

# Distributed Systems

## CS 15-440

Lecture 25, April 30, 2013

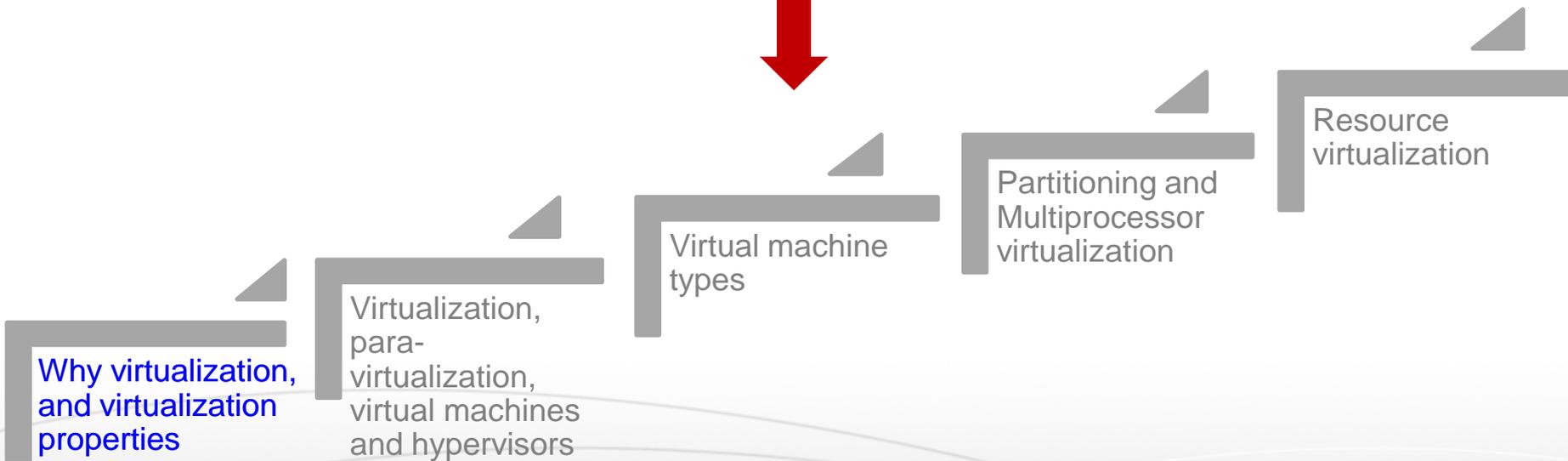
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Borrowed from our good friends in Doha:

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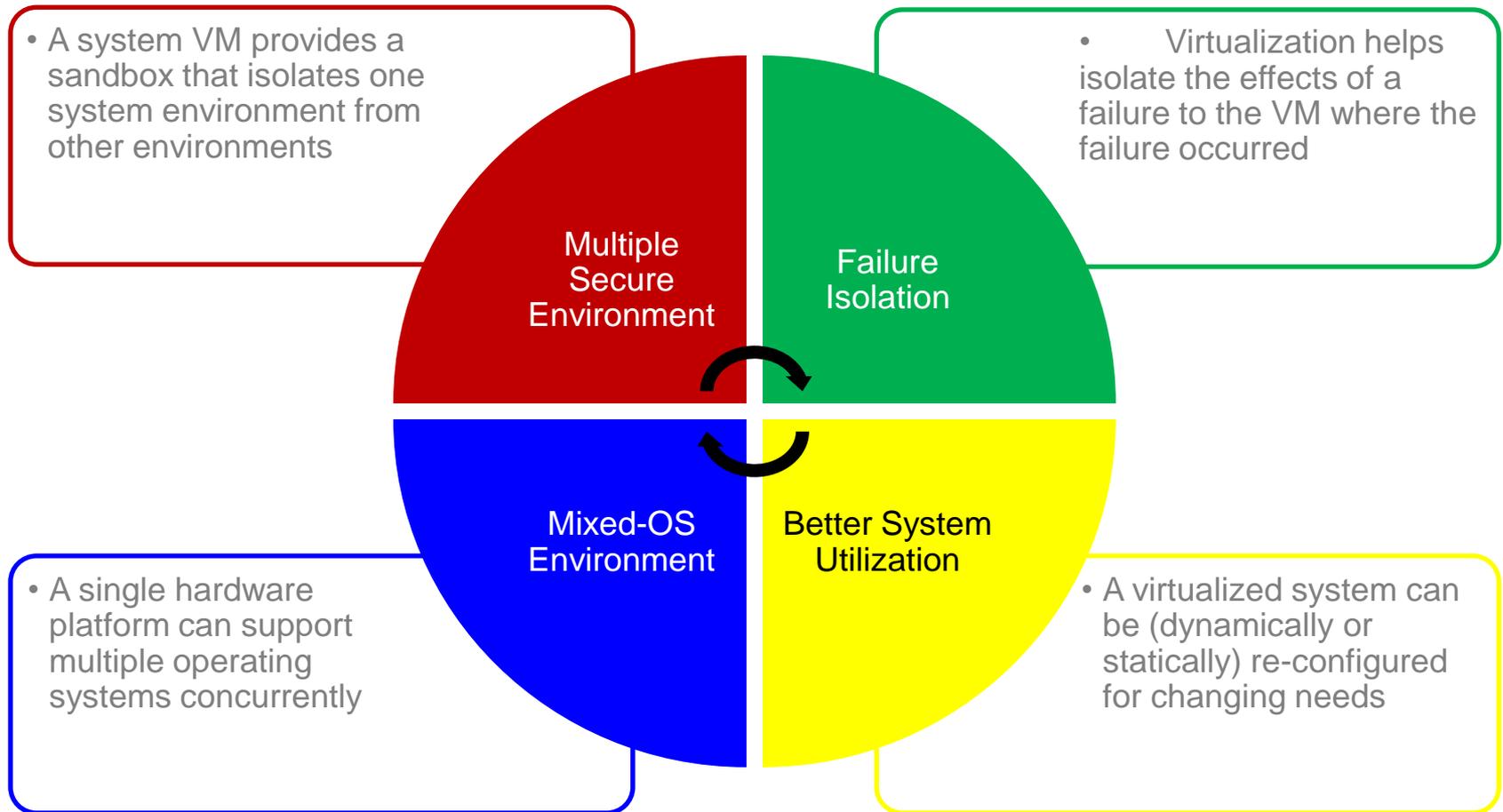
# Objectives

Discussion on Virtualization



# Benefits of Virtualization

- Here are some of the benefits that are typically provided by a virtualized system



# Operating Systems Limitations

- OSs provide a way of virtualizing hardware resources among *processes*
- This may help isolate *processes* from one another
- However, this does not provide a virtual machine to a user who may wish to run a different OS
- Having hardware resources managed by a single OS limits the flexibility of the system in terms of available software, security, and failure isolation
- Virtualization typically provides a way of relaxing constraints and increasing flexibility

# Virtualization Properties

- Fault Isolation
- Software Isolation
- Performance Isolation (accomplished through scheduling and resource allocation)

Isolation

1

- All VM state can be captured into a file (i.e., you can operate on VM by operating on file— cp, rm)
- Complexity is proportional to virtual HW model and independent of guest software configuration

Encapsulation

2

- All guest actions go through the virtualizing software which can inspect, modify, and deny operations

Interposition

3

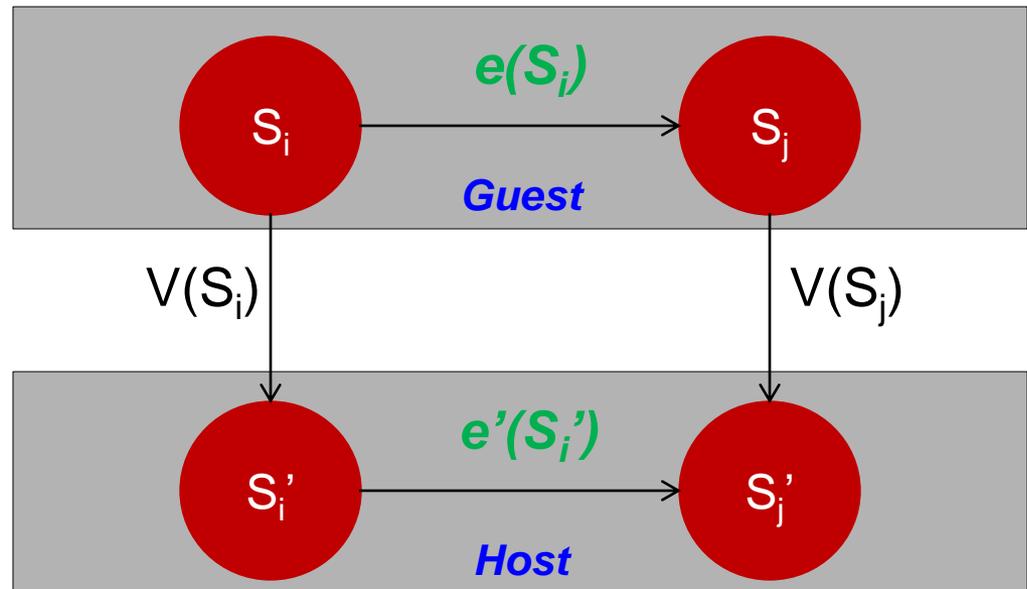
# What is Virtualization?

- Informally, a virtualized system (or subsystem) is a mapping of its interface, and all resources visible through that interface, to the interface and resources of a real system
- Formally, virtualization involves the construction of an isomorphism that maps a virtual *guest* system to a real *host* system (Popek and Goldberg 1974)

✓ Function  $V$  maps the guest state to the host state

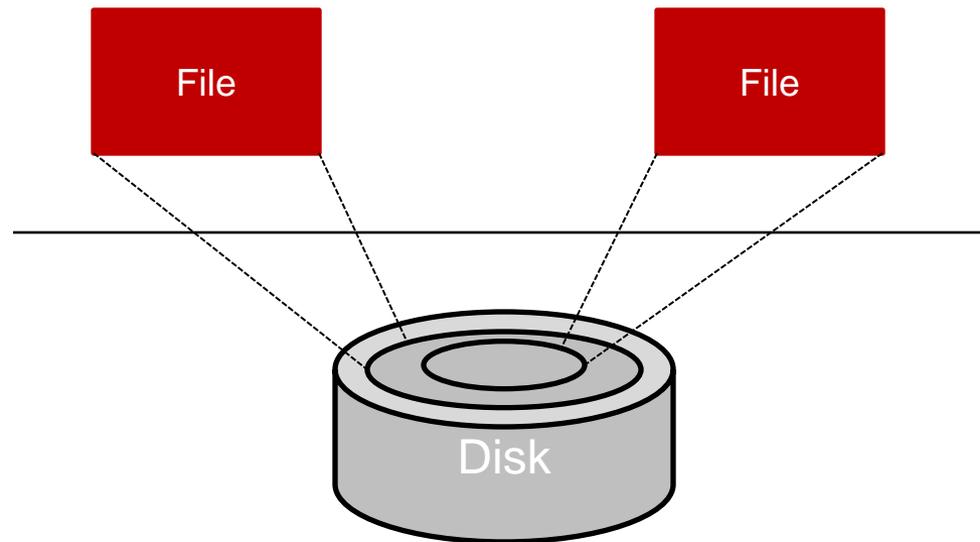
✓ For a sequence of operations,  $e$ , that modifies a guest state, there is a corresponding  $e'$  in the host that performs an equivalent modification

✓ How can this be managed?



# Abstraction

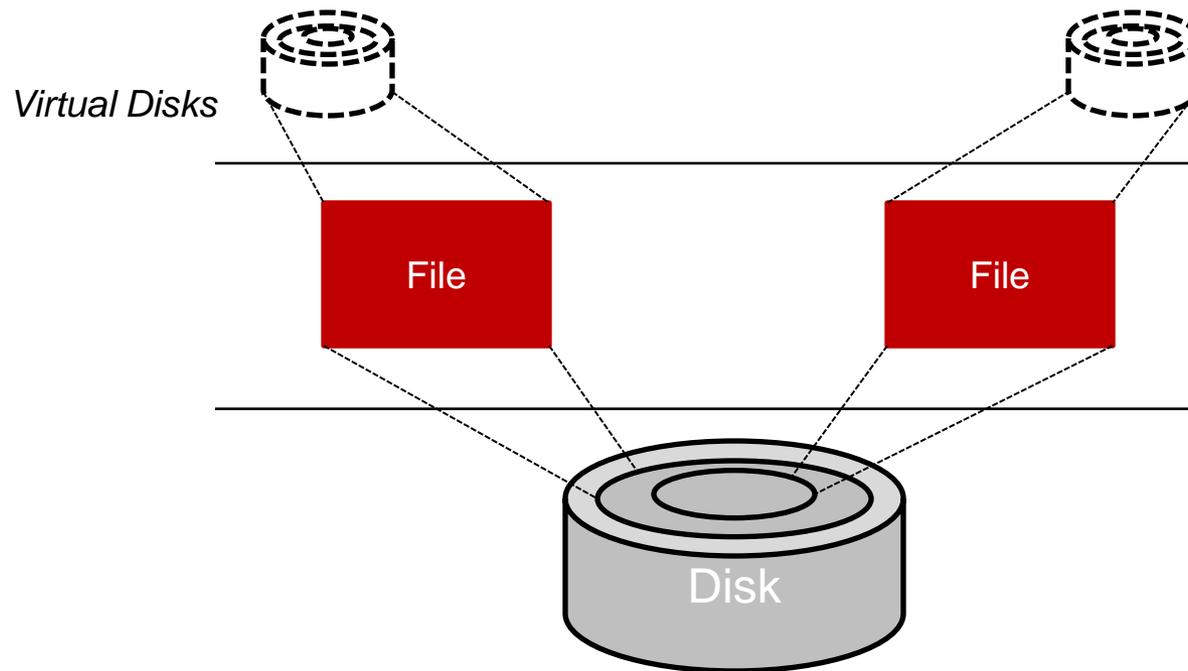
- The key to managing complexity in computer systems is their division into *levels of abstraction* separated by *well-defined interfaces*
- Levels of abstraction allow implementation details at lower levels of a design to be ignored or simplified



- ✓ Files are an abstraction of a Disk
- ✓ A level of abstraction provides a simplified interface to underlying resources

# Virtualization and Abstraction

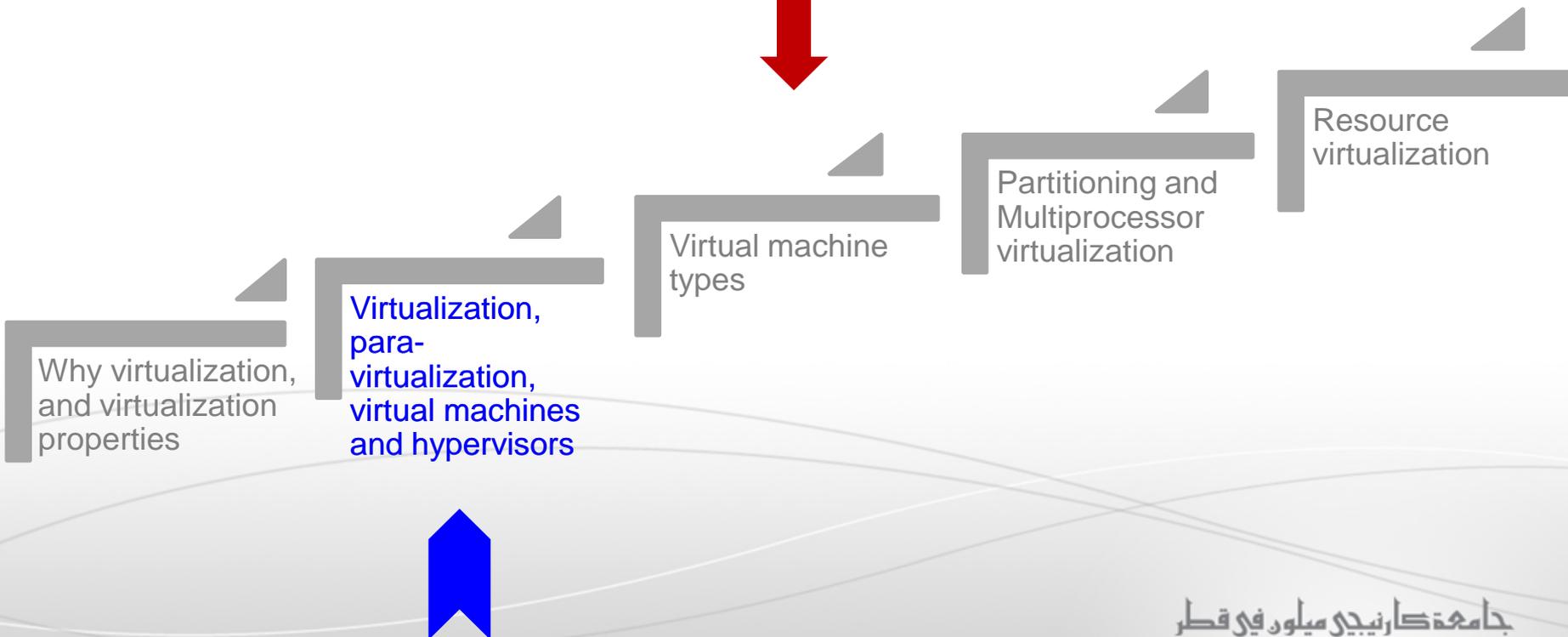
- Virtualization uses abstraction but is different in that it doesn't necessarily hide details; the level of detail in a virtual system is often the same as that in the underlying real system



✓ *Virtualization provides a different interface and/or resources at the same level of abstraction*

# Objectives

Discussion on Virtualization

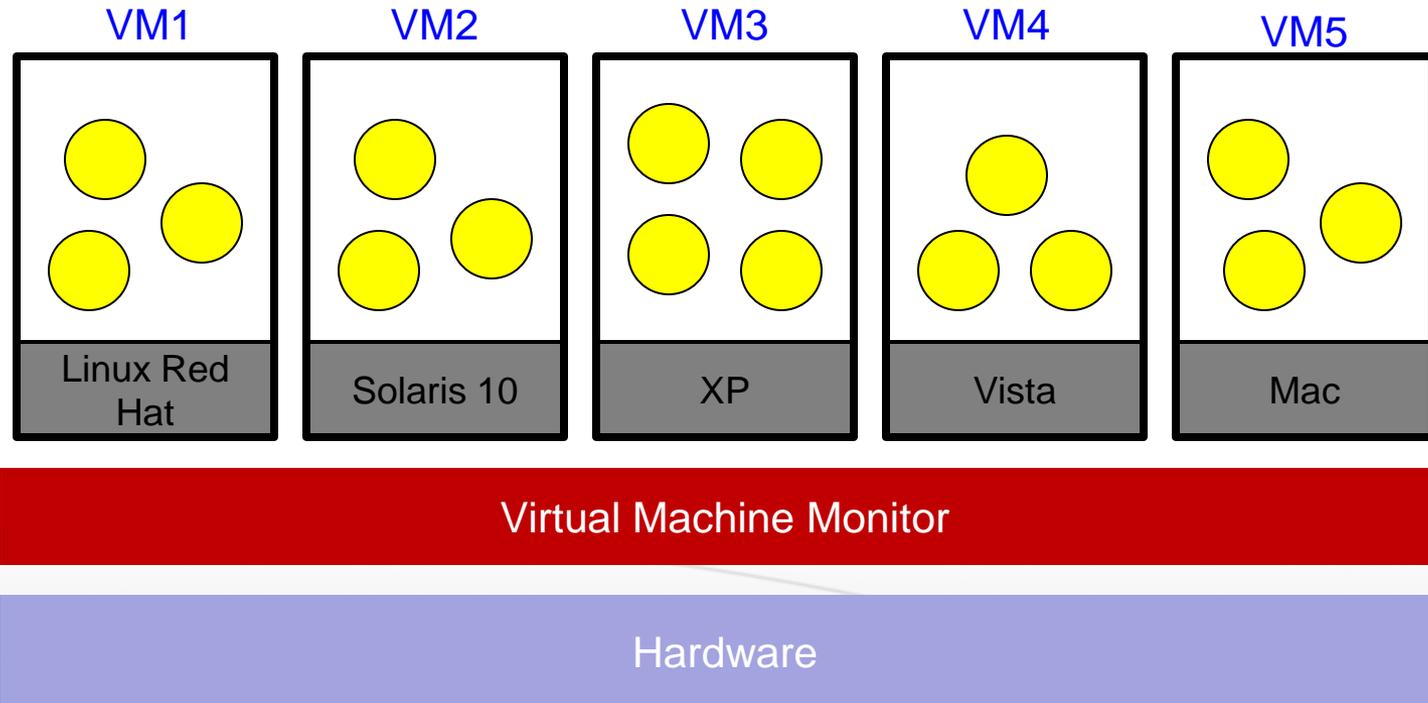


# Virtual Machines and Hypervisors

- The concept of virtualization can be applied not only to subsystems such as disks, but to an entire machine denoted as a **virtual machine (VM)**
- A VM is implemented by adding a layer of software to a real machine so as to support the desired VM's architecture
- This layer of software is often referred to as **virtual machine monitor (VMM)**
- Early VMMs are implemented in firmware
- Today, VMMs are often implemented as a co-designed firmware-software layer, referred to as the **hypervisor**

# A Mixed OS Environment

- Multiple VMs can be implemented on a single hardware platform to provide individuals or user groups with their own OS environments



# Full Virtualization

- Traditional VMMs provide **full-virtualization**:
  - The functionality provided is identical to the underlying physical hardware
  - The functionality is exposed to the VMs
  - They allow unmodified guest OSs to execute on the VMs
    - This might result in some performance degradation
  - E.g., *VMWare* provides full virtualization

# Para-Virtualization

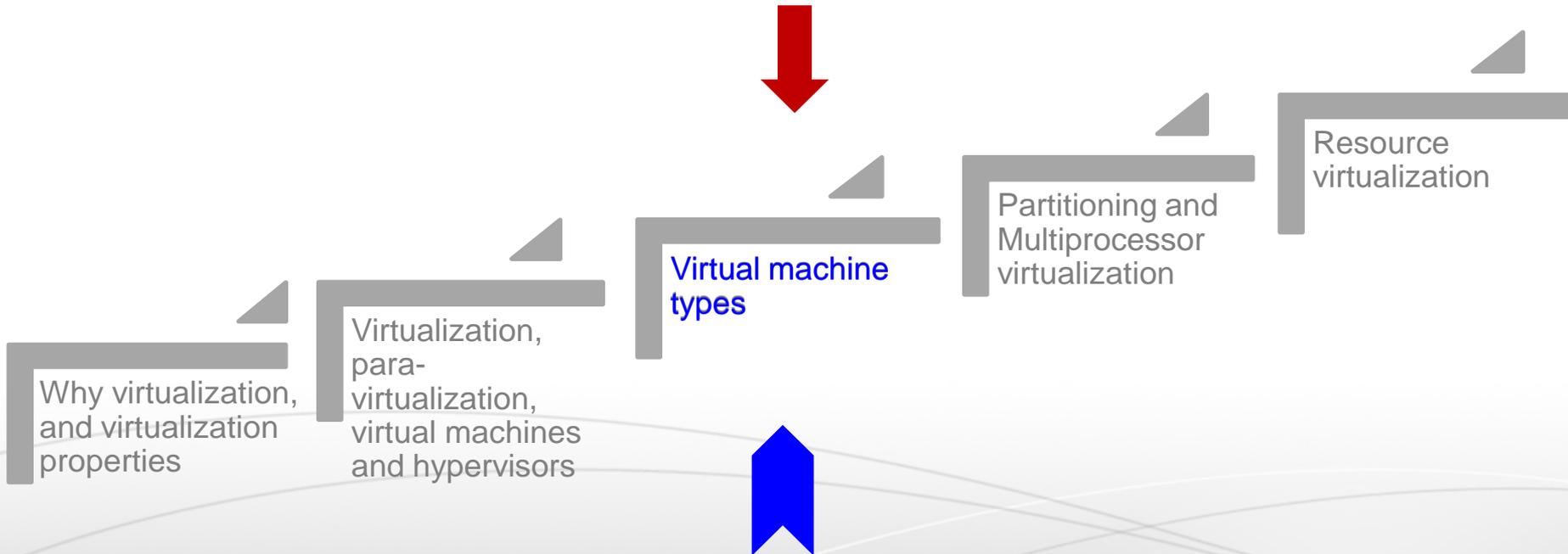
- Other types of VMMs provide **para-virtualization**:
  - They provide a virtual hardware abstraction that is similar, but not identical to the real hardware
  - They modify the guest OS to cooperate with the VMM
  - They result in lower overhead leading to better performance
  - E.g., *Xen* provides both para-virtualization as well as full-virtualization

# Virtualization and Emulation

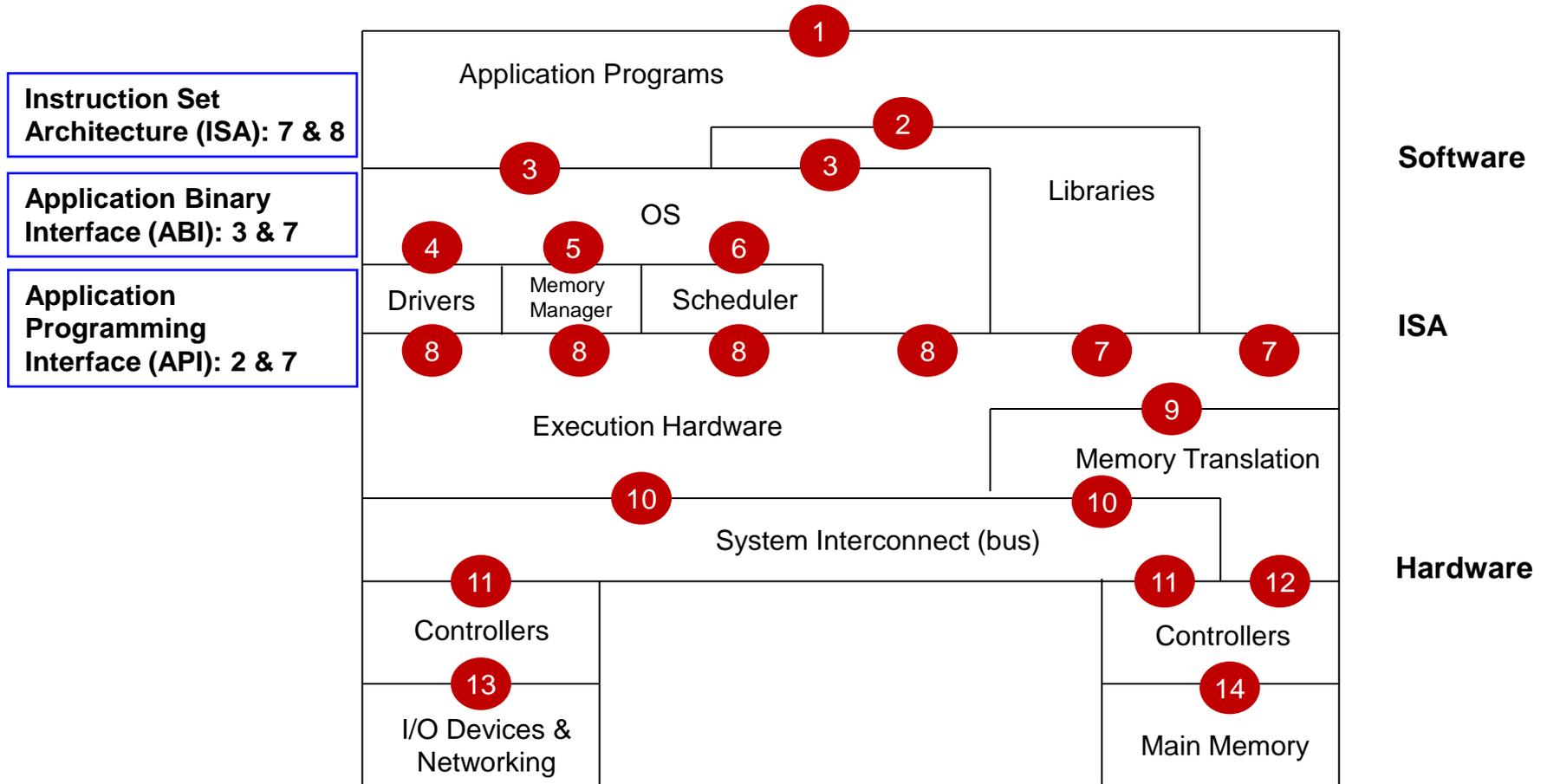
- VMs can employ *emulation techniques* to support cross-platform software compatibility
- Compatibility can be provided either at the system level (e.g., to run a Windows OS on Macintosh) or at the program or process level (e.g., to run Excel on a Sun Solaris/SPARC platform)
- Emulation is the process of implementing the interface and functionality of one system on a system having a different interface and functionality
- It can be argued that virtualization itself is simply a form of emulation

# Objectives

Discussion on Virtualization



# Background: Computer System Architectures



# Types of Virtual Machines

- As there is a process perspective and a system perspective of machines, there are also process-level and system-level VMs
- Virtual machines can be of two types:

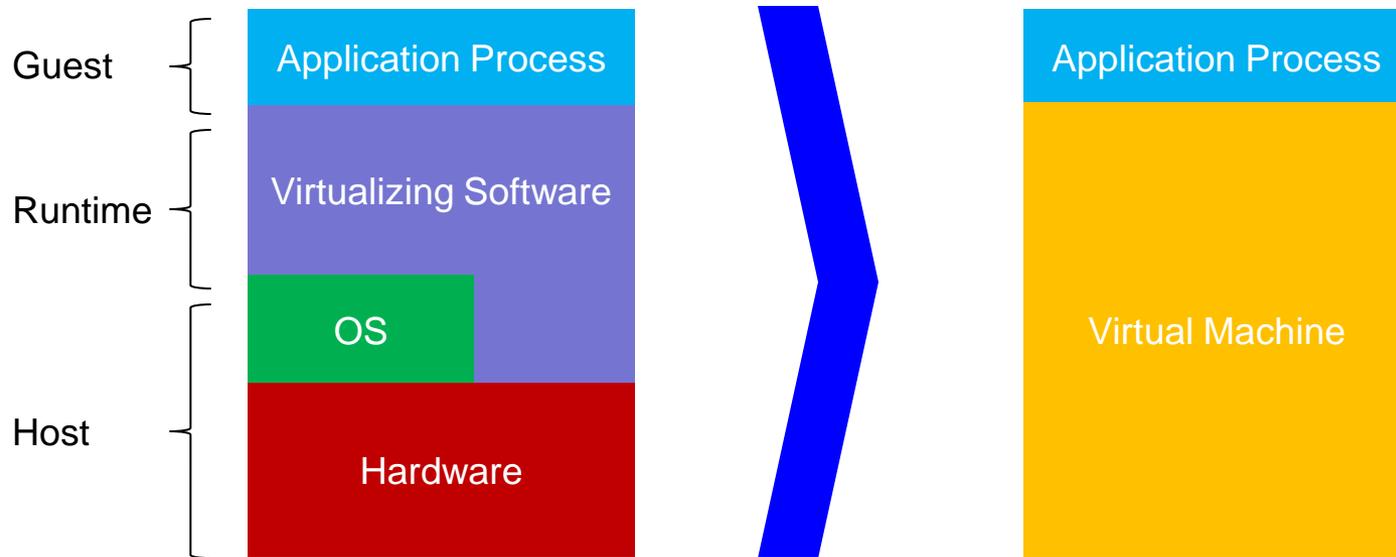
## 1. Process VM

- Capable of supporting an individual process

## 2. System VM

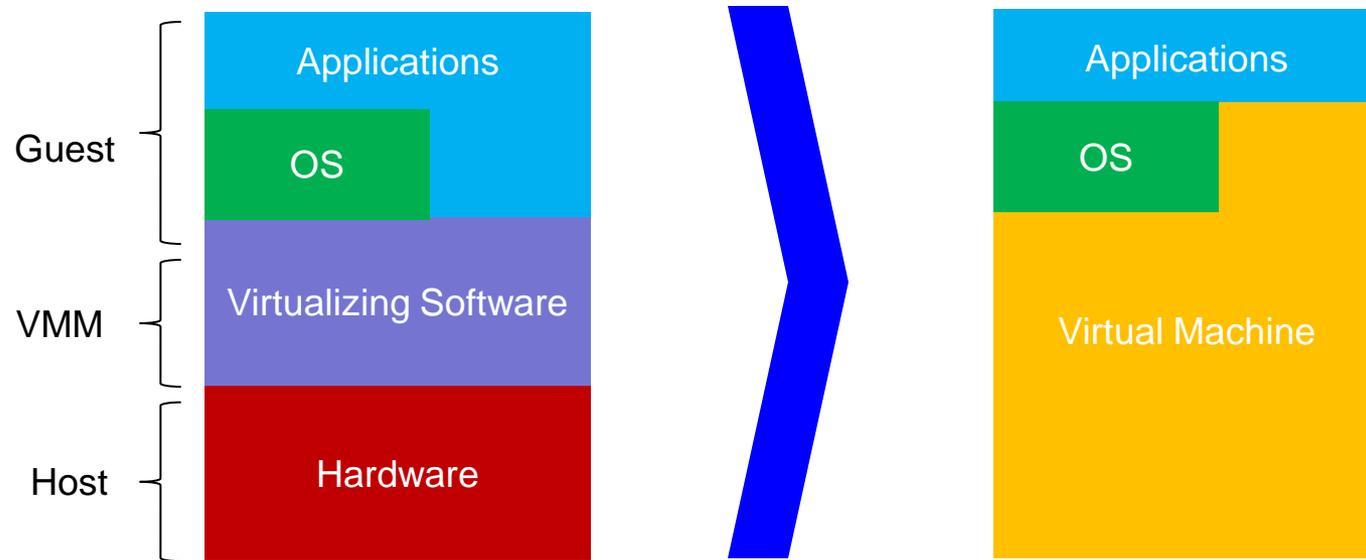
- Provides a complete system environment
- Supports an OS with potentially many types of processes

# Process Virtual Machine



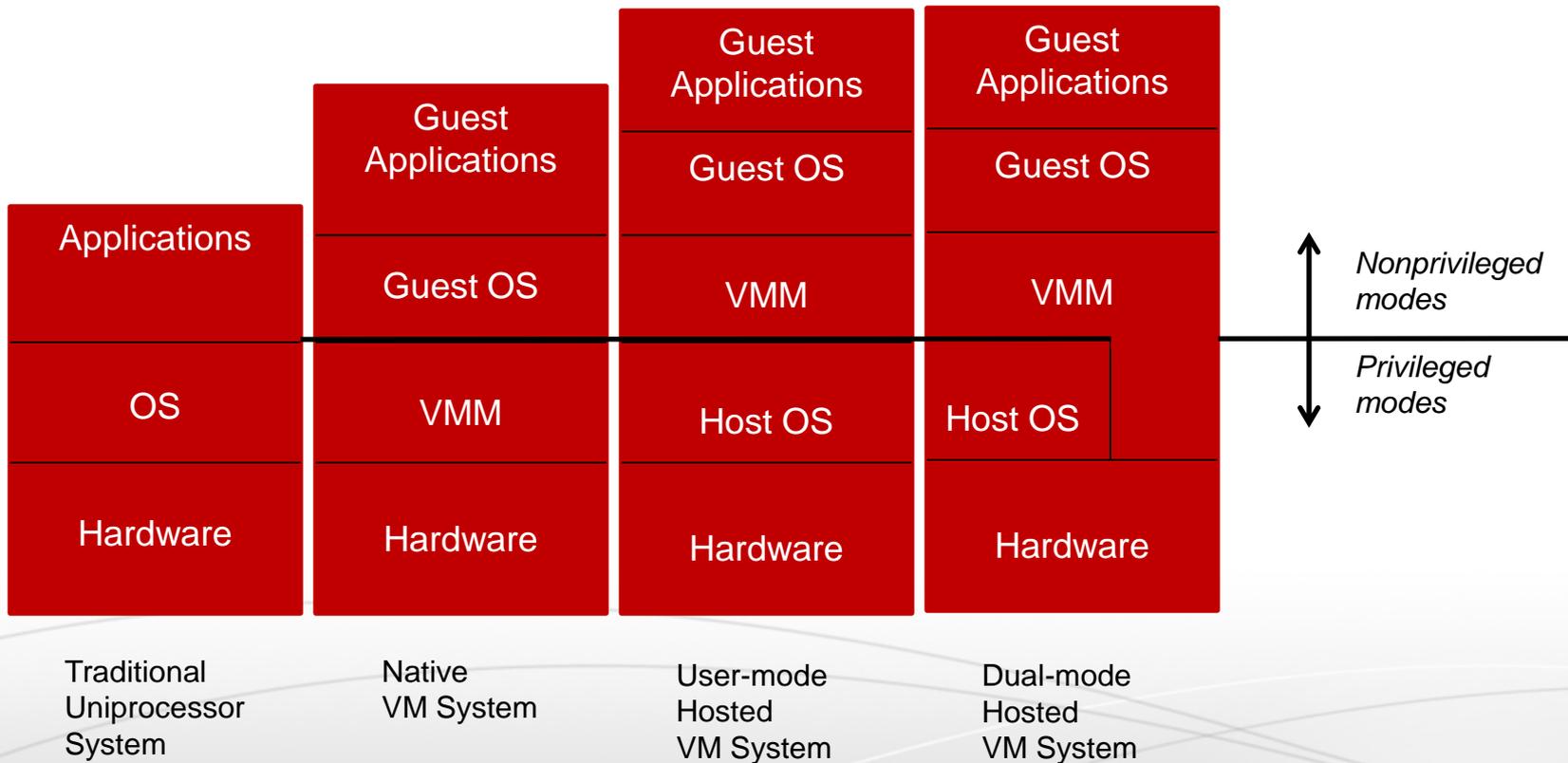
- ✓ Runtime is placed at the ABI interface
- ✓ Runtime emulates both user-level instructions and OS system calls

# System Virtual Machine

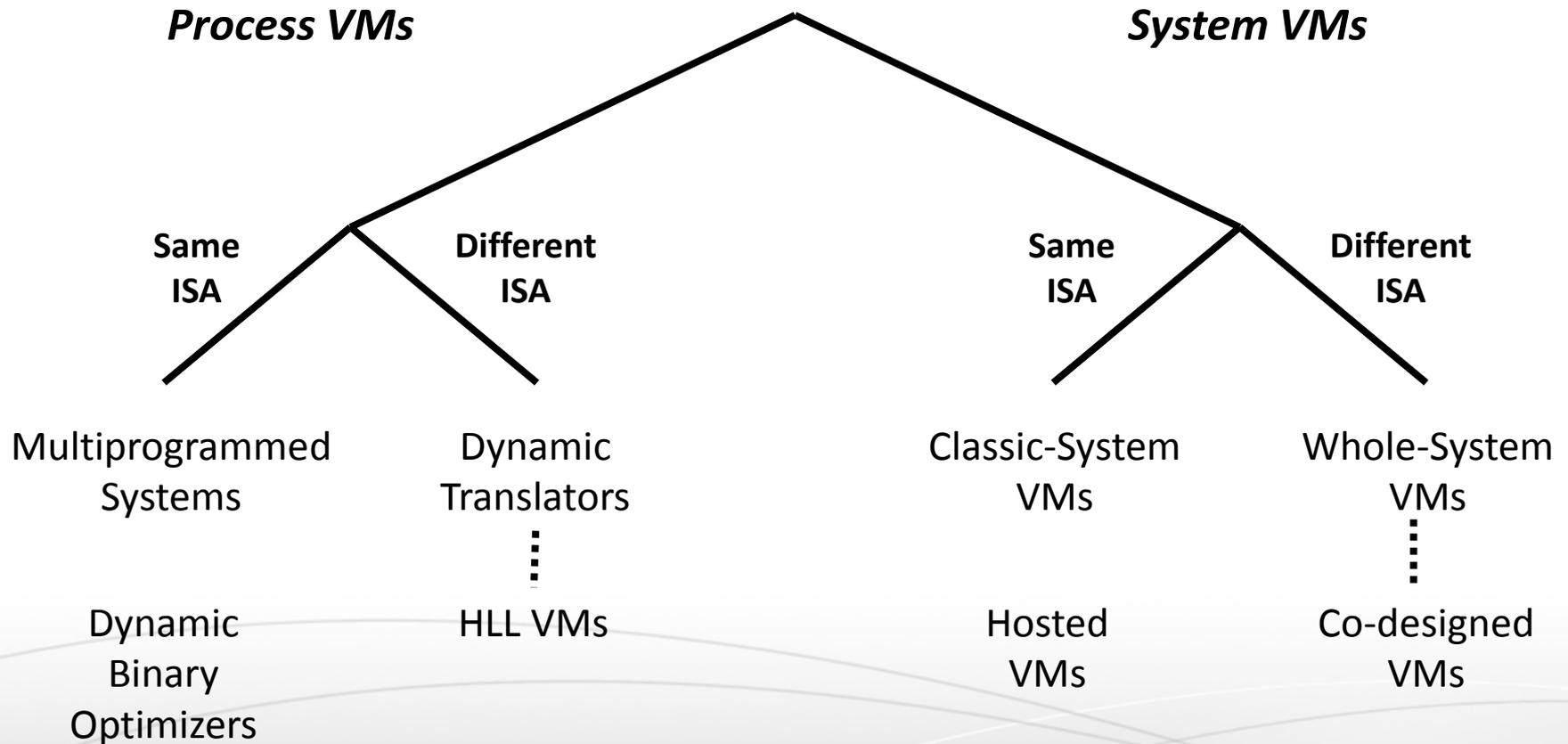


- ✓ VMM emulates the ISA used by one hardware platform to another, forming a system VM
- ✓ A system VM is capable of executing a system software environment developed for a different set of hardware

# Native and Hosted VM Systems



# A Taxonomy



# The Versatility of VMs

Java Application

*JVM*

Linux IA-32

*VMWare*

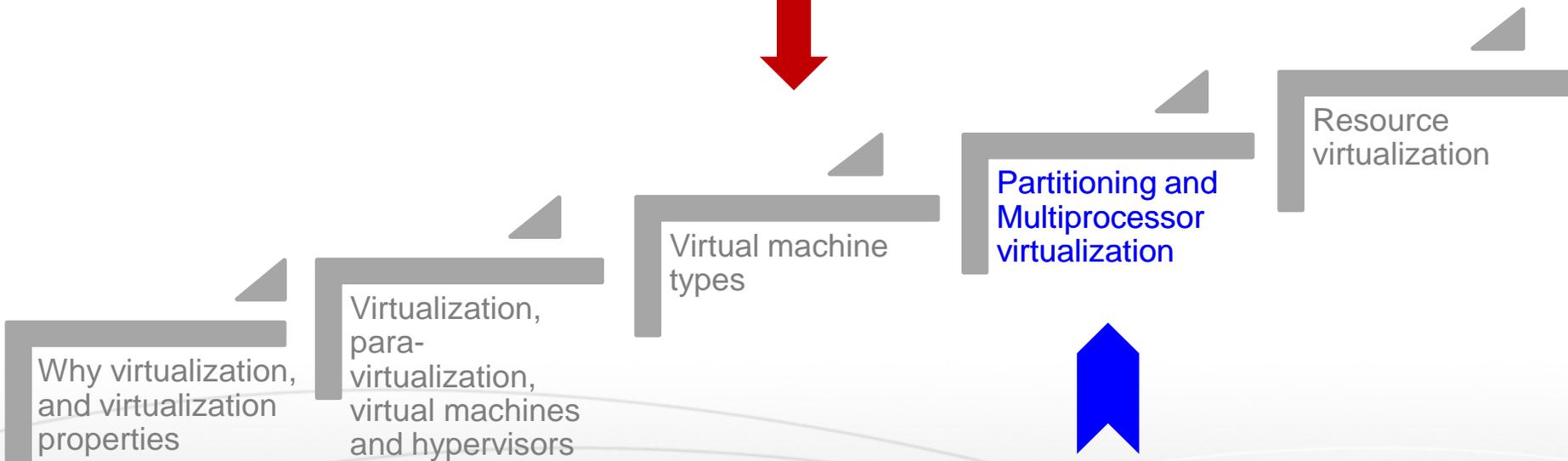
Windows IA-32

*Code Morphing*

Crusoe VLIW

# Objectives

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# Multiprocessor Systems

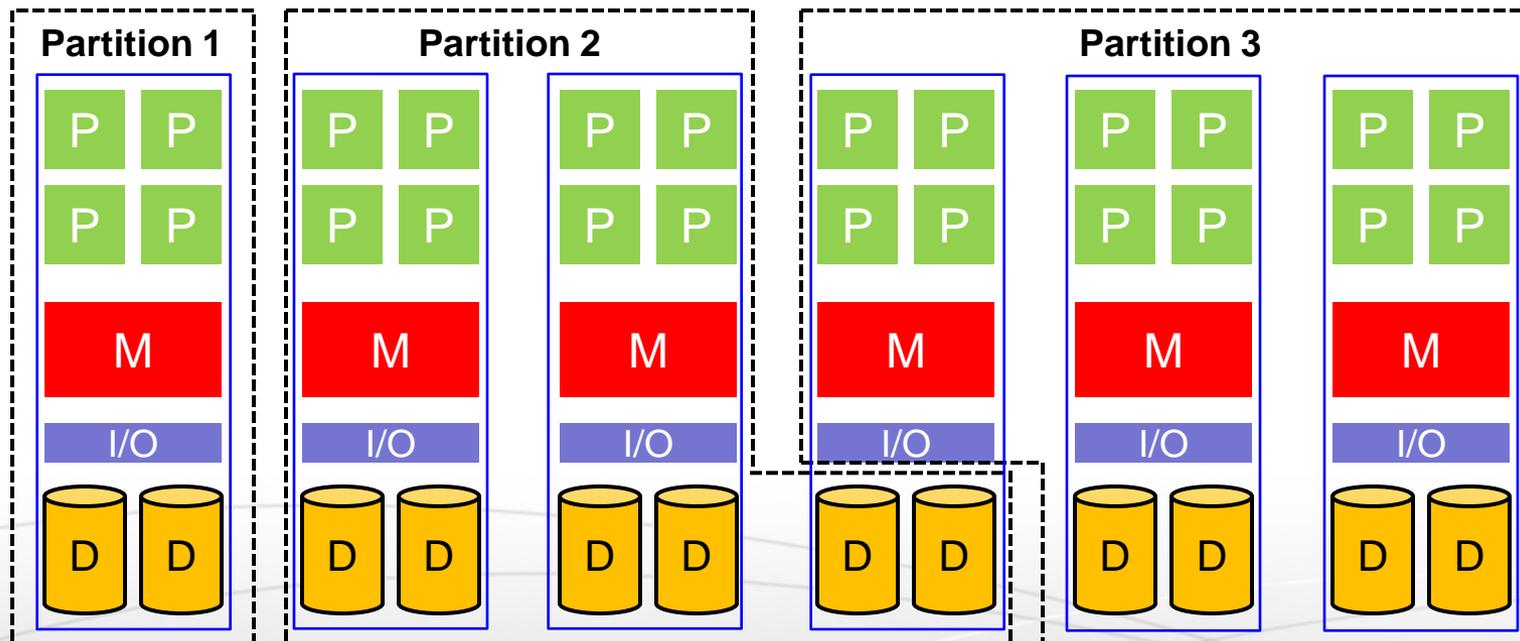
- Multiprocessor systems might have 1000s of processors connected to TBs of memory and PBs of disk capacity
- Often there is a mismatch between the ideal number of processors an application needs and the actual number of physical processors available
- It is more often the case that applications cannot exploit more than a fraction of the processors available. This is mainly because of:
  - Limitations in the parallelism available in the programs
  - Limitations in the scalability of applications due to the overhead of communication between processors

# Partitioning

- The increasing availability of multiprocessor systems has led to the examination of techniques that can help *utilize* them more effectively
- Techniques have been developed in which the multiprocessor system can be **partitioned** into multiple **partitions**
  - A partition is given a subset of the resources available on the system
- Hence, using partitioning, multiple applications can simultaneously exploit the available resources of the system
- Partitioning can be achieved:
  - Either **in-space** (referred to as **physical partitioning**)
  - Or **in-time** (referred to as **logical partitioning**)

# Physical Partitioning

- With physical partitioning, each partition is assigned resources that are physically distinct from the resources used by the other partitions



# Physical Partitioning

- Physical partitioning allows a partition to own its resources physically
- It is not permissible for two partitions to share the resources of a single system board
- Partitions are configured by a *central control unit* that receives commands from the console of the system admin and provisions hardware resources accordingly
- The number of partitions that can be supported in physically partitioned systems is limited to the number of available physical processors

# Physical Partitioning- Advantages

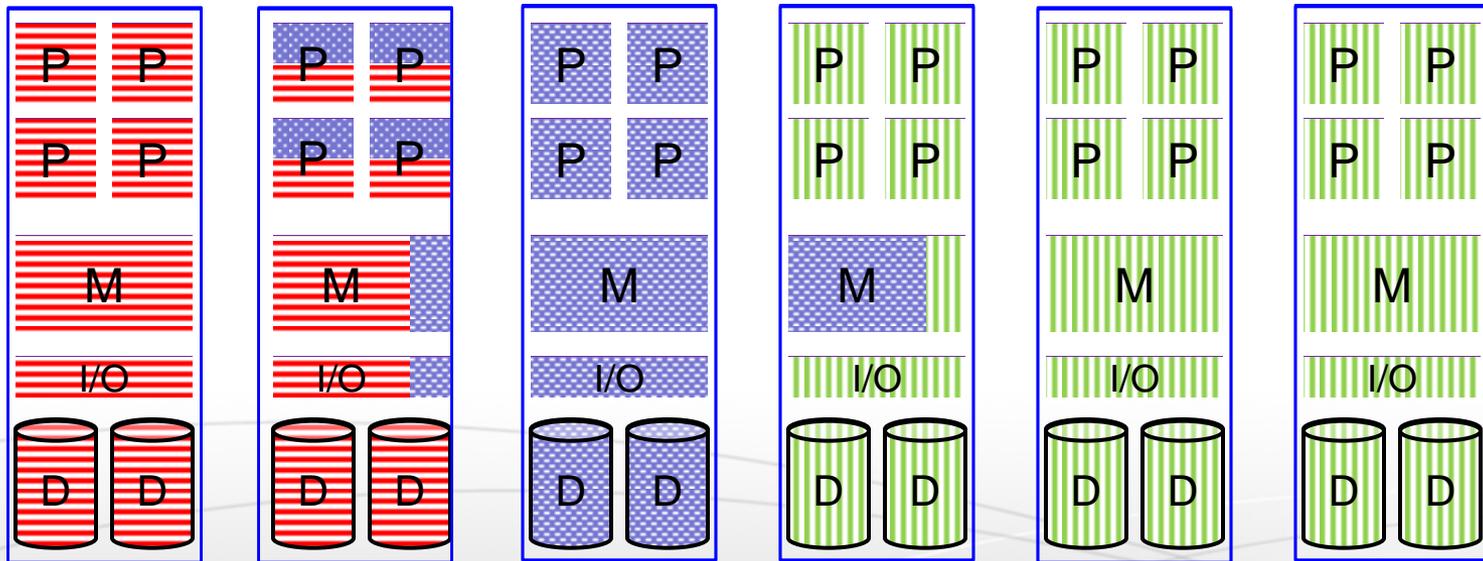
- Physical partitioning provides:
  - **Failure Isolation**: it ensures that in the event of a failure, only the part of the physical system that houses the failing partition will be affected
  - **Better security isolation**: Each partition is protected from the possibility of intentional or unintentional denial-of-service attacks by other partitions
  - **Better ability to meet system-level objectives** (these result from contracts between system owners and users of the system)
  - **Easier management of resources**: no need of sophisticated algorithms for scheduling and management of resources

# Physical Partitioning- Disadvantages

- While physical partitioning has a number of attractive features, it has some major disadvantages:
  - **System utilization:** Physical partitioning is probably not the ideal solution if system utilization is to be optimized
    - It is often the case that each of the physical partitions is underutilized
  - **Load balancing:** with physical partitioning, dynamic workload balancing becomes difficult to implement

# Logical Partitioning

- With logical partitioning, partitions share some of the physical resources, usually in a *time-multiplexed* manner



# Logical Partitioning

- With logical partitioning it is permissible for two partitions to share the resources of a single system board
- Logical partitioning makes it possible to partition an ***n-way*** system into a system with more than ***n*** partitions, if so desired
- Logical partitioning is more flexible than physical partitioning but needs additional mechanisms to provide safe and efficient way of sharing resources
- Logical partitioning is usually done through a VMM or a hypervisor and provides what is referred to as ***multiprocessor virtualization***

# Multiprocessor Virtualization

- A virtualized multiprocessor gives the appearance of a system that may or may not reflect the exact configuration of the underlying physical system

