

# Symbolic Integration

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## Integrals of a Logarithmic Extension

In this section we outline the decision procedure algorithm for the log extension defined by

$$D(t) = D(u) / u$$

$$\int \frac{\partial}{\partial x} u(x) \frac{dx}{u(x)} = \log u(x)$$

Applying the Euclidean division algorithm and then Hermite's algorithm we obtain

$$\int \frac{a(t)}{b(t)} dx = \text{Rational\_Function}(t) + \int c(t) + \int \frac{p(t)}{q(t)}$$

where  $q$  is squarefree. The first integral is called the polynomial part and the second - the rational part. Unlike to rational function integration, the integration of the polynomial part is the hardest of two.

### Polynomial Part

We are to integrate

$$\int c(t)$$

where  $c$  is a polynomial in  $t(x)$ . In other words,  $c$  is a polynomial in logs

$$c(t) = a_k t^k + \dots + a_0$$

Integrating this, yields

$$\int c(t) = \int a_k t^k + \dots + a_0 = b_{k+1} t^{k+1} + \dots + b_0 + \sum_{k=1}^n d_k * \log(u_k)$$

Coefficients  $b_p(x)$  can be computed by differentiating both sides wrt  $x$

$$c(\theta) = \frac{d}{dx} \left( b_{k+1}(x) t(x)^{k+1} + \dots + b_0(x) + \sum_{k=1}^n d_k * \log(u_k(x)) \right)$$

or

$$a_k t^k + \dots + a_0 = b_{k+1}' t^{k+1} + \left[ (k+1) b_{k+1} t' + b_k' \right] t^k + \dots + \left[ 2 b_2 t' + b_1' \right] t + b_1 t' + b_0' + \frac{d}{dx} \sum_{k=1}^n d_k * \log u_k$$

Equating coefficients by  $t$ , we get a system of equations

$$\begin{aligned} b_{k+1}' &= 0 \\ (k+1) b_{k+1} t' + b_k' &= a_k \\ \dots &\dots \\ 2 b_2 t' + b_1' &= a_1 \\ b_1 t' + b_0' + \frac{d}{dx} \sum_{k=1}^n d_k * \log u_k &= a_0 \end{aligned}$$

We solve the system from top to bottom, recursively applying the integration algorithm. Though each step requires integration, the integrand lies in a lower field extension.

The algorithm is terminated when integration cannot be done or the whole system has no solution.

Note, integration itself might require a new LOG extension. In this case we terminate algorithm as well. Though, we could allow a new LOG extension at the last step.

### ■ Example.

Consider

$$\int \log^2(x) + 2 \log(x) dx$$

Extending  $\mathbb{Q}[x]$  by  $t = \log(x)$ , we obtain

$$\int t^2 + 2 t dx = b_3 t^3 + b_2 t^2 + b_1 t + b_0 + \sum_{k=1}^n d_k * \log(u_k(x))$$

where  $b_k$  are free of  $t = \log(x)$ . Differentiating both sides wrt  $x$

$$t^2 + 2 t = b_3' t^3 + 3 b_3 t' t^2 + b_2' t^2 + 2 b_2 t' t + b_1' t + b_1 t' + b_0' + \frac{d}{dx} \sum_{k=1}^n d_k * \log(u_k(x))$$

and equating coefficients by  $t$ , we get a system of equations

$$\begin{aligned}
 b_3' &= 0 \\
 3 b_3 t' + b_2' &= 1 \\
 2 b_2 t' + b_1' &= 2 \\
 b_1 t' + b_0' + \frac{d}{dx} \sum_{k=1}^n d_k * \log(u_k) &= 0
 \end{aligned}$$

From the first equation

$$b_3(x) = c_3$$

From the second equation

$$\begin{aligned}
 b_2' &= 1 - 3 c_3 t' \\
 b_2(x) &= x - 3 c_3 t + c_2
 \end{aligned}$$

Since  $t = \log(x)$ , we conclude that  $c_3 = 0$ .

From the third equation

$$\begin{aligned}
 2(x + c_2) t' + b_1' &= 2 \\
 2 + 2 c_2 t' + b_1' &= 2 \\
 2 c_2 t' + b_1' &= 0 \\
 2 c_2 t + b_1 &= c_1
 \end{aligned}$$

Since  $t = \log(x)$ , we conclude that  $c_2 = 0$ .

From the last equation

$$\begin{aligned}
 c_1 t' + b_0' + \frac{d}{dx} \sum_{k=1}^n d_k * \log(u_k) &= 0 \\
 c_1 t + b_0 + \sum_{k=1}^n d_k * \log(u_k) &= c_0
 \end{aligned}$$

Equating coefficients, we obtain  $c_1 = 0$  and  $b_0 = c_0$  and all  $d_k = 0$

Putting all these together

$$\begin{aligned}
 \int t^2 + 2 t dx &= b_3 t^3 + b_2 t^2 + b_1 t + b_0 + \sum_{k=1}^n d_k * \log(u_k) \\
 \int \log^2(x) + 2 \log(x) dx &= x t^2 = x \log^2 x
 \end{aligned}$$

## Logarithmic Part

To integrate

$$\int \frac{p(\theta)}{q(\theta)}$$

we use Rothstein-Trager's algorithm:

$$\int \frac{p}{q} = \sum_{k=1}^n c_k * \log(d_k)$$

where  $c_k$  are the distinct roots of

$$R(z) = \text{res}_{\theta}(q(\theta), p(\theta) - z * q'(\theta))$$

and  $d_k$  are the polynomials

$$d_k = \text{GCD}(q(\theta), p(\theta) - c_k * q'(\theta))$$

In a case of transcendental extension, roots of  $R(z)$  might depend not only of  $z$  but  $x$  as well. In the former case, we say that the integral cannot be done in elementary functions.

### ■ Example.

Consider

$$\int \frac{\log(x) - 1}{\log^2(x) - x^2} dx$$

We extend  $Q(x)$  by  $t = \log(x)$ . Then the integrand can be rewritten as

$$\frac{\log(x) - 1}{\log^2(x) - x^2} = \frac{t - 1}{t^2 - x^2}$$

Compute a resultant of two polynomials (note that  $t$  is a function of  $x$ )

$$t^2 - x^2 \quad \text{and} \quad t - 1 - z \frac{d}{dx} (t^2 - x^2)$$

Since

$$\frac{d}{dx} (t(x)^2 - x^2) = 2 t(x) t'(x) - 2 x = \frac{2 t}{x} - 2 x$$

we get

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Resultant[t^2 - x^2, t - 1 - z (2 t / x - 2 x), t] // Simplify
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$$(-1 + x^2) (-1 + 4 z^2)$$

Find roots wrt z

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Solve[% == 0, z]
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$$\left\{ \left\{ z \rightarrow -\frac{1}{2} \right\}, \left\{ z \rightarrow \frac{1}{2} \right\} \right\}$$

Find GCDs:

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PolynomialGCD[t^2 - x^2, t - 1 - 1/2 (2 t / x - 2 x)]
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$$1 + \frac{t}{x}$$

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PolynomialGCD[t^2 - x^2, t - 1 + 1/2 (2 t / x - 2 x)]
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$$1 - \frac{t}{x}$$

Therefore, the integral is

$$\frac{1}{2} \log\left(\frac{t}{x} + 1\right) - \frac{1}{2} \log\left(1 - \frac{t}{x}\right)$$

or

$$\int \frac{\log(x) - 1}{\log^2(x) - x^2} dx = \frac{1}{2} \log\left(\frac{\log(x)}{x} + 1\right) - \frac{1}{2} \log\left(1 - \frac{\log(x)}{x}\right)$$

### ■ Example.

Consider

$$\int \frac{1}{x + \log(x)} dx$$

The integrand is  $\frac{1}{x+t}$ , where  $t = \log(x)$  in the field  $Q(x, t)$ . Compute a resultant of two polynomials

$$t + x \quad \text{and} \quad 1 - z \frac{d}{dx} (x + t(x))$$

which is

$$\mathbf{Resultant} \left[ \mathbf{x + t}, \mathbf{1 - z \left( 1 + \frac{1}{x} \right)}, \mathbf{t} \right]$$

$$- \frac{-x + z + xz}{x}$$

Find roots

$$\mathbf{Solve} [\% == 0, \mathbf{z}]$$

$$\left\{ \left\{ z \rightarrow \frac{x}{1+x} \right\} \right\}$$

Roots are NOT free of  $x$ , so the integral is not elementary.

## Tower of Extensions

Given an integrand

$$f \in K(x, \theta_1, \theta_2, \dots, \theta_n)$$

where each  $\theta_k$  is transcendental. The integrand may be manipulated as the rational functions of  $\theta_k$ . Let us choose the last extension  $\theta = \theta_n$

$$f(\theta) = \frac{p(\theta)}{q(\theta)} \in F_{n-1}(\theta) = K(x, \theta_1, \theta_2, \dots, \theta_{n-1})$$

and integrate  $f$  with respect to  $\theta_n$ . The algorithm is recursive, namely, when we integrate  $f(\theta)$  we will be recursively invoke integration in the lower field  $F_{n-1}$ .

### ■ Example

Compute the integral

$$\int \left[ \frac{2 \log(x+1)}{x} + \log(x) \left( \log(x+1) + \frac{1}{x+1} \right) \right] dx$$

Extensions

$$t_1 = \log(x), \quad t_2 = \log(x+1)$$

The integral

$$\int \frac{2 t_2}{x} + t_1 \left( t_2 + \frac{1}{x+1} \right) dx$$

It is polynomial in  $t_1$

$$\int \frac{2 t_2}{x} + t_1 \left( t_2 + \frac{1}{x+1} \right) = b_2 t_1^2 + b_1 t_1 + b_0 + R \quad (1)$$

where  $R$  is free of  $t_1$ :

$$R = \sum_{k=1}^n d_k * \log(u_k)$$

We differentiate both sides of (1), to get

$$\frac{2 t_2}{x} + t_1 \left( t_2 + \frac{1}{x+1} \right) = b_2' t_1^2 + (b_1' + 2 b_2 t_1') t_1 + b_0' + b_1 t_1' + \frac{d}{d x} R$$

Equating coefficients by  $t_1$ , we get the following system of equations

$$b_2' = 0$$

$$2 b_2 t_1' + b_1' = \frac{1}{x+1} + t_2$$

$$b_0' + b_1 t_1' + \frac{d}{d x} R = \frac{2 t_2}{x}$$

From the first equation, we find

$$b_2 = c_2$$

From the second equation, we find

$$2 c_2 t_1 + b_1 = \int \left( \frac{1}{x+1} + t_2 \right)$$

Apply the algorithm recursively (we skip this step) to evaluate the integral on the right side. We obtain

$$2 c_2 t_1 + b_1(x) = x t_2 - x + 2 t_2 + c_1$$

We conclude

$$c_2 = 0$$

$$b_1(x) = x t_2 - x + 2 t_2 + c_1$$

From the last equation

$$b_0' + b_1 t_1' + \frac{d}{dx} R = \frac{2t_2}{x}$$

$$b_0' + \frac{x t_2 - x + 2 t_2 + c_1}{x} - \frac{2t_2}{x} = -\frac{d}{dx} R$$

$$b_0 + \int \frac{t_2 x - x + c_1}{x} = -R$$

Again, we need to apply the algorithm recursively to evaluate the above integral

$$b_0 + t_2 x - 2 x + t_2 + t_1 c_1 = -R$$

Since  $R$  is free of  $t_1$ , we conclude that  $c_1 = 0$ . We also set  $b_0$  to  $c_2$ .

Summing all the above,

$$\int \frac{2t_2}{x} + t_1 \left( t_2 + \frac{1}{x+1} \right) = b_2 t_1^2 + b_1 t_1 + b_0 + R \quad (2)$$

$$\int \frac{2t_2}{x} + t_1 \left( t_2 + \frac{1}{x+1} \right) = (x t_2 - x + 2 t_2) t_1 - (t_2 x - 2 x + t_2)$$

or

$$\int \frac{2 \log(x+1)}{x} + \log(x) \left( \log(x+1) + \frac{1}{x+1} \right) dx$$

=

$$(x \log(x+1) - x + 2 \log(x+1)) \log(x) - x \log(x+1) - \log(x+1) + 2x$$

## References

- [1]. M. Bronstein, *Symbolic Integration - Transcendental Functions*, Algorithms and Computations in Mathematics, 2nd edition, Vol 1, Springer-Verlag, 2005