

### Homework 8

1. Complete the proof from the text of Kripke completeness for the positive fragment of IPC as follows:
  - (a) Show that for any poset  $I$ , the exponential poset  $\mathbf{2}^I$  is a Heyting algebra. (Hint: the limits and colimits are “pointwise”, and the Heyting implication  $p \Rightarrow q$  is defined at  $i \in I$  by  $(p \Rightarrow q)(i) = \top$  iff for all  $j \leq i, p(j) \leq q(j)$ ).
  - (b) Show that for any poset CCC  $\mathbf{A}$ , the map  $y : \mathbf{A} \rightarrow \mathbf{2}^{\mathbf{A}^{\text{op}}}$  defined in the text is indeed (i) monotone, (ii) injective, and (iii) preserves CCC structure.
2. Verify the claim in the text that the products  $A \times B$  in categories  $\mathbf{Sets}^I$  of  $I$ -indexed sets ( $I$  a poset) can be computed “pointwise”. Show, moreover, that the same is true for all limits and colimits.
3. Prove that every functor  $F : \mathbf{C} \rightarrow \mathbf{D}$  can be factored as

$$\mathbf{C} \xrightarrow{E} \mathbf{E} \xrightarrow{M} \mathbf{D}$$

in the following two ways:

- (a)  $E : \mathbf{C} \rightarrow \mathbf{E}$  is bijective on objects and full, and  $M : \mathbf{E} \rightarrow \mathbf{D}$  is faithful;
- (b)  $E : \mathbf{C} \rightarrow \mathbf{E}$  surjective on objects and  $M : \mathbf{E} \rightarrow \mathbf{D}$  is full and faithful.

When do the two factorizations agree?

4. Using the fact that representable functors preserve limits, show that for any sets  $A, B, C$ :

$$A^{B+C} \cong A^B \times A^C.$$

Conclude that for any sets  $A, B$  with power sets  $P(A), P(B)$ , etc.:

$$P(A + B) \cong P(A) \times P(B).$$

5. \* Consider the (covariant) composite functor,

$$\mathcal{F} = \mathcal{P}^{\mathbf{BA}} \circ \text{Ult}^{\text{op}} : \mathbf{BA} \rightarrow \mathbf{Sets}^{\text{op}} \rightarrow \mathbf{BA}$$

taking each Boolean algebra  $B$  to the powerset algebra of sets of ultrafilters in  $B$ . Note that,

$$\mathcal{F}(B) \cong \text{Hom}_{\mathbf{Sets}}(\text{Hom}_{\mathbf{BA}}(B, \mathbf{2}), \mathbf{2})$$

is a sort of “double-dual” Boolean algebra. There is always a homomorphism,

$$\phi_B : B \rightarrow \mathcal{F}(B)$$

given by  $\phi_B(b) = \{\mathcal{V} \in \text{Ult}(B) \mid b \in \mathcal{V}\}$ . Show that for any Boolean homomorphism  $h : A \rightarrow B$ , the following square commutes.

$$\begin{array}{ccc} A & \xrightarrow{\phi_A} & \mathcal{F}(A) \\ \downarrow h & & \downarrow \mathcal{F}(h) \\ B & \xrightarrow{\phi_B} & \mathcal{F}(B) \end{array}$$