

HOMEWORK #5
Due Wednesday, September 26

1. Finish reading section 1.3, and start reading section 1.4 in van Dalen.
2. A *binary truth function* is a function $f(x, y)$ that takes values of x and y in the set $\{0, 1\}$ to a value in the set $\{0, 1\}$. Note that a binary truth function is defined uniquely by its truth table.
 - a. How many different binary truth functions are there?
 - b. Two binary truth functions don't depend on any of their arguments: the constant 0 function and the constant 1 function. How many binary truth functions depend only on one of their two arguments?
 - c. We've already seen a number of binary truth functions that depend on both arguments, namely those corresponding to the connectives \wedge , \vee , \rightarrow , \leftrightarrow , \oplus (exclusive or), \mid (nand, or the sheffer stroke), and \downarrow (nor). (The last three are defined by $p \oplus q \equiv \neg(p \leftrightarrow q)$, $p \mid q \equiv \neg(p \wedge q)$, $p \downarrow q \equiv \neg(p \vee q)$.)

What are the remaining ones? You can define them in words, in terms of the other connectives, or with truth tables.

- ★ 3. Show that if k is a natural number and $\varphi_1, \dots, \varphi_k$ are propositional formulas, then $\llbracket \varphi_1 \wedge \dots \wedge \varphi_k \rrbracket_v = 1$ if and only if $\llbracket \varphi_i \rrbracket_v = 1$ for each i from 1 to k . Remember that, for example, $\varphi_1 \wedge \varphi_2 \wedge \varphi_3$ is an abbreviation for $((\varphi_1 \wedge \varphi_2) \wedge \varphi_3)$. Do this carefully, using only the definition of $\llbracket \cdot \rrbracket_v$.
- ★ 4. Show that if $\varphi_1, \dots, \varphi_k$ and ψ are in PROP, then the following is true:

$$\{\varphi_1, \dots, \varphi_k\} \models \psi \text{ if and only if } \models \varphi_1 \wedge \dots \wedge \varphi_k \rightarrow \psi.$$

Once again, do this carefully, using the definition of semantic entailment.

5. Show that if $\{\varphi\} \models \psi$ and $\{\psi\} \models \theta$ then $\{\varphi\} \models \theta$.
- ★ 6. Do problem 1a on page 20 of van Dalen.
7. Do problems 2, 3, 5, and 6 on page 21 of van Dalen.
8. Do problem 1 on page 27.

★ 9. Use “algebraic means” (as in the notes and on page 23 of the textbook) to do problem 2 on page 28 of van Dalen.

10. Use “algebraic means” to show that the following are all tautologies:

a. $((\varphi \wedge \neg\psi) \vee \psi) \leftrightarrow (\varphi \vee \psi)$

b. $(\varphi \rightarrow \neg\varphi) \rightarrow \neg\varphi$

c. $(\varphi \rightarrow \psi) \leftrightarrow (\neg\psi \rightarrow \neg\varphi)$

d. $\varphi \rightarrow (\psi \rightarrow \varphi \wedge \psi)$

★ 11. Let \equiv be any equivalence relation on a set X . For any element a in X , let $[a]$ denote the equivalence class of a , defined by

$$[a] = \{b \in X \mid b \equiv a\}.$$

Show that for any elements a and b of X , $a \equiv b$ if and only if $[a] = [b]$. (Remember that two sets said to be equal if and only if they have exactly the same elements.)

★ 12. Now let \equiv denote equivalence modulo 5 on the natural numbers. In other words, $a \equiv b$ holds iff $a - b$ is a multiple of 5, that is, iff there is an integer c such that $a - b = 5c$.

a. Define the operation of addition \oplus on equivalence classes by

$$[a] \oplus [b] = [a + b].$$

Show that this operation is well defined, that is, if $a \equiv a'$ and $b \equiv b'$ then $a + b \equiv a' + b'$.

b. Define exponentiation \uparrow on equivalence classes by

$$[a] \uparrow [b] = [a^b].$$

Show that exponentiation is *not* well-defined.

○ 13. In the problem above, show that multiplication on equivalence classes, defined by $[a] \otimes [b] = [a \times b]$, is well-defined.