

Gröbner Bases

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Graph Coloring with Gröbner bases

Theorem. *The problem of computing a groebner basis is NP Complete.*

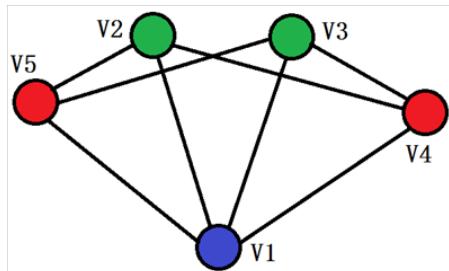
Show it's in NP. For a given finite set $G = \{g_1, g_2, \dots, g_n\}$, we just need to use Buchberger's S-pair criterion: check if for any i and j , the remainder of division $S(g_i, g_j)$ by G is zero. This means, we need to perform the division algorithm $O(n^2)$ times. The division algorithm can be completed in polynomial time, so the whole check will finish in polynomial time.

Show it's NP hard. We should reduce a known NP complete problem to this problem in polynomial time. For simplicity, we choose graph 3-coloring problem.

For a given graph $G = \langle V, E \rangle$, we set $n = |V|$ variables named x_1, \dots, x_n . We choose three colors as roots of the cubic equation $x_i^3 - 1$ for each vertex x_i . Two vertices on an edge must have different colors $x_i^3 = x_j^3$. This translates $x_i^3 - x_j^3 = (x_i - x_j)(x_i^2 + x_i x_j + x_j^2) = 0$. Therefore, we assign $x_i^2 + x_i x_j + x_j^2$ for each edge in E .

Next we compute the reducece Gröbner basis. If the basis is $\langle 1 \rangle$ - there is no coloring. Otherwise, we choose color accordingly.

Example.



```
In[8]:= A = {{0, 1, 1, 1, 1}, {1, 0, 0, 1, 1}, {1, 0, 0, 1, 1}, {1, 1, 1, 0, 0}}
```

```
In[27]:= basis = {}; n = 5;
For[i = 1, i <= n, i++, AppendTo[basis, x[i]^3 - 1]]
```

```
In[29]:= For[i = 1, i < n, i++,
  For[j = i + 1, j <= n, j++,
    If[A[[i, j]] == 1, AppendTo[basis, x[i]^2 + x[i] x[j] + x[j]^2]]]]
```

```
In[77]:= basis
```

```
Out[77]= {-1 + x[1]^3, -1 + x[2]^3, -1 + x[3]^3, -1 + x[4]^3,
-1 + x[5]^3, x[1]^2 + x[1] x[2] + x[2]^2, x[1]^2 + x[1] x[3] + x[3]^2,
x[1]^2 + x[1] x[4] + x[4]^2, x[1]^2 + x[1] x[5] + x[5]^2, x[2]^2 + x[2] x[4] + x[4]^2,
x[2]^2 + x[2] x[5] + x[5]^2, x[3]^2 + x[3] x[4] + x[4]^2, x[3]^2 + x[3] x[5] + x[5]^2}
```

```
In[78]:= gb = GroebnerBasis[basis, Array[x, n]]
```

```
Out[78]= {-1 + x[5]^3, x[4] - x[5],
x[3]^2 + x[3] x[5] + x[5]^2, x[2] - x[3], x[1] + x[3] + x[5]}
```

The reduce Gröbner basis: $\{x_5^3 - 1, x_4 - x_5, x_3^2 + x_3 x_5 + x_5^2, x_2 - x_3, x_1 + x_3 + x_5\}$

■ Graph Coloring Ideal

Define

$$V_{n,k} = \langle x_i^k - 1, x_i \in V \rangle$$

$$E_{n,k} = \langle x_i^{k-1} + x_i^{k-2} x_j + \dots + x_i x_j^{k-2} + x_j^{k-1}, (x_i, x_j) \in E \rangle$$

Then $V_{n,k} + E_{n,k}$ is called the k-coloring ideal.

An advantage of Gröbner bases is that it allows us to decide if a graph G is uniquely k -colorable up to the permutation.

Theorem. A graph G with n vertices $x_1, x_2, x_3, z_1, z_2, \dots, z_{n-3}$ is uniquely 3-colorable iff a correspondent Gröbner basis has the following form

$$\langle x_3^3 - 1, x_2^2 + x_2 x_3 + x_3^2, x_1 + x_2 + x_3, f(z_1), \dots, f(z_{n-3}) \rangle$$

where

$$f(z_i) = \begin{cases} z_i - x_j & j = 2, 3, \text{ and } z_i \text{ has the same color as } x_j \\ z_i + x_2 + x_3 & \text{if } z_i \text{ has the same color as } x_1 \end{cases}$$

The uniqueness follows from the fact that the first three equations in the Groebner basis uniquely define colors for three vertices, the other linear equations $f(z_i)$ uniquely define the rest.

■ Demo

```
In[21]:= ChooseColor[x_] :=
```

```

Module[{}, Switch[x, 1, Red, -(-1)^1/3, Green, -1 + (-1)^1/3, Blue,
_, Black]];

GraphColoring[g_Graph] := Module[
{i, j, n, rg, A, basis = {}, gb, tmp, ansRule = {}, ansList = {}, cg},
(* rg: random generated graph; cg: color graph *)

rg = g;
n = VertexCount[g];

cg = Graph[Range[n], EdgeList[rg], VertexLabels → Table[i → vi, {i, n}],
VertexStyle → Black];

(* Get the adjacency matrix *)
A = Normal[AdjacencyMatrix[rg]];

(* Construct the 3-coloring ideal *)
For[i = 1, i ≤ n, i++, AppendTo[basis, x[i]^3 - 1]];
For[i = 1, i < n, i++,
For[j = i + 1, j ≤ n, j++,
If[A[[i, j]] == 1, AppendTo[basis, x[i]^2 + x[i] x[j] + x[j]^2]];
]
];
(* I guess there is a more elegant way... *)

gb = GroebnerBasis[basis, Array[x, n]];
(* calculate the Groebner basis of 3-coloring ideal wrt lex order *)
If[TrueQ[gb[[1]] == 1], Print["No solution"];
Return[{rg, A, gb, ansRule, ansList, cg}]];
(* The graph is not 3-colorable iff 3-coloring ideal = <1> *)

tmp = gb;
For[i = 1;
j = n, j ≥ 1 && i ≤ Length[gb], i++,
(* Use i to pick polynomials. Use j to pick variables. *)
If[TrueQ[tmp[[i]] == 0], (* If tmp[[i]] is 0, do nothing;
else solve x[j] and back substitute x[j] to GB. Always
choose the first root. *)
,
rule = First[Solve[tmp[[i]] == 0, x[j]]];
tmp = Simplify[tmp /. rule];
PrependTo[ansRule, rule];
PrependTo[ansList, (rule /. Rule → List)[[1, 2]] // FullSimplify]; j-- ]
];

```

```

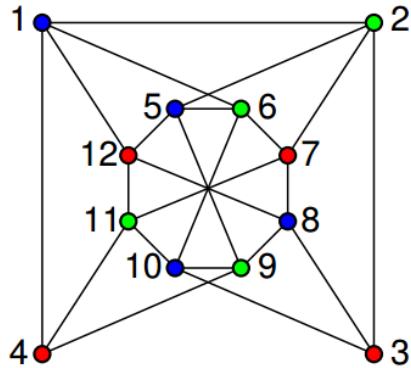
];
cg = Graph[Range[n], EdgeList[rg], VertexLabels → Table[i → vi, {i, n}],
    VertexStyle → Table[i → ChooseColor[ansList[[i]]], {i, n}]];
{rg, A, gb, ansRule, ansList, cg}
]

```

In[23]:= `Graph1 = Graph[Range[9], {1 → 2, 1 → 4, 1 → 5, 1 → 6, 2 → 7, 2 → 3,
 2 → 5, 3 → 4, 3 → 6, 3 → 9, 5 → 6, 6 → 8, 7 → 8, 8 → 9},
 VertexLabels → Table[i → vi, {i, 9}]];`

In[24]:= `GraphColoring[Graph1]`

In 2001 Chao provided a 12-vertex 22-edges graph that uniquely colored.

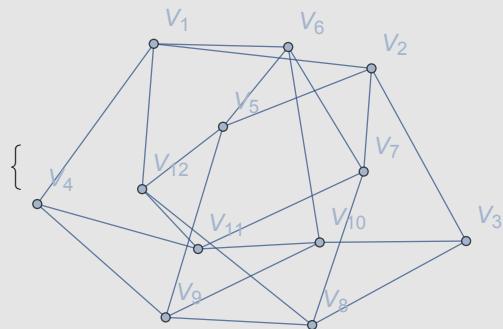


In[25]:= `ChaoChenGraph =
 Graph[Range[12], {1 → 2, 1 → 6, 1 → 12, 1 → 4, 2 → 5, 2 → 7, 2 → 3,
 3 → 8, 3 → 10, 4 → 9, 4 → 11, 5 → 6, 6 → 7, 7 → 8, 8 → 9, 9 → 10,
 10 → 11, 11 → 12, 12 → 5, 5 → 9, 6 → 10, 7 → 11, 8 → 12},
 VertexLabels → Table[i → vi, {i, 12}]];`

In[79]:=

GraphColoring [ChaoChenGraph]

Out[79]=



References

- [1] C. J. Hillar and D. Windfeldt, *Algebraic Characterization of Uniquely vertex Colorable Graphs*, J. Comb. Theory., 98(2), 2008, 400-414.