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## 1. Primitive Recursion (50)

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### Background

Recall that [bounded search](#) is defined as follows. Let  $g : \mathbb{N}^{n+1} \rightarrow \mathbb{N}$ . Then  $f = \text{BS}[g] : \mathbb{N}^{n+1} \rightarrow \mathbb{N}$  is the function defined by

$$f(x, \mathbf{y}) = \begin{cases} \min(z < x \mid g(z, \mathbf{y}) = 0) & \text{if } z \text{ exists,} \\ x & \text{otherwise.} \end{cases}$$

### Task

- A. Show that  $f(x, \mathbf{y}) = \sum_{z < h(x)} g(z, \mathbf{y})$  is primitive recursive when  $h$  is primitive recursive and strictly monotonic.
- B. What if  $h$  is not monotonic?
- C. Show that  $\text{BS}[g]$  is primitive recursive whenever  $g$  is.

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## 2. A Recursion (50)

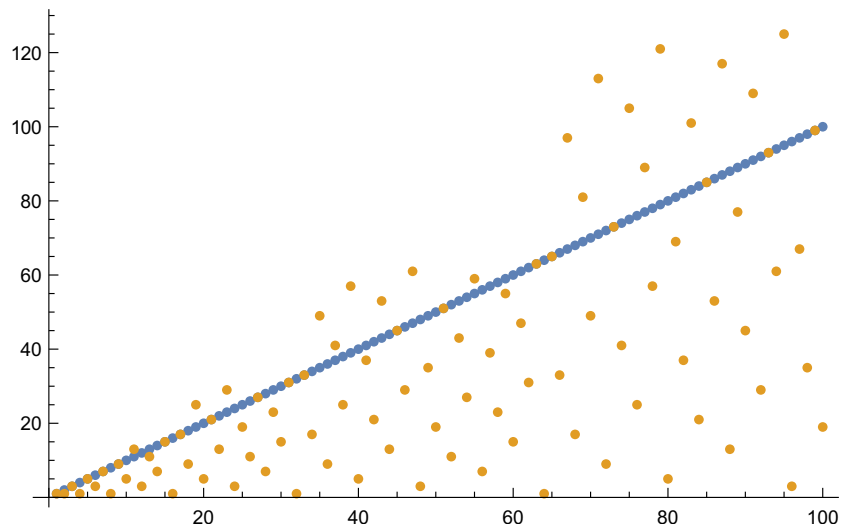
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### Background

Consider the following function  $f$ , presumably defined on the positive integers.

$$\begin{aligned}f(1) &= 1 \\f(3) &= 3 \\f(2n) &= f(n) \\f(4n + 1) &= 2f(2n + 1) - f(n) \\f(4n + 3) &= 3f(2n + 1) - 2f(n)\end{aligned}$$

For what it's worth, here is a plot of the first few values.



### Task

- Consider small values of  $f$  and conjecture an explicit, non-recursive definition of  $f$ .
- Prove that your definition is correct and conclude that  $f$  is indeed a function from  $\mathbb{N}_+$  to  $\mathbb{N}_+$ .
- Is  $f$  primitive recursive?

### Comment

Implement  $f$  and experiment.

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### 3. Primitive Recursive Word Functions (40)

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#### Background

We defined primitive recursive functions on the naturals. A similar definition would also work for words over some alphabet  $\Sigma$ . We write  $\varepsilon$  for the empty word and  $\Sigma^*$  for the set of all words over  $\Sigma$ . Consider the clone of word functions generated by the basic functions

- **Constant empty word**  $E : (\Sigma^*)^0 \rightarrow \Sigma^*$ ,  $E() = \varepsilon$ ,
- **Append functions**  $S_a : \Sigma^* \rightarrow \Sigma^*$ ,  $S(x) = xa$  where  $a \in \Sigma$ .

and closed under **primitive recursion** over words: suppose we have a function  $g : (\Sigma^*)^n \rightarrow \Sigma^*$  and a family of functions  $h_a : (\Sigma^*)^{n+2} \rightarrow \Sigma^*$ , where  $a \in \Sigma$ . We can then define a new function  $f : (\Sigma^*)^{n+1} \rightarrow \Sigma^*$  by

$$\begin{aligned} f(\varepsilon, \mathbf{y}) &= g(\mathbf{y}) \\ f(xa, \mathbf{y}) &= h_a(x, f(x, \mathbf{y}), \mathbf{y}) \quad a \in \Sigma \end{aligned}$$

We will call the members of this clone the **word primitive recursive (w.p.r.)** functions.

#### Task

1. Show that the reversal operation  $\text{rev}(x) = x_n x_{n-1} \dots x_1$  is w.p.r.
2. Show that the prepend operations  $\text{pre}_a(x) = ax$  are w.p.r.
3. Show that the concatenation operation  $\text{cat}(x, y) = xy$  is w.p.r.
4. Prove that every primitive recursive function is also a word primitive recursive function. By this we mean that for every p.r. function  $f : \mathbb{N}^k \rightarrow \mathbb{N}$  there is a w.p.r. function  $F : (\Sigma^*)^k \rightarrow \Sigma^*$  so that  $f(\mathbf{x}) = D(F(C(\mathbf{x})))$  where  $C$  and  $D$  are simple coding and decoding functions (between numbers and words).
5. Prove the opposite direction: every w.p.r. is already p.r., using coding and decoding as in the last problem.

#### Comment

For the last part, don't get bogged down in tons of technical details, just explain how one would go about proving this.