

MULTI-RESOURCE ALLOCATION

FAIRNESS-EFFICIENCY TRADEOFFS IN A UNIFYING FRAMEWORK

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What is Fairness?

- Politics, economics, sociology, engineering...

How do you allocate a resource to different users?

- Variance, Jain's index, entropy (see TR for references)...
- Isoelastic or α -fairness
- Unifying axiomatic theory of decomposable fairness measures

$$\text{sgn}(1 - \beta) \left(\sum_{i=1}^n \left(\frac{x_i}{\sum_{j=1}^n x_j} \right)^{1-\beta} \right)^{\frac{1}{\beta}} \left(\sum_{i=1}^n x_i \right)^{\lambda}$$

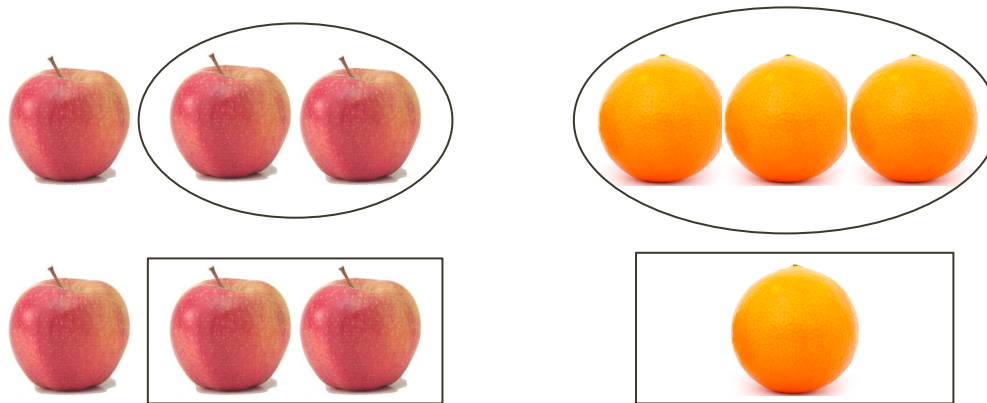
Our Question

- Suppose you have **multiple** non-substitutable resources.
 - Memory
 - CPU
 - Bandwidth
- They combine to make something...
 - Jobs in a datacenter
- that multiple people want.
 - Different bundles of resource requirements
- But the resources are finite.



How do you allocate the resources to different people?

Two-Resource Example



Generalized Fairness on Jobs (GFJ)

$$f_{\beta, \lambda}(\mathbf{x}) = \text{sgn}(1 - \beta) \left(\sum_{i=1}^n \left(\frac{x_i}{\sum_{j=1}^n x_j} \right)^{1-\beta} \right)^{\frac{1}{\beta}} \left(\sum_{i=1}^n x_i \right)^{\lambda}$$

Fairness

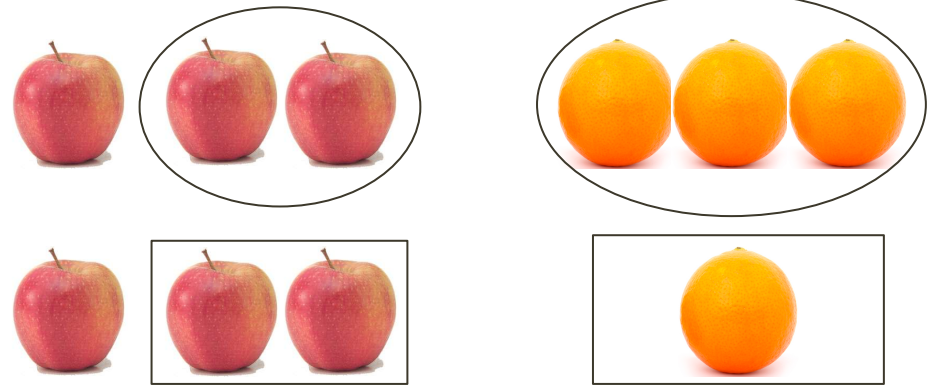
Efficiency

- Unique family of functions: β and λ parameters
 - β : type of fairness
 - λ : importance of efficiency

Defining “Fairness”

- An equal allocation?

- 1 job for each user



- But not efficient

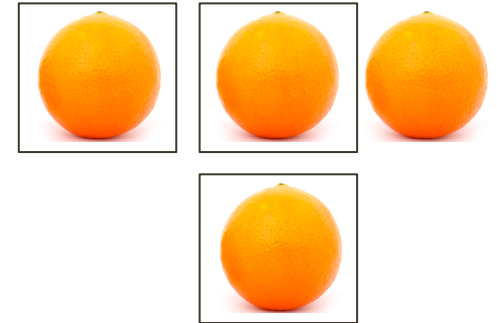
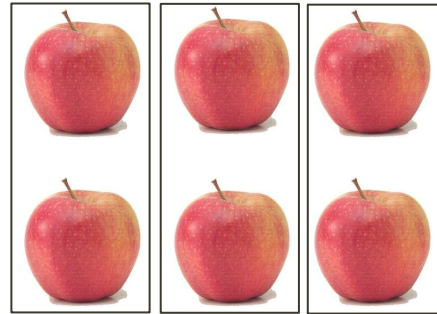
- Ranking the fairness of different allocations

$$\text{sgn}(1 - \beta) \left(\sum_{i=1}^n \left(\frac{x_i}{\sum_{j=1}^n x_j} \right)^{1-\beta} \right)^{\frac{1}{\beta}}$$

Defining “Efficiency”

- Maximize the total number of jobs?

- 0 jobs to user 1
- 3 jobs to user 2

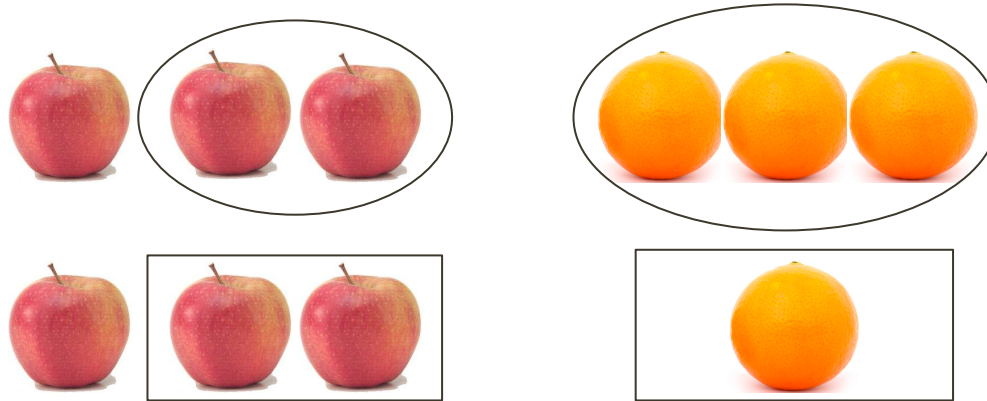


- But not that fair

- Ranking the efficiency of different allocations

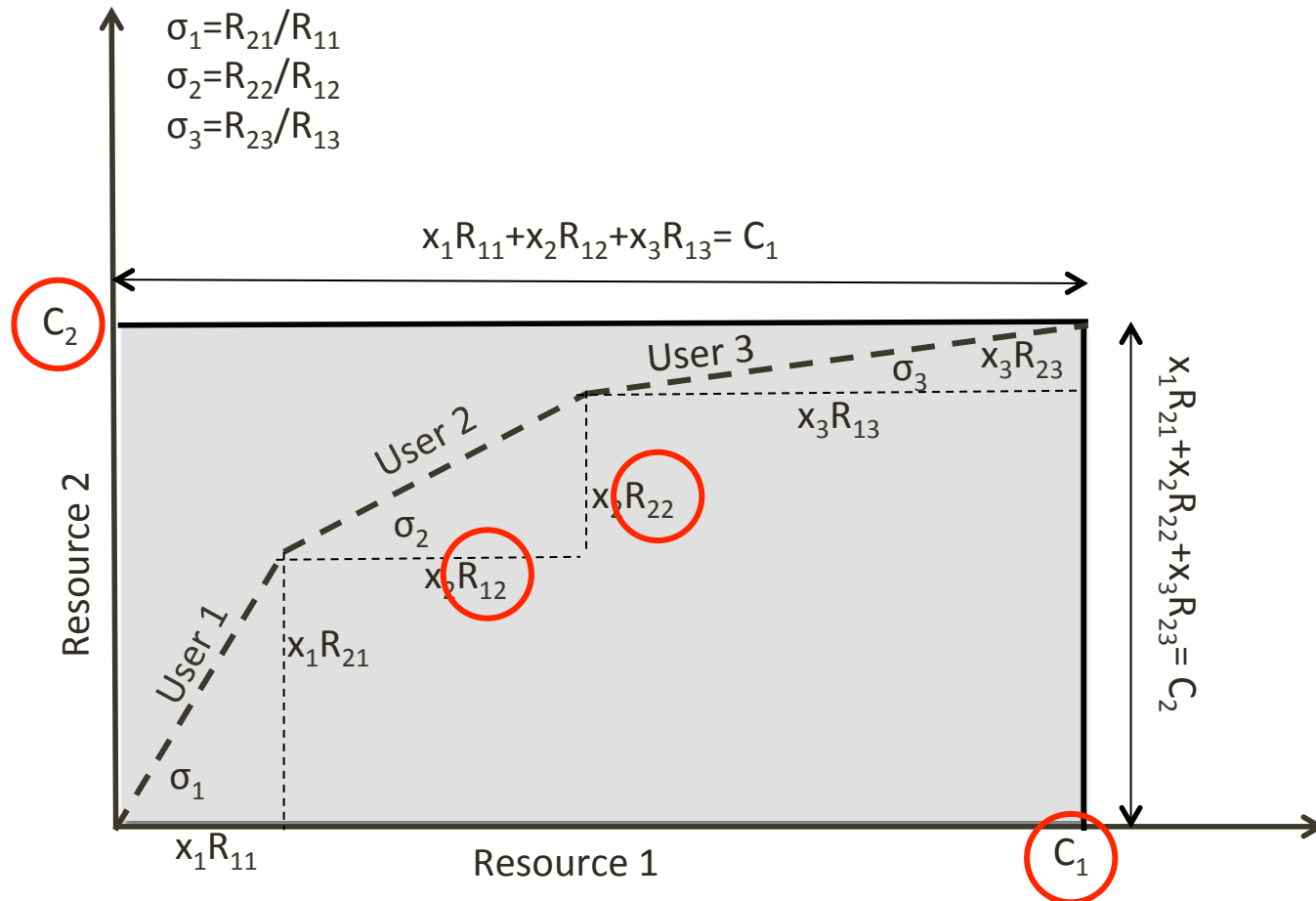
$$\left(\sum_{i=1}^n x_i \right)^\lambda$$

Heterogeneous Users



- Different users need different mixes of resources...
- Is it fair to treat them the same way?

Visualizing Heterogeneity



3 Users, 2 Resources

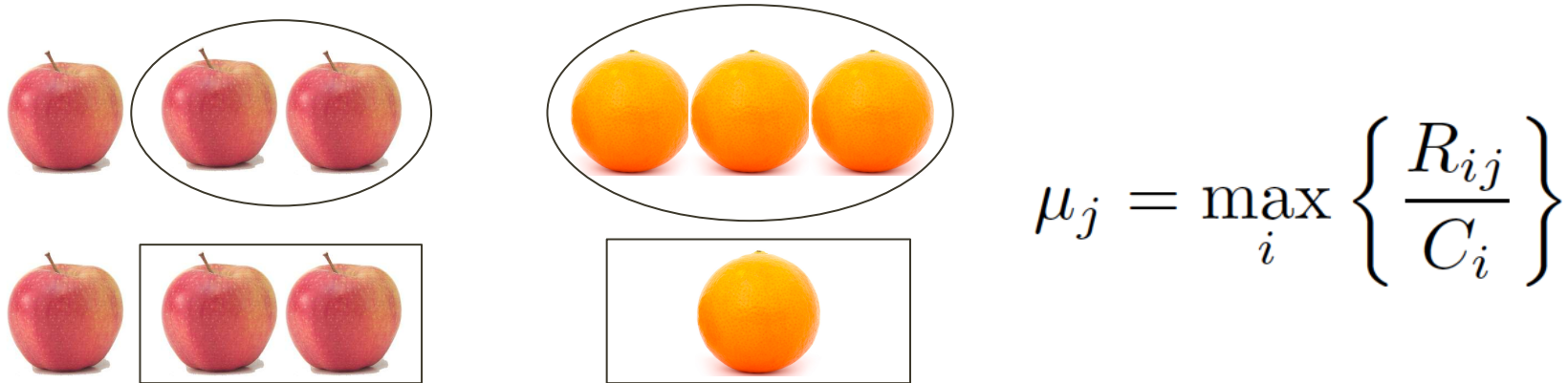
Dominant Shares

- Dominant shares $\mu_j x_j$ for each user

$$\mu_j = \max_i \left\{ \frac{R_{ij}}{C_i} \right\}$$

- Maximum share of any resource

Calculating Dominant Shares



$$\mu_j = \max_i \left\{ \frac{R_{ij}}{C_i} \right\}$$

- 2 of 6 apples and 3 of 4 oranges: $\mu_1 = \max \left(\frac{1}{3}, \frac{3}{4} \right)$
- 2 of 6 apples and 1 of 4 oranges: $\mu_2 = \max \left(\frac{1}{3}, \frac{1}{4} \right)$

Fairness on Dominant Shares (FDS)

- Use dominant shares instead of number of jobs
- If μ is larger, equal dominant shares for smaller number of jobs

$$\text{sgn}(1 - \beta) \left(\sum_{j=1}^n \left(\frac{\mu_j x_j}{\sum_{k=1}^n \mu_k x_k} \right)^{1-\beta} \right)^{\frac{1}{\beta}} \left(\sum_{j=1}^n \mu_j x_j \right)^{\lambda}$$

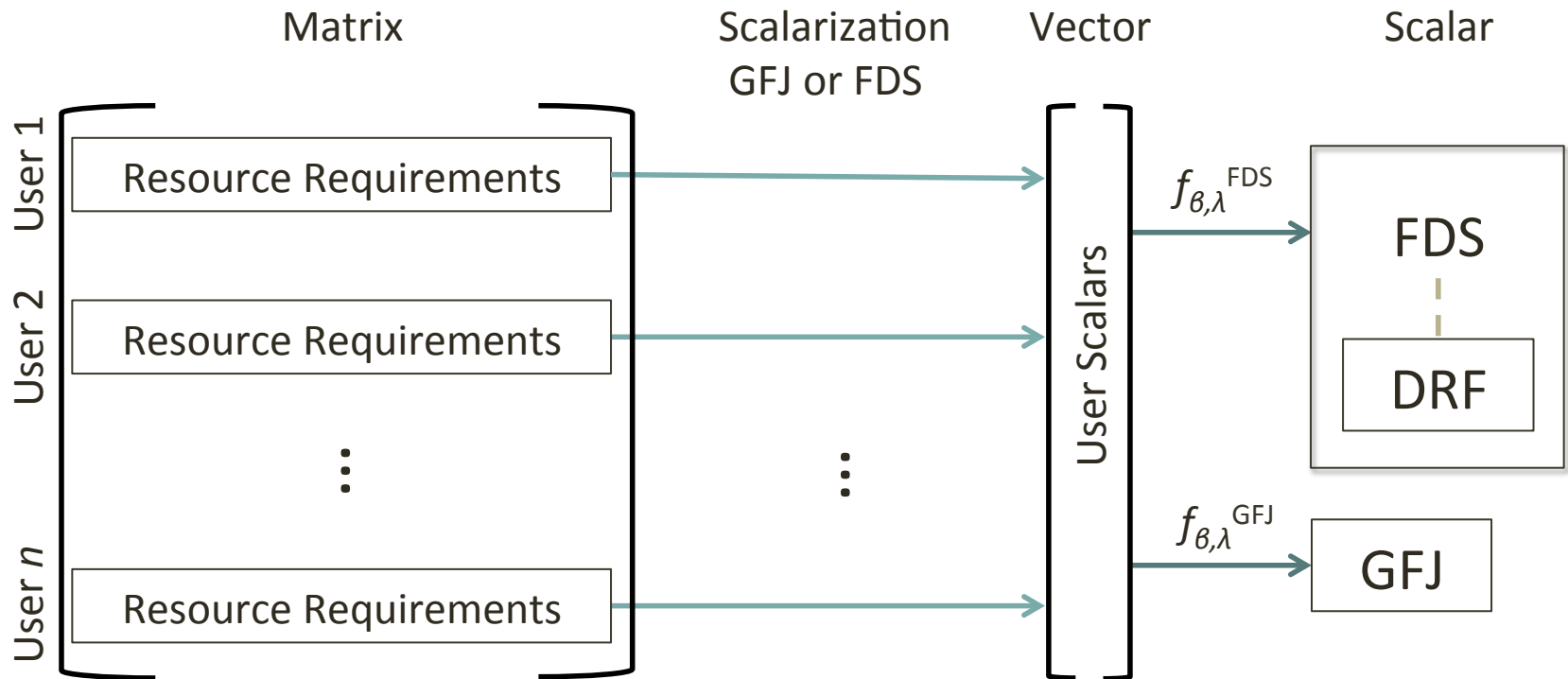
GFJ

Generalized Fairness on Jobs

FDS

Fairness on Dominant Shares

Resource Scalarization



Desirable Properties Of Fairness Functions

Property 1: Pareto-Efficiency

- $f(x) > f(y)$ whenever the allocation x Pareto-dominates y .
 - $x_i \geq y_i$ for all entries i , with strict inequality for some i
- Not an axiom: needs to be proven
- Does not hold for all parameter combinations

Parameter Conditions

- Necessary and sufficient conditions

$$|\lambda| \geq \left| \frac{1-\beta}{\beta} \right| \quad \beta > 0$$

- Holds for FDS and GFJ
 - Comes from the same conditions for single-resource fairness
- If $\lambda = \frac{1-\beta}{\beta}$ and $\beta > 0$, fairness becomes α -fairness with $\alpha = \beta$.

Property 2: Sharing Incentive

- Each user receives at least a $\frac{1}{n}$ share of some resource.
 - Dominant share is over $\frac{1}{n}$
- Users don't want to share the resources equally.
- Does it hold?

Parameter Conditions

- Sufficient conditions:

$$\text{FDS} \quad \lambda = \frac{1-\beta}{\beta} \quad \beta > 1$$

- Counterexamples exist:

$$\lambda = \frac{1-\beta}{\beta} \quad \text{FDS} \quad 0 < \beta < 1 \quad \text{GFJ} \quad \beta > 0$$

Property 3: Envy-Freeness

- A user can process more jobs with his own rather than another user's resource allocation.
 - Users don't want to switch allocations.
- Does it hold?

Parameter Conditions

- Sufficient conditions:

$$\text{FDS} \quad \lambda = \frac{1-\beta}{\beta} \quad \beta > 1$$

- Counterexamples exist:

$$\lambda = \frac{1-\beta}{\beta} \quad \text{FDS} \quad 0 < \beta < 1 \quad \text{GFJ} \quad \beta > 0$$

Sufficient Conditions

Fairness	Pareto-Efficiency	Sharing Incentive	Envy-Freeness
FDS	$ \lambda \geq \left \frac{1-\beta}{\beta} \right , \beta > 0$	$\lambda = \frac{1-\beta}{\beta}, \beta > 1$ $\lambda = 0, \text{ any } \beta$	$\lambda = \frac{1-\beta}{\beta}, \beta > 1$ $\lambda = 0, \text{ any } \beta$
GFJ	$ \lambda \geq \left \frac{1-\beta}{\beta} \right , \beta > 0$	–	–

Existence of a Counterexample

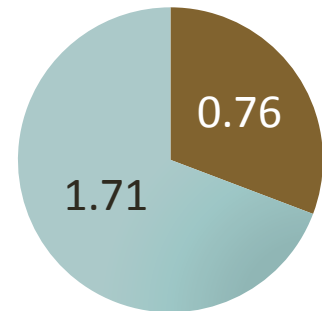
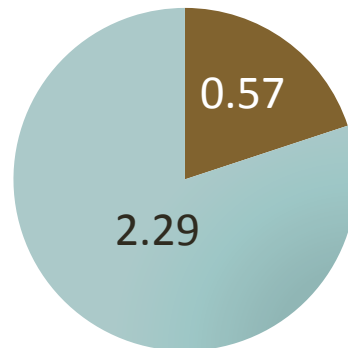
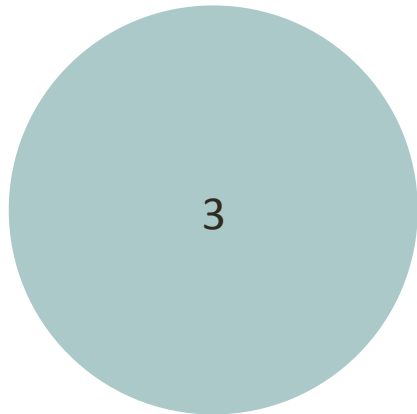
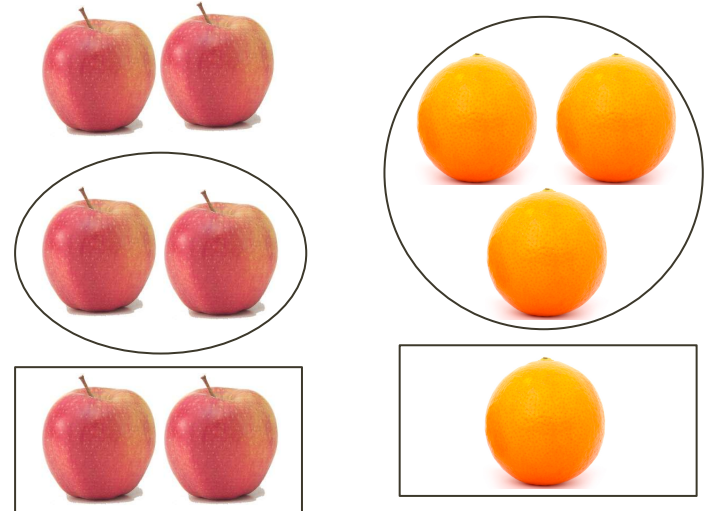
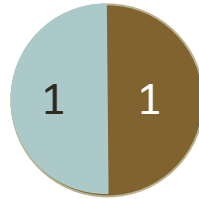
Fairness	Sharing Incentive		Envy-Freeness	
FDS	$\lambda = \frac{1-\beta}{\beta}, 0 < \beta < 1$	$\lambda = \infty, \text{ any } \beta$	$\lambda = \frac{1-\beta}{\beta}, 0 < \beta < 1$	$\lambda = \infty, \text{ any } \beta$
GFJ	$\lambda = \frac{1-\beta}{\beta}, \beta > 0$ $ \lambda < \frac{ 1-\beta }{\beta}, \beta > 1$	$\lambda = \infty \text{ or } 0, \text{ any } \beta$ $ \lambda > \frac{ 1-\beta }{\beta}, 0 < \beta < 1$	$\lambda = \frac{1-\beta}{\beta}, \beta > 0$ $ \lambda < \frac{ 1-\beta }{\beta}, \beta > 1$	$\lambda = \infty \text{ or } 0, \text{ any } \beta$ $ \lambda > \frac{ 1-\beta }{\beta}, 0 < \beta < 1$

What about Efficiency?

Fair, Efficient, or Both?

User 1

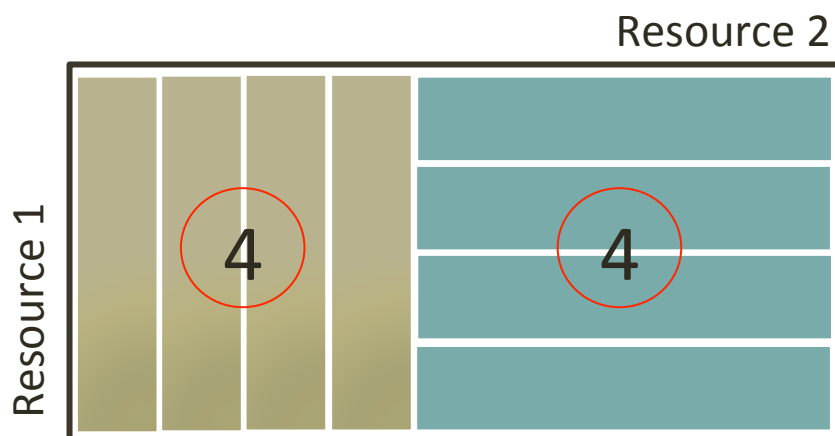
User 2



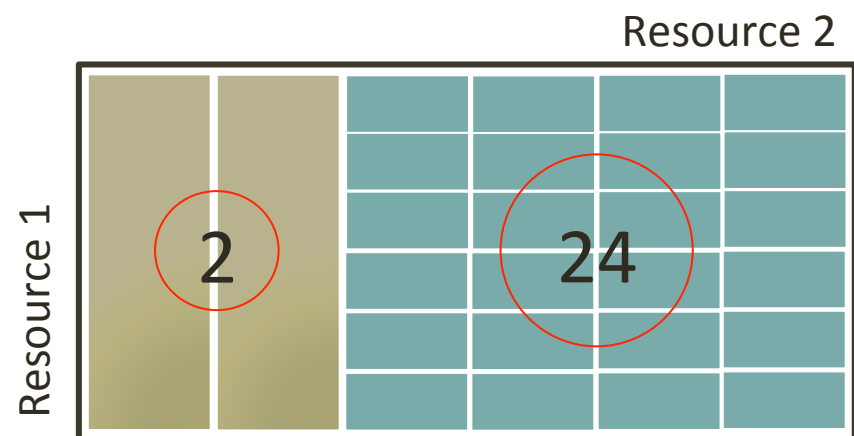
Existence of a Tradeoff

- Nonlinear, non-separable, multidimensional, continuous state-space knapsack problem
 - Maximize fairness function subject to multiple linear capacity constraints
 - Allow fractional jobs

No Tradeoff



Tradeoff Exists



Equal Allocations at Maximum Efficiency

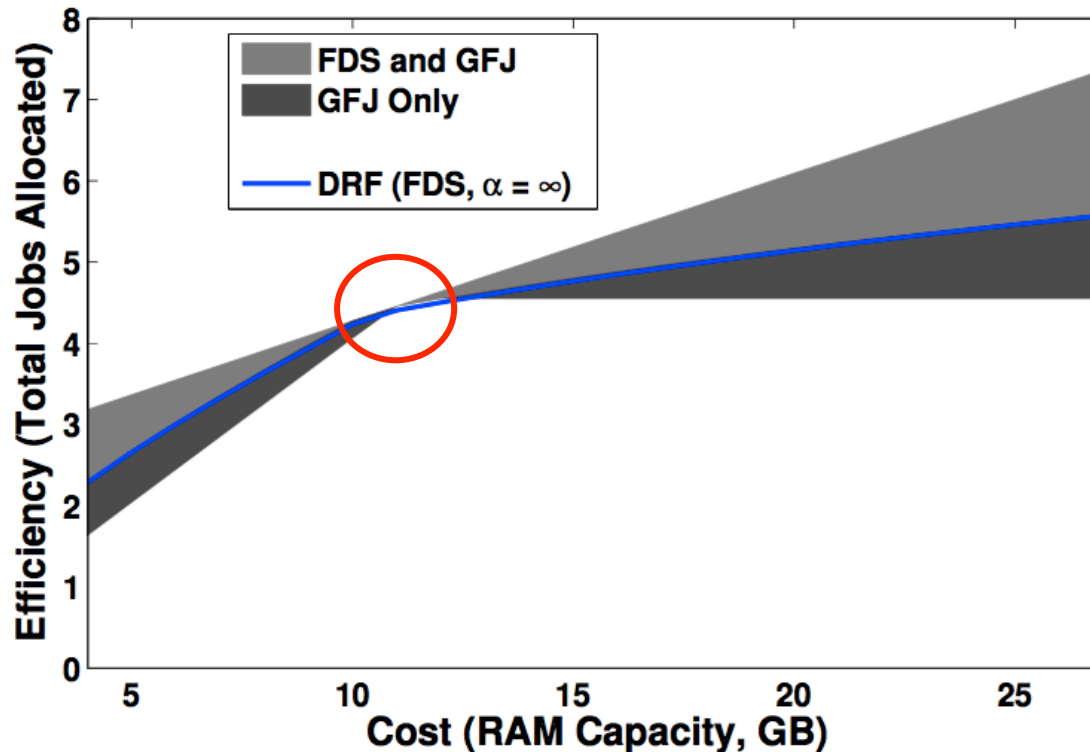
- Number of tight resource constraints = number of users

$$\sum_{j=1}^n \gamma_{ij} x_j \leq 1 \quad \forall i.$$

$$\text{FDS} \quad \sum_{j=1}^n \frac{\gamma_{ij}}{\mu_j} = \rho$$

$$\text{GFJ} \quad \sum_{j=1}^n \gamma_{ij} = r$$

Efficiency Operating Range



$$x + 3y \leq 9 \quad \text{CPUs}$$

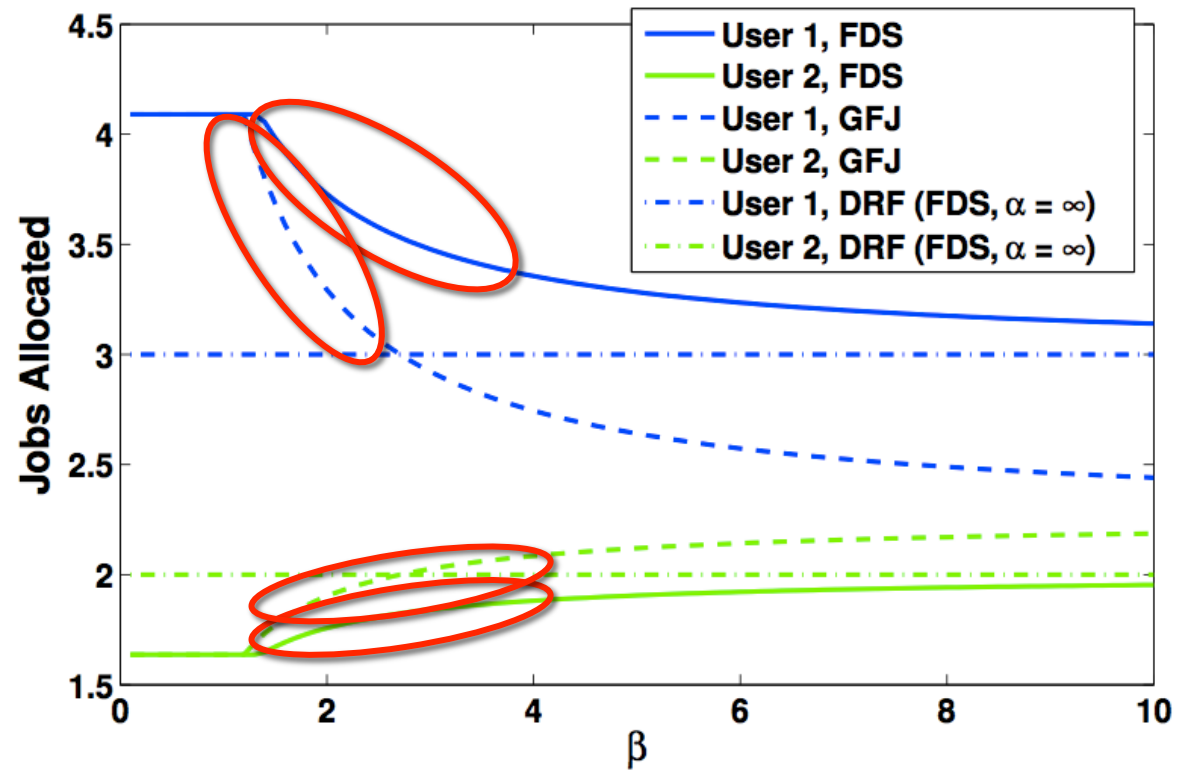
$$4x + y \leq 18 \quad \text{GB}$$

Optimal allocations for a range of β and λ

Dominant Resource Fairness (DRF): max-min fairness on dominant shares

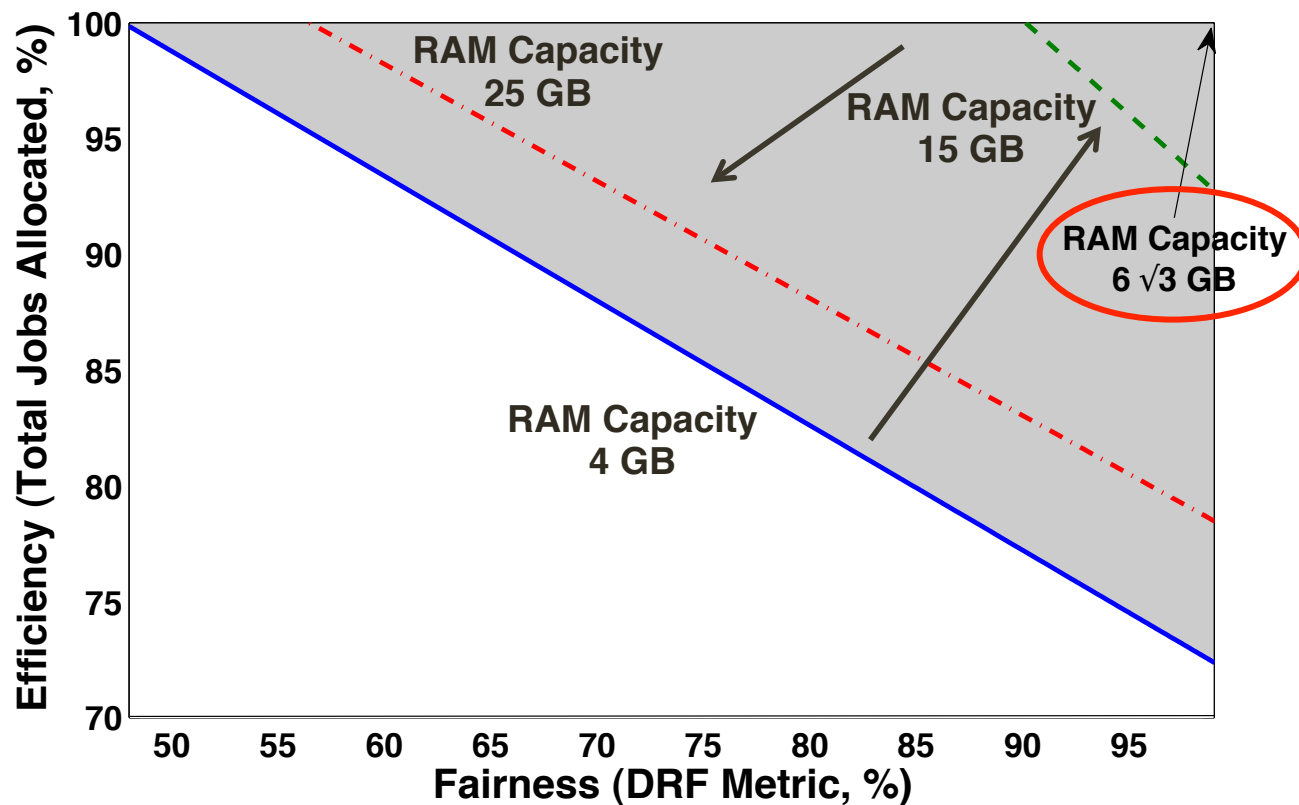
Job Allocation

$$\lambda = \frac{1-\beta}{\beta}$$



Optimal allocations for $\alpha = \beta$ -fairness

Numerical Example



- Fairness: DRF-fairness divided by maximal DRF value
- Efficiency: Total jobs divided by maximum number of jobs

Psychological Perceptions

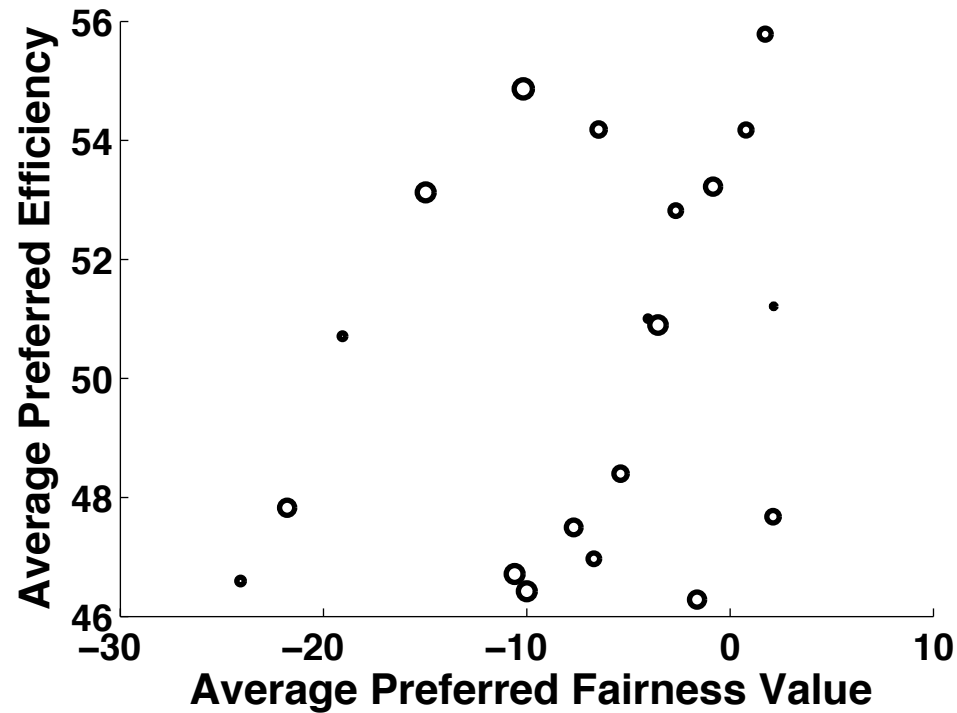
C. Joe-Wong, et al. Multi-Resource Allocation: Fairness-Efficiency Tradeoffs in a Unifying Framework.
Tech report, available <http://www.princeton.edu/~chiangm/multiresourcefairness.pdf>

Parameter Values

- What parameters are compatible with the responses?
 - Do they satisfy Pareto-efficiency, etc.?
- Do people agree with each other?
- Online survey asking people to rank datacenter allocations

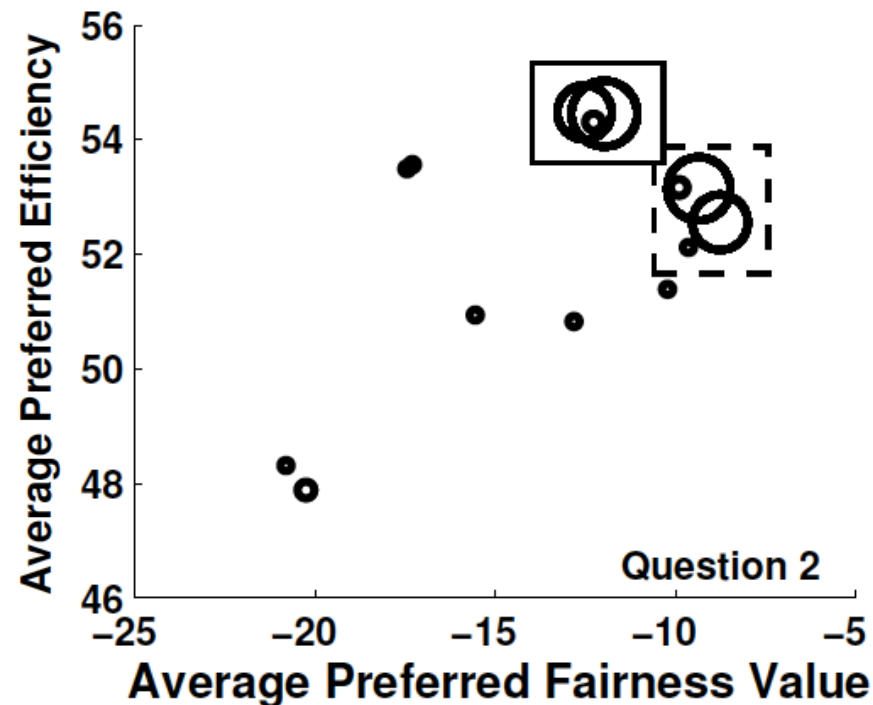
Allocation Options	Allocated to Client A			Allocated to Client B			Total no. of Jobs Completed
	CPU	TB	No. of Jobs Completed for Client A	CPU	TB	No. of Jobs Completed for Client B	
Allocation 1	24	96	24	84	28	28	52

Are People Very Different?

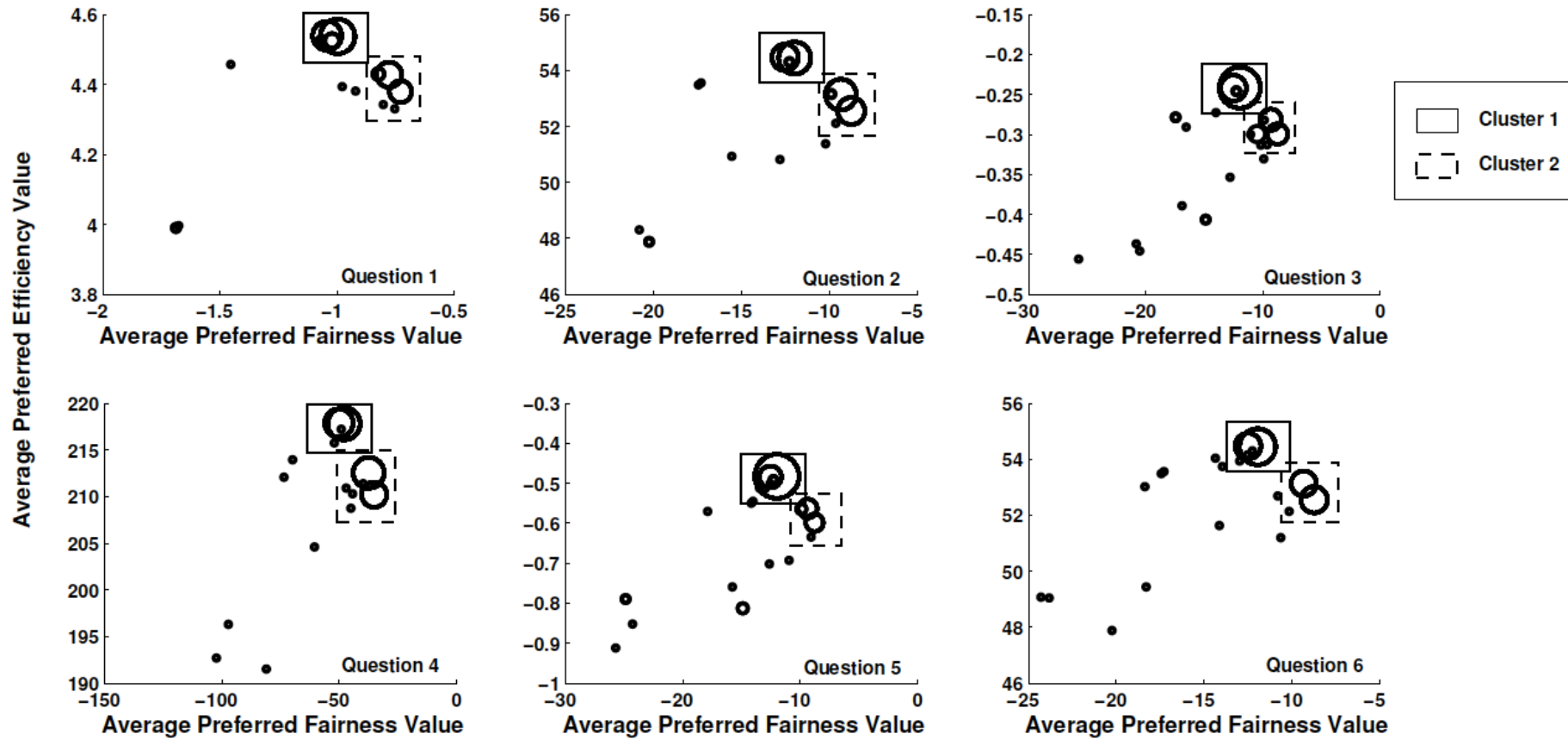


Actually They're Pretty Similar

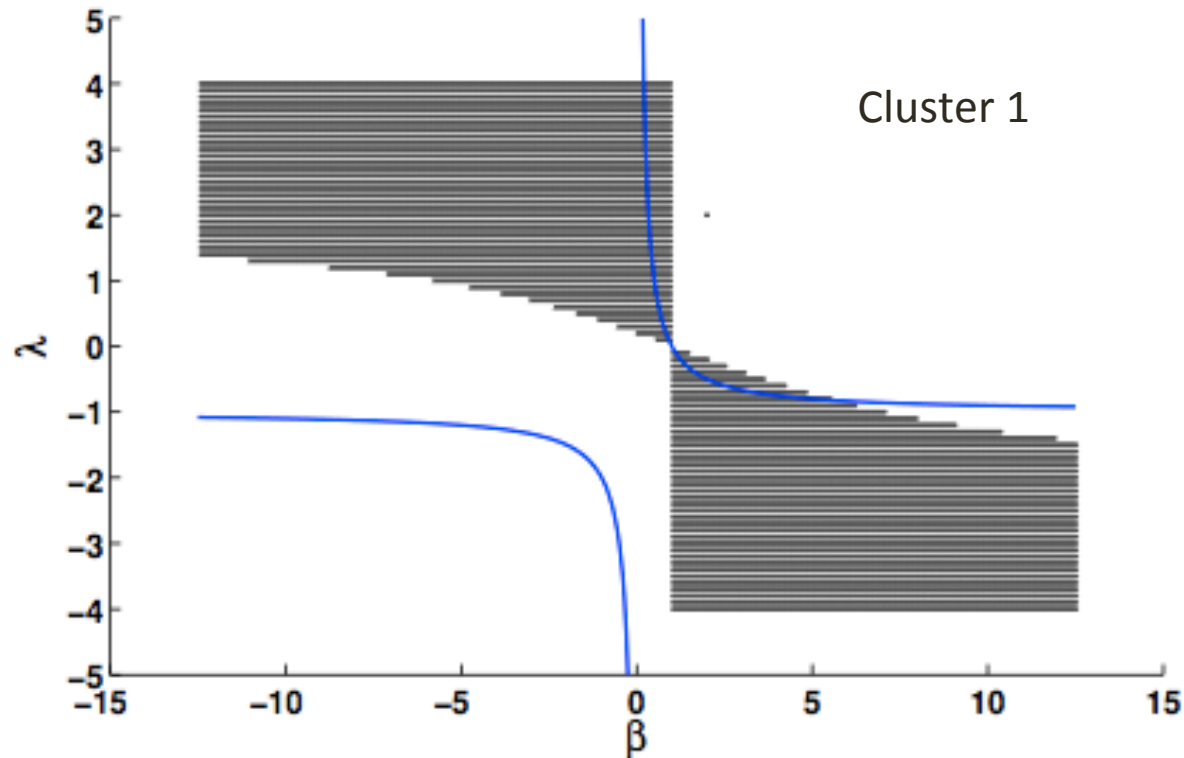
- Cluster 1 prefers efficiency to fairness
- Cluster 2 prefers fairness to efficiency



All Responses



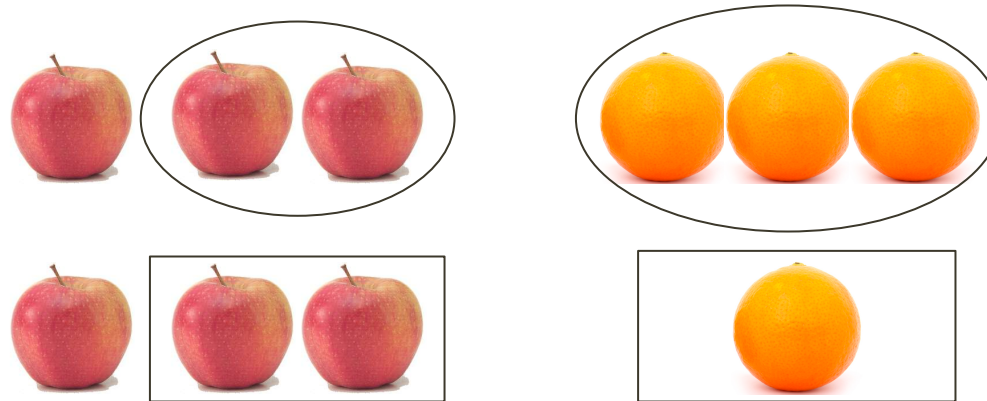
Compatible Parameters (GFJ)



The darker the square, the more participant rankings were compatible.

Lines represent Pareto-efficient frontiers.

Back to the Motivation



Questions Answered

- How do we define fairness?
 - GFJ and FDS
- Are these properties satisfied?
 - Pareto-efficiency
 - Envy-freeness
 - Sharing incentive
- Does a fairness-efficiency tradeoff exist?
- What parameters are consistent with actual preferences?
 - Users fall into 2 clusters



Photo: Cam Barker

Thank you!
Questions?
