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# **Security: An Overview of Cryptographic Techniques**

15-640/440

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# Cryptography, Cryptographic Protocols and Key Distribution

- Authentication
- Mutual Authentication
- Private/Symmetric Keys
- Public Keys
- Key Distribution

# What do we need for a secure communication channel?

- Authentication (Who am I talking to?)
  - Confidentiality (Is my data hidden?)
  - Integrity (Has my data been modified?)
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- Availability (Can I reach the destination?)

# What is cryptography?

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"cryptography is about communication in the presence of adversaries."

- Ron Rivest

“cryptography is using math and other crazy tricks to approximate magic”

- Unknown TA

# What is cryptography?

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Tools to help us build secure communication channels that provide:

- 1) Authentication
- 2) Integrity
- 3) Confidentiality

# Cryptography As a Tool

- Using cryptography securely is not simple
- Designing cryptographic schemes correctly is near impossible.

Today we want to give you an idea of what can be done with cryptography.

Take a security course if you think you may use it in the future

# The Great Divide

**Symmetric** Crypto  
(Private key)  
(E.g., AES)

**Asymmetric** Crypto  
(Public key)  
(E.g., RSA)

**Shared secret**  
between parties?

Yes

No

**Speed** of crypto  
operations

Fast

Slow

# Symmetric Key: Confidentiality

## Motivating Example:

You and a friend share a key  $K$  of  $L$  random bits, and want to secretly share message  $M$  also  $L$  bits long.

## Scheme:

You send her the  $xor(M, K)$  and then she “decrypts” using  $xor(M, K)$  again.

- 1) Do you get the right message to your friend?
- 2) Can an adversary recover the message  $M$ ?
- 3) Can adversary recover the key  $K$ ?



# Symmetric Key: Confidentiality

- One-time Pad (OTP) is secure but usually impractical
  - ◆ Key is as long as the message
  - ◆ Keys cannot be reused (why?)

In practice, two types of ciphers are used that require constant length keys:

## **Stream Ciphers:**

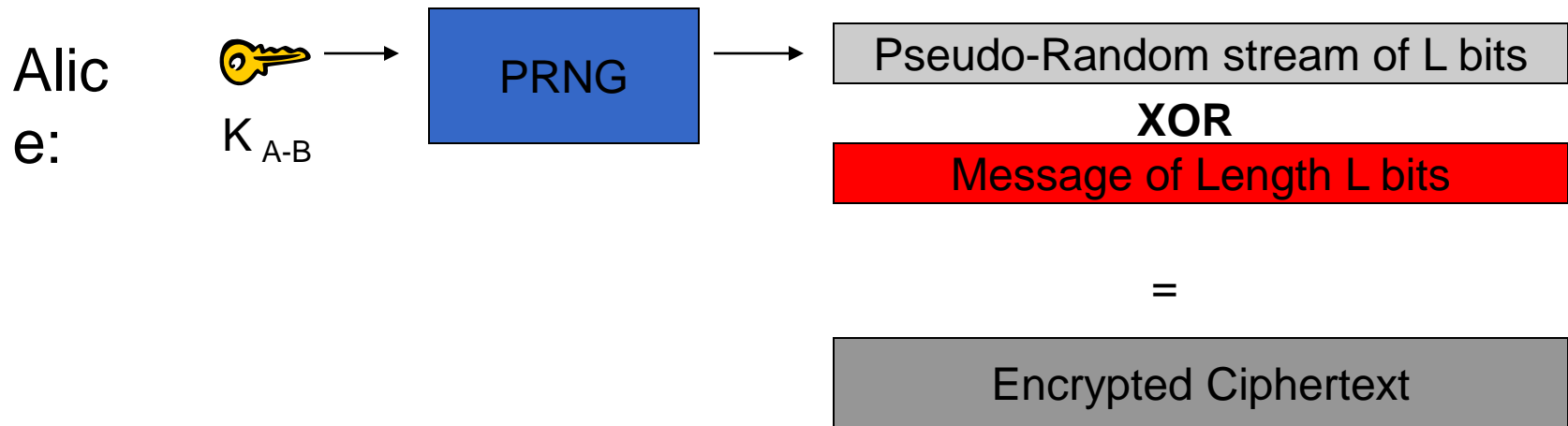
Ex: RC4, A5

## **Block Ciphers:**

Ex: DES, AES,  
Blowfish

# Symmetric Key: Confidentiality

- Stream Ciphers (ex: RC4)



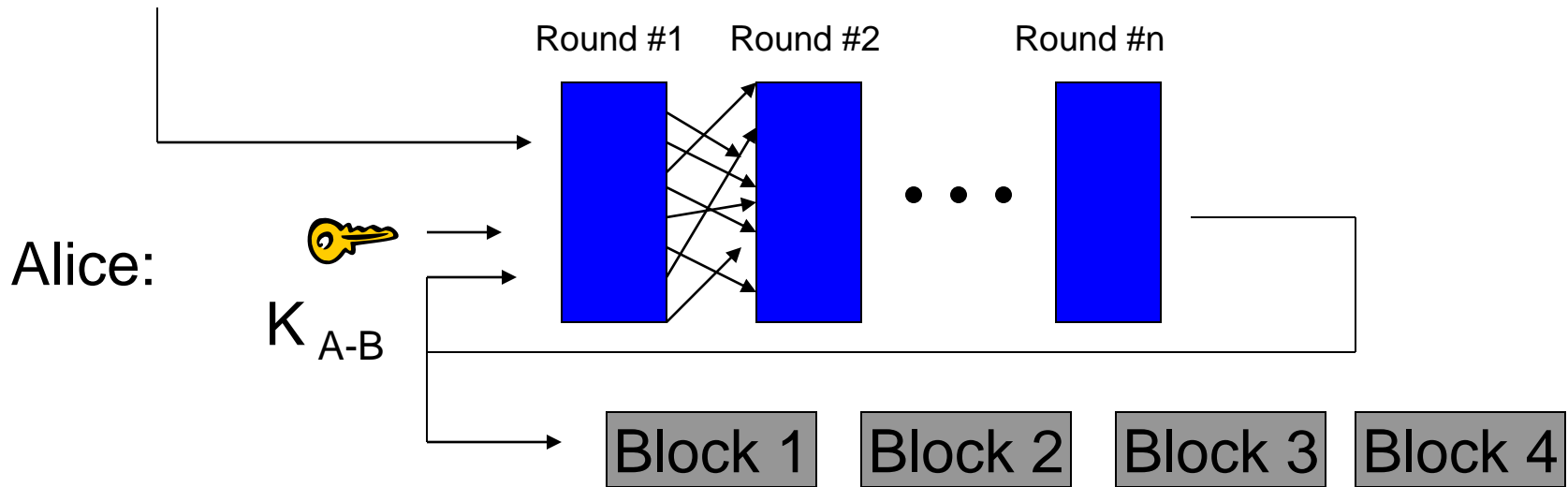
Bob uses  $K_{A-B}$  as PRNG seed, and XORs encrypted text to get the message back (just like OTP).

# Symmetric Key: Confidentiality

- Block Ciphers (ex: AES)

Block 1   Block 2   Block 3   Block 4

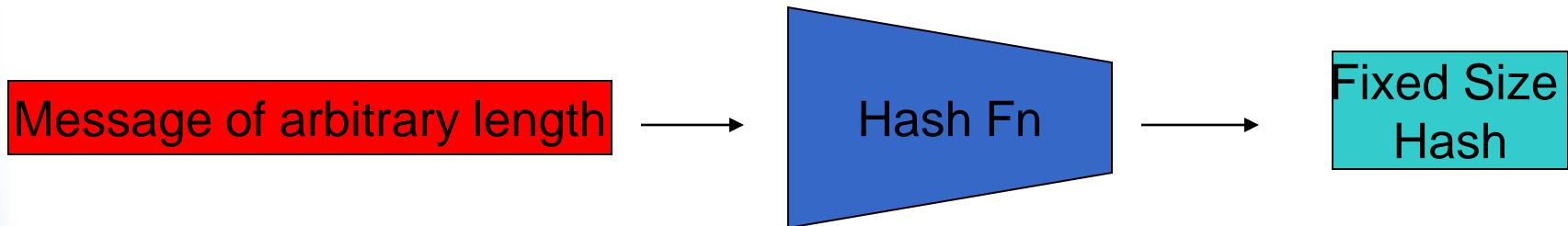
(fixed block size,  
e.g. 128 bits)



Bob breaks the ciphertext into blocks, feeds it through decryption engine using  $K_{A-B}$  to recover the message.

# Cryptographic Hash Functions

- **Consistent**  
hash(X) always yields same result
- **One-way**  
given Y, can't find X s.t. hash(X) = Y
- **Collision resistant**  
given hash(W) = Z, can't find X such that hash(X) = Z

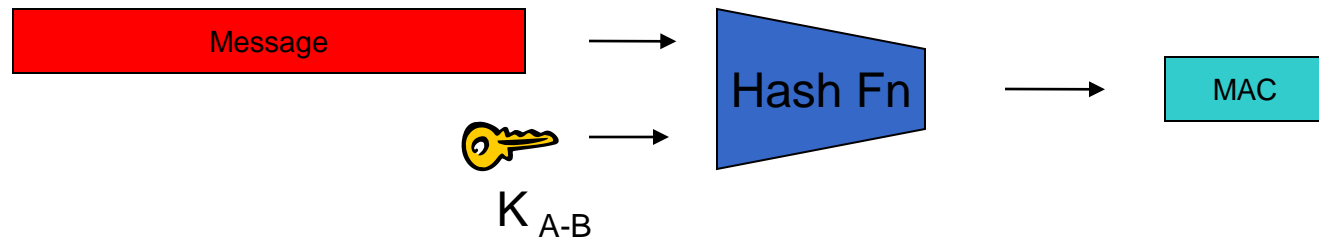


# Symmetric Key: Integrity

- Hash Message Authentication Code (HMAC)

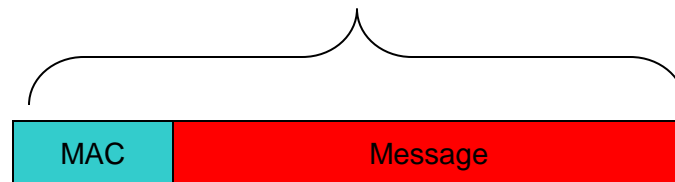
Step #1:

Alice creates  
MAC



Step #2

Alice Transmits Message & MAC



Step #3

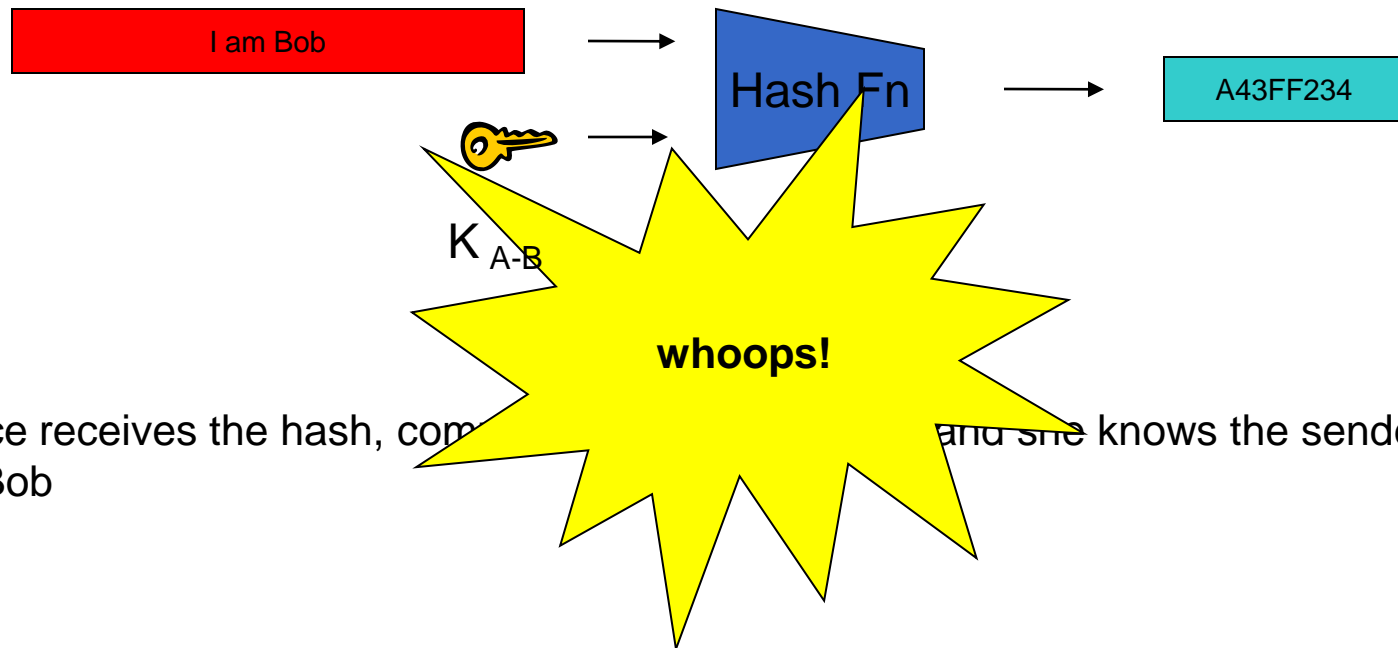
Bob computes MAC with  
message and  $K_{A-B}$  to verify.

Why is this secure?

How do properties of a hash function help us?

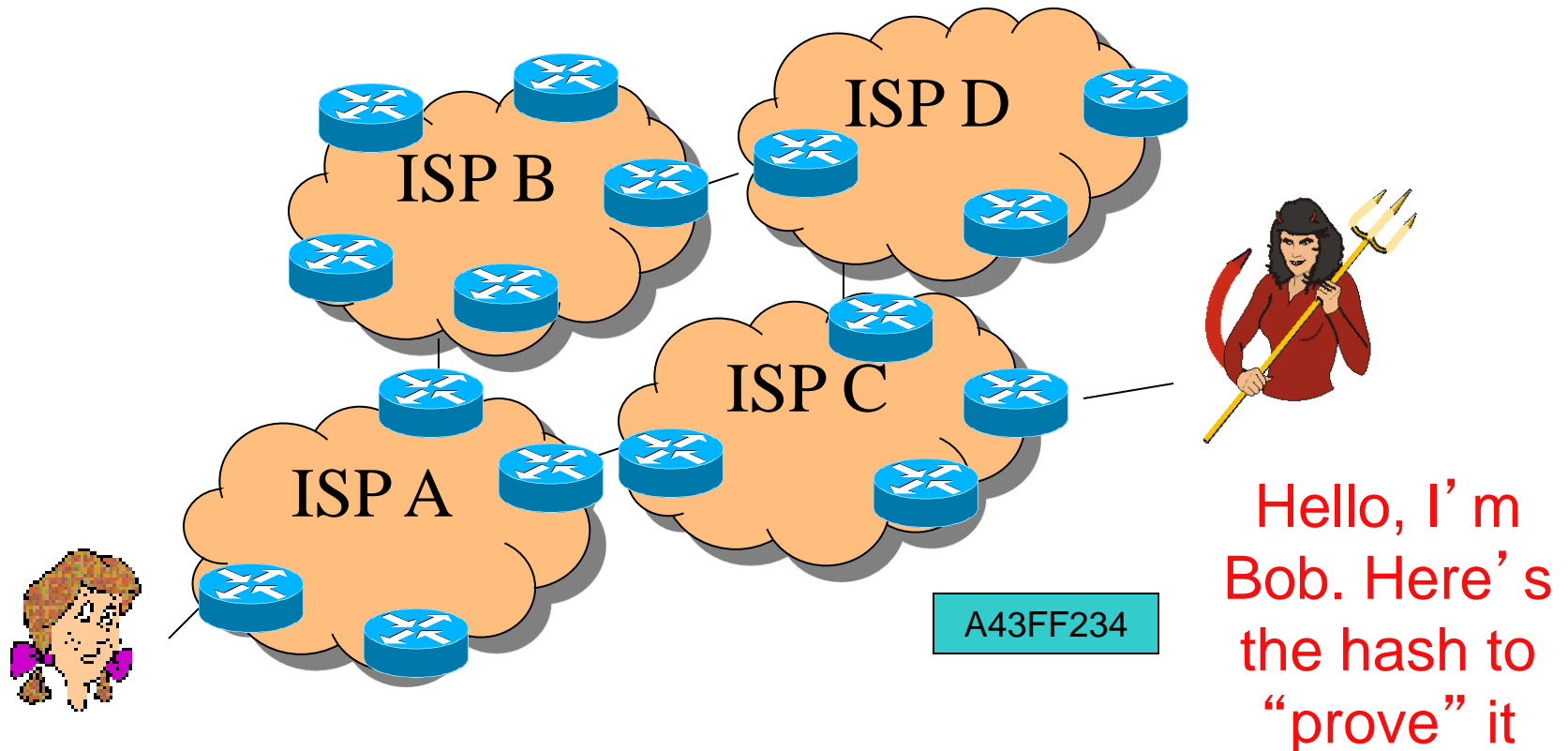
# Symmetric Key: Authentication

- You already know how to do this!  
(hint: think about how we showed integrity)



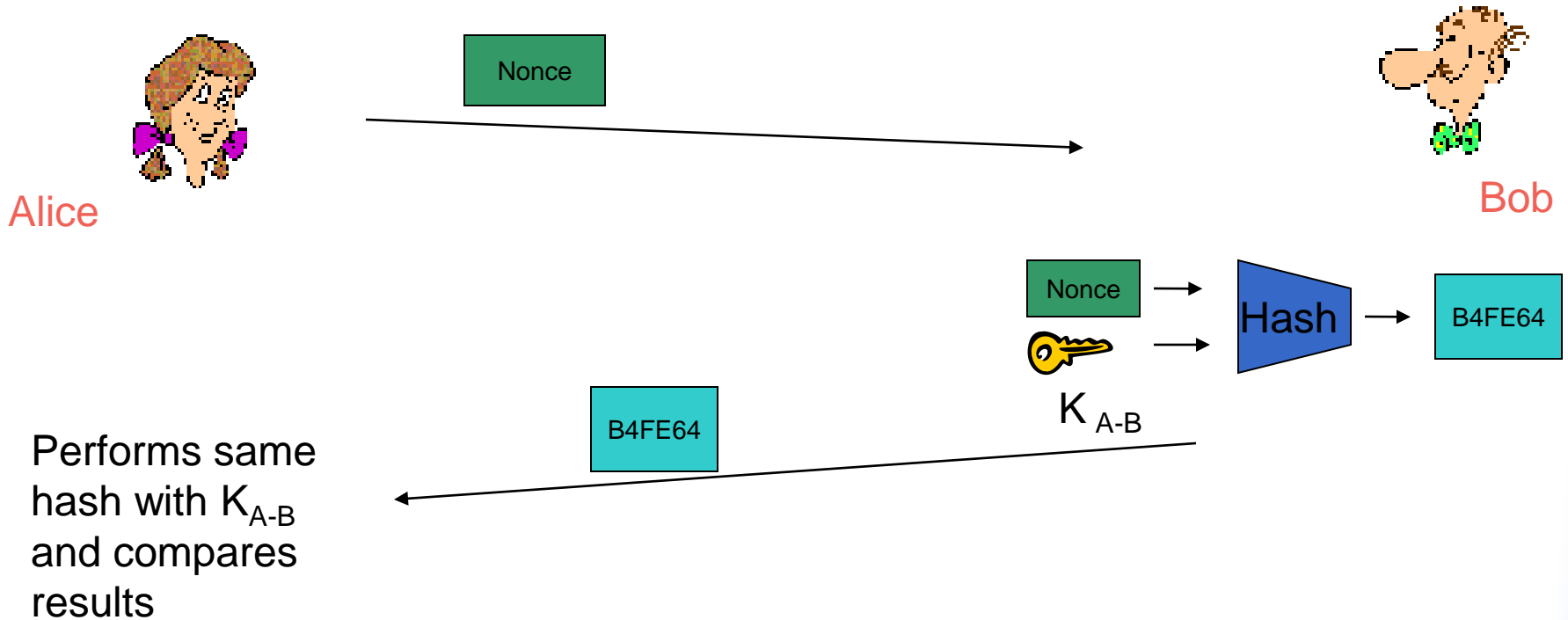
# Symmetric Key: Authentication

What if Mallory overhears the hash sent by Bob, and then “replays” it later?



# Symmetric Key: Authentication

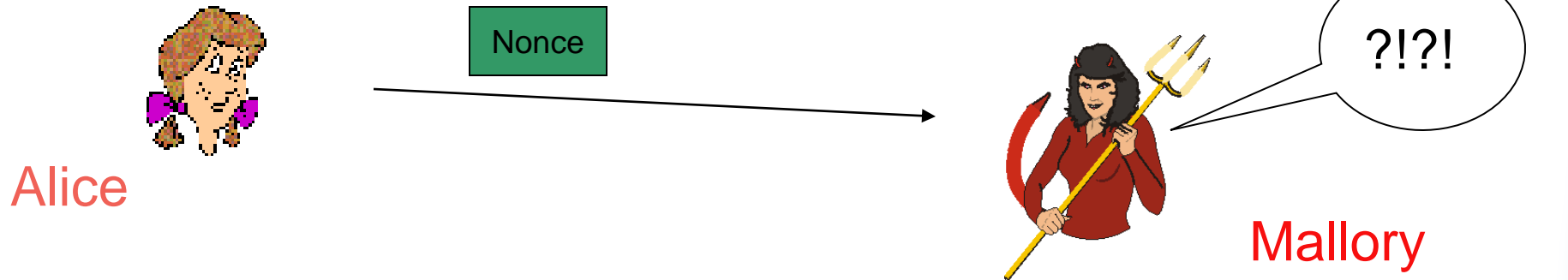
- A “Nonce”
  - ♦ A random bitstring used only once. Alice sends nonce to Bob as a “challenge”. Bob Replies with “fresh” MAC result.





# Symmetric Key: Authentication

- A “Nonce”
  - ♦ A random bitstring used only once. Alice sends nonce to Bob as a “challenge”. Bob Replies with “fresh” MAC result.



If Alice sends Mallory a nonce, she cannot compute the corresponding MAC without  $K_{A-B}$

# Symmetric Key Crypto Review

- Confidentiality: Stream & Block Ciphers
- Integrity: HMAC
- Authentication: HMAC and Nonce

**Questions??**

**Are we done? Not Really:**

**1) Number of keys scales as  $O(n^2)$**

**2) How to securely share keys in the first place?**

# Asymmetric Key Crypto:

- Instead of shared keys, each person has a “key pair”



$K_B$  Bob's public key



$K_B^{-1}$  Bob's private key

- The keys are inverses, so:  $K_B^{-1} (K_B (m)) = m$

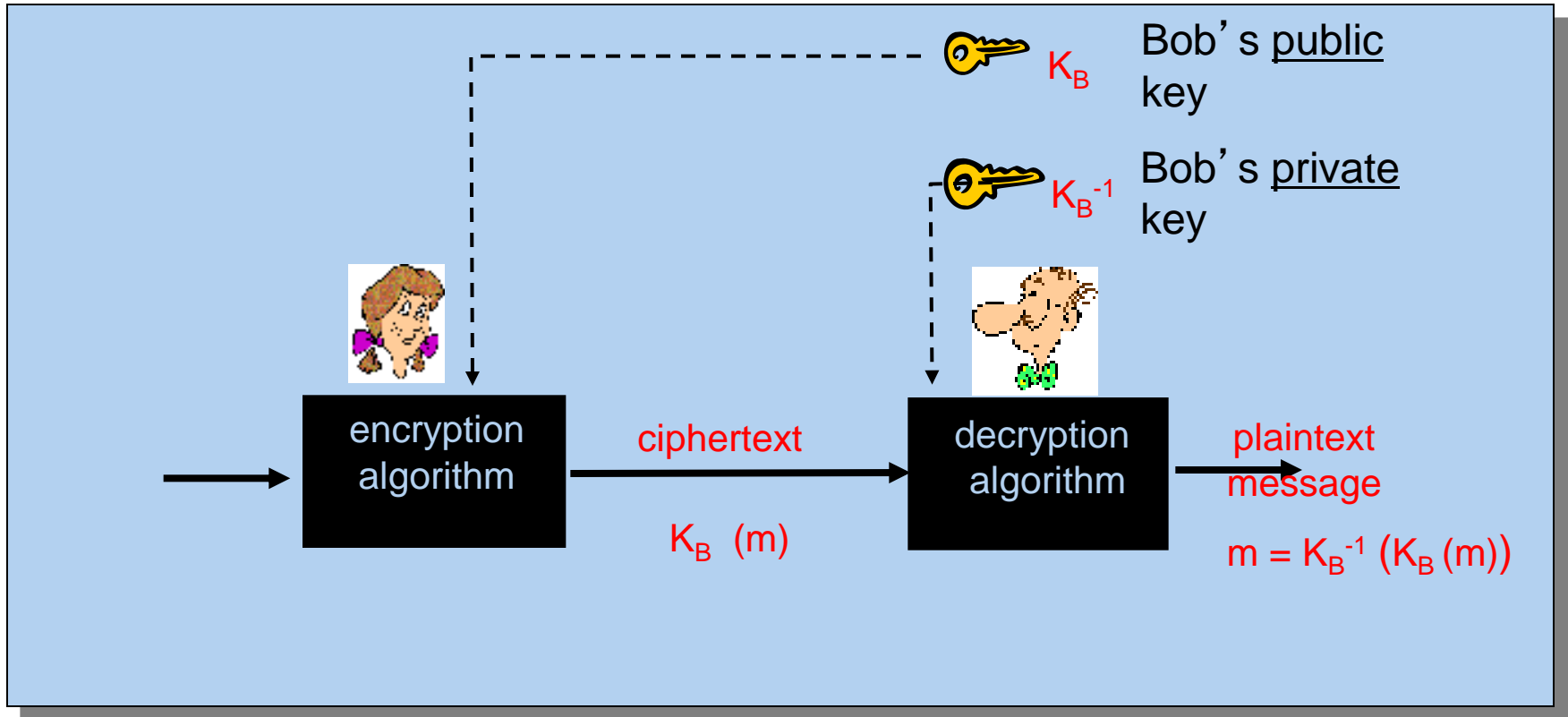
# Asymmetric Key Crypto:

- It is believed to be computationally unfeasible to derive  $K_B^{-1}$  from  $K_B$  or to find any way to get  $M$  from  $K_B(M)$  other than using  $K_B^{-1}$  .

=>  $K_B$  can safely be made public.

Note: We will not explain the computation that  $K_B(m)$  entails, but rather treat these functions as black boxes with the desired properties.

# Asymmetric Key: Confidentiality



# Asymmetric Key: Sign & Verify

- If we are given a message  $M$ , and a value  $S$  such that  $K_B(S) = M$ , what can we conclude?
- The message must be from Bob, because it must be the case that  $S = K_B^{-1}(M)$ , and only Bob has  $K_B^{-1}$  !
- This gives us two primitives:
  - Sign  $(M) = K_B^{-1}(M) = \text{Signature } S$
  - Verify  $(S, M) = \text{test}( K_B(S) == M )$

# Asymmetric Key Review:

- Confidentiality: Encrypt with Public Key of Receiver
- Integrity: Sign message with private key of the sender
- Authentication: Entity being authenticated signs a nonce with private key, signature is then verified with the public key

But, these operations are computationally expensive\*

# Biometrics

- Nice in some respects
  - No need to distribute
  - Reducible to digital form
  - Unique in practice
- Hard to duplicate?
  - Used via binary representation
  - Warm gelatin fingers or slip-on finger-pads molded to prints?
  - Artificial eyeballs made to match scans?
  - Pictures? Videos w/blinking?
- Change over time?
  - Injury?
  - Aging?
- **Not replaceable or revocable**
  - What happens when “stolen?”
  - Are you “Deleted”?!?!?
  - (Well, you do have 10 fingers, two retinas, one nose, etc)



# Multi-Factor, Human Factors

- Best systems use more than one factor
  - Something you know
  - Something piece of you
  - Biometrics + Password/Q&A Challenge, Etc
  - More natural factors better than fewer unnatural challenges
  - More weak factors may be stronger than fewer stronger factors
- Human factors are critical
  - Too many password restrictions? Too many passwords?
    - Write them down on Post-Its Notes!

# Summary – Part II

- Symmetric (pre-shared key, fast) and asymmetric (key pairs, slow) primitives provide:
  - Confidentiality
  - Integrity
  - Authentication
- “Hybrid Encryption” leverages strengths of both.
- Great complexity exists in securely acquiring keys.
- Crypto is hard to get right, so use tools from others, don't design your own (e.g. TLS).