Artificial Intelligence Methods for Social Good M2-I [Game Theory]: Basics of Game Theory

> 08-537 (9-unit) and 08-737 (12-unit) Instructor: Fei Fang <u>feifang@cmu.edu</u> Wean Hall 4126

## Quiz I: Recap: Optimization Problem

- Given coordinates of n residential areas in a city (assuming 2-D plane), denoted as x<sup>1</sup>, ..., x<sup>n</sup>, the government wants to find a location that minimizes the sum of (Euclidean) distances to all residential areas to build a hospital. The optimization problem can be written as
  - A:  $\min_{x} \sum_{i} |x^{i} x|$
  - B:  $\min_{x} \sum_{i} ||x^{i} x||_{2}$
  - C:  $\min_{x} \sum_{i} (x^{i} x)^{2}$
  - D: none of above

#### From Games to Game Theory



- The study of mathematical models of conflict and cooperation between intelligent decision makers
- Used in economics, political science etc







#### Outline

- Basic Concepts in Games
- Basic Solution Concepts
- Compute Nash Equilibrium
- Compute Strong Stackelberg Equilibrium

## Learning Objectives

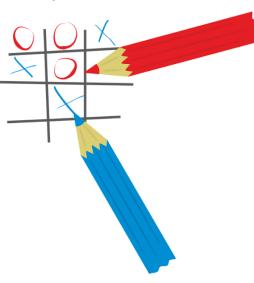
- Understand the concept of
  - Game, Player, Action, Strategy, Payoff, Expected utility, Best response
  - Dominant Strategy, Maxmin Strategy, Minmax Strategy
  - Nash Equilibrium
  - Stackelberg Equilibrium, Strong Stackelberg Equilibrium
- Describe Minimax Theory
- Formulate the following problem as an optimization problem
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#### Let's Play! Classical Games

- Exp I: Rock-Paper-Scissors (RPS)
  - Rock beats Scissors
  - Scissors beats Paper
  - Paper beats Rock
- Exp 2: Prisoner's Dilemma (PD)
  - If both Cooperate: I year in jail each
  - If one Defect, one Cooperate: 0 year for (D), 3 years for (C)
  - If both Defect: 2 years in jail each

#### Let's Play! Classical Games

- Exp 3: Battle of Sexes (BoS)
  - ▶ If football together: Alex ☺☺, Berry ☺
  - ▶ If concert together: Alex ☺, Berry ☺☺
  - ▶ If not together: Alex ☺, Berry ☺
- Tic-Tac-Toe (TTT)



- Def I: Game
  - Players
  - Actions
  - Payoffs

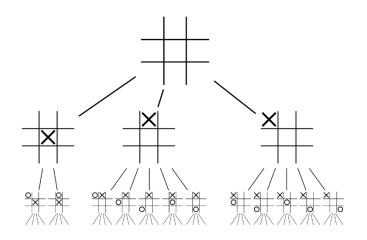
#### Representation

- Normal form (Matrix form, Strategic form, Standard form)
  - Move simultaneously
  - Bimatrix game (Two-player)
    Exp 2: PD

	Player 2		
		Cooperate	Defect
Player I	Cooperate	-1,-1	-3,0
	Defect	0,-3	-2,-2

#### Extensive form

- Timing, Sequence of move
- Game tree
- Information
- Natural description



- Pure Strategy
  - Choose one action deterministically
- Def 2: Mixed Strategy
  - Play randomly
  - Support: chosen with non-zero probability
- Def 3: Expected utility
  - Average utility weighted by probability

#### Quiz 2: Basic Concepts in Games

In Exp I (Rock-Paper-Scissors), if  $s_1 = (\frac{1}{3}, \frac{2}{3}, 0), s_2 = (0, \frac{1}{2}, \frac{1}{2}), \text{ what is } u?$  u = (0,0)  $u = (-\frac{1}{3}, \frac{1}{3})$   $u = (-\frac{1}{2}, \frac{1}{3})$   $u = (-\frac{1}{2}, \frac{1}{2})$   $u = (-\frac{1}{2}, \frac{1}{2})$ 

		Rock	Paper	Scissors
/er	Rock	0,0	-1,1	١,-١
Player	Paper	١,-١	0,0	-1,1
	Scissor	-1,1	۱,-۱	0,0

- Def 4: Best Response
  - Set of actions or strategies leading to highest expected utility
- Thm I (Nash 1951): A mixed strategy is BR iff all actions in the support are BR

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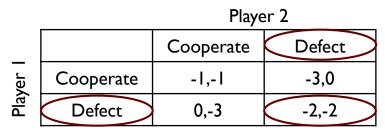
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- How should a player play?
- Def 5: Dominant Strategy
  - One strategy is always better/never worse/never worse and sometimes better than any other strategy
  - Focus on single player's strategy
  - Exp 2: PD

	Player 2		
		Cooperate	Defect
/er	Cooperate	-1,-1	-3,0
Player	Defect	0,-3	-2,-2

- Dominant strategy equilibrium/solution
  - Every player plays a dominant strategy
  - Focus on strategy profile for all players

- Def 6: Nash Equilibrium
  - Every player's strategy is a best response to others' strategy profile
  - Focus on strategy profile for all players
  - One cannot gain by unilateral deviation
  - Pure Strategy Nash Equilibrium (PSNE)
  - Mixed Strategy Nash Equilibrium (MSNE)
  - Exp 2: PD



- Thm 2 (Nash 1951): NE always exists in finite games
  - Brouwer's fixed point theorem

- Def 7: Maxmin Strategy (applicable to multiplayer games)
  - Maximize worst case expected utility
  - Focus on single player's strategy
- Def 8: Minmax Strategy (two-player games only)
  - Minimize best case expected utility for the other player (just want to harm your opponent)
  - Focus on single player's strategy

- Thm 3: (Minimax Theorem, von Neumann 1928) Minmax=Maxmin in 2-player zero-sum games
- Further, Minmax=Maxmin=NE (Nash 1951)

- Exp 4: Power of Commitment
  - NE utility=(2,1)
  - If leader (player I) commits to playing b, then player has to play d, leading to a utility of 3 for leader
  - If leader (player 1) commits to playing a and b uniformly randomly, then player still has to play d, leading to a utility of 3.5 for leader

	Player 2		
_		С	d
Player	а	2,1	4,0
Ы	b	١,0	3,2

- Def 9: Best Response Function
  - A mapping from a strategy of one player to a strategy of another player in the best response set
- Def 10: Stackelberg Equilibrium
  - Leader vs follower game
  - Leader commits to a strategy
  - Follower responds according a best response function
  - Focus on strategy profile for all players

#### Def I I: Strong Stackelberg Equilibrium (SSE)

Follower breaks tie in favor of the leader

- Def I I: Strong Stackelberg Equilibrium (SSE)
  - Follower breaks tie in favor of the leader
  - Leader can induce the follower to do so by perturbing the strategy in the right direction
  - SSE always exist in two-player finite games

### Quiz 3: Basic Solution Concepts

- What is the relationship of leader's expected utility in SSE and NE in two-player games?
  - $u^{SSE} \ge u^{NE}$
  - $u^{SSE} = u^{NE}$
  - $u^{SSE} \leq u^{NE}$
  - none of the above

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- Find pure strategy Nash Equilibrium (PSNE)
  - Enumerate all action profile
  - Check if no incentive to deviate

- Find all Nash Equilibrium (two-player)
  - Special case: Zero-sum game
    - Polynomial time solvable (minmax or maxmin LP)
  - General case
    - PPAD-Complete (Chen & Deng, 2006)
      - □ Unlikely to have polynomial time algorithm
      - □ Conjecture: slightly easier than NP-Complete problems

#### Find all Nash Equilibrium (two-player)

- Support Enumeration Method
  - Enumerate support pair
  - For each possible support pair
    - Compute the probability so as to keep the other player indifferent among actions in the support
    - □ Check if no incentive to deviate
    - □ Or combine the two steps: solve an LP

- Exp 3 (Battle of Sexes)
  - Enumerate support pair
    - Support size=1: PSNE!
    - Support size=2:Alex: (Football, Concert), Berry: (Football, Concert)
    - Different support size: no NE exist (#constraints>#variables)
  - Compute the probability so as to keep the other player indifferent among actions in the support
  - Why? (check Thm I)

	Berry		
		Football	Concert
Alex	Football	2,1	0,0
A	Concert	0,0	١,2

## Quiz 4: Compute Nash Equilibrium

- What is the probability of Berry choosing Football in NE of Exp 3 (Battle of Sexes) with support size=2?
  - 01  $\frac{1}{3}$  $\frac{2}{3}$

		Berry	
		Football	Concert
Alex	Football	2,1	0,0
A	Concert	0,0	١,2

## Find all Nash Equilibrium (two-player)

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#### Lemke-Howson Algorithm

- Linear Complementarity (LCP) formulation (another special class of optimization problem)
- Solve by pivoting on support (similar to Simplex algorithm)
- In practice, available solvers/packages: nashpy (python), gambit project (<u>http://www.gambit-project.org/</u>)

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## Compute Strong Stackelberg Equilibrium

- Find Strong Stackelberg Equilibrium (not restricted to pure strategy)
  - Special case (zero-sum): SSE=NE=Minmax=Maxmin
  - General case: Polynomial time solvable
    - Multiple Linear Programming
    - Mixed Integer Linear Programming
      - □ Exponential in theory
      - □ Efficient in practice in some cases
  - Solvers available? Course project idea (9 units only): develop a SSE solver package in Python

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## • Key take-away:

- There are various solution concepts in games
- In two-player zero-sum games, many solution concepts lead to same strategy and utility
- Finding NE or SSE can be formulated as one or more optimization problem

# Summary

Solution Concepts	Key Algorithm In Class
Minmax/Maxmin	LP
Nash Equilibrium	LP for zero-sum, Support enumeration for general-sum
Strong Stackelberg Equilibrium	LP for zero-sum, multiple LP or MILP for general-sum

#### Game Theory: Additional Resources

- Text book
  - Algorithmic Game Theory 1st Edition, Chapters 1-3
  - Noam Nisan (Editor), Tim Roughgarden (Editor), Eva Tardos (Editor), Vijay V.Vazirani (Editor)
  - http://www.cs.cmu.edu/~sandholm/cs15-892F13/algorithmicgame-theory.pdf
- Online course
  - https://www.youtube.com/user/gametheoryonline