\bar{s}$  be such that  $m \geq C_\beta \cdot \bar{s} \cdot \log n$ . Then with probability at least  $1 - 6/n - 6e^{-\beta}$  the solution  $\hat{x}$  to the Lasso or the Dantzig selector with  $\lambda = 10\sqrt{\log n}$  satisfies

$$\|\hat{x} - \bar{x}\|_1 \leq \min_{1 \leq s \leq \bar{s}} C \cdot (1 + \alpha^2) \cdot \left( \|\bar{x} - \bar{x}_s\|_1 + s\sigma \sqrt{\frac{\log n}{m}} \right)$$

for some constant  $C$ .

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## Optimality of Compressive Sensing

### Back to Gaussian sensing

$m$  Gaussian measurements and  $\ell_1$  decoding:

$$\|\hat{x} - x\|_2 \lesssim \frac{\|x - x_s\|_1}{\sqrt{s}}, \quad s \approx \frac{m}{\log(n/m)}.$$

### Question

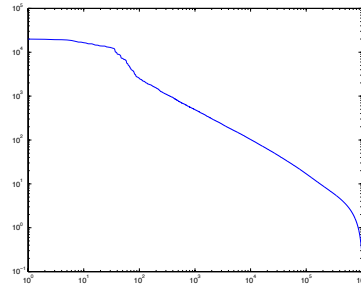
Can we do better with other measurements or other algorithms?

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## Signals with power law



image



sorted wavelet coefficients

Power law decay:  $|x|_{(1)} \geq |x|_{(2)} \geq \dots \geq |x|_{(n)}$

$$|x|_{(k)} \leq \frac{C}{k^p}$$

### Model

$\ell_p$  ball  $\mathcal{B}_p := \{x : \|x\|_p \leq 1\}$ .

Discuss case  $p = 1$  but same discussion applies to  $0 \leq p \leq 1$ .

## Gelfand widths

### Theorem (Donoho)

$$d_m(\mathcal{B}) \leq E_m(\mathcal{B}) \leq C \cdot d_m(\mathcal{B}),$$

where  $d_m(\mathcal{B})$  is the  $m$ -width of  $\mathcal{B}$ :

$$d_m(\mathcal{B}) := \inf_V \left\{ \sup_{x \in \mathcal{B}} \|P_V x\|_2 : \text{codim}(V) < m \right\}$$

### Theorem (Kashin, Garnaev-Gluskin)

For  $\ell_1$  ball

$$C_1 \cdot \sqrt{\frac{\log(n/m) + 1}{m}} \leq d_m(\mathcal{B}) \leq C_2 \cdot \sqrt{\frac{\log(n/m) + 1}{m}}.$$

Compressive sensing achieves the limits of performance.

## Recovery of $\ell_1$ ball $\mathcal{B}$

### Gaussian sensing

- Suppose unknown vector is in  $\mathcal{B}$
- Take  $m$  Gaussian measurements

$$\|\hat{x} - x\|_2 \lesssim \sqrt{\frac{\log(n/m) + 1}{m}}.$$

### Ideal sensing

Best we can hope from  $m$  linear measurements:

$$E_m(\mathcal{B}) = \inf_{D, E} \sup_{x \in \mathcal{B}} \|x - D(F(x))\|.$$

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## References for today's material

- Slides for this minicourse:  
<http://andrew.cmu.edu/user/jfp/UNencuentro>
- E. Candès & Y. Plan, "A probabilistic and RIPless theory of compressed sensing," *IEEE Trans. Inf. Theory*, vol 57, no. 11, pp. 7235–7254, Nov. 2011.
- E. Candès, J. Romberg, & T. Tao, "Robust uncertainty principles: Exact signal reconstruction from highly incomplete frequency information," *IEEE Trans. Inf. Theory*, vol 52, no. 2, pp. 489–509, Feb. 2006.
- E. Candès & T. Tao, "Decoding by linear programming," *IEEE Trans. Inf. Theory*, vol 51, no. 12, pp. 4203–4215, Dec. 2005.
- D. Donoho, "Compressed sensing," *IEEE Trans. Inf. Theory*, vol 52, no. 4, pp. 1289–1306, April 2006.
- Papers on compressive sensing: <http://dsp.rice.edu/cs>

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